## Errata for the

## Amiga Hardware Manual

The attached documents contain various updates to the Amiga Hardware Manual. The documents are:

Changes and Additions

Read this first; it contains miscellaneous corrections and is a guide to all of the updates.

Examples

Contains example assembly code sequences for the Copper, playfields, and sprites.

Addendum to Chapter 5: Audio Hardware Describes the audio state machine.

Addendum to Chapter 6: Blitter Hardware Describes blitter DMA and how the operation of the blitter affects the performance of the rest of the system.

Appendix F This replaces the present Appendix F.

Appendix G This replaces the present Appendix G.

## Changes and Additions to the Amiga Hardware Manual

Changes and additions have been made to the following chapters and appendixes of the Amiga Hardware Manual.

#### Chapter 2: Coprocessor Hardware

Add the examples in the attached document called "Examples for the Amiga Hardware Manual".

## Chapter 3: Playfield Hardware

Add the examples in the attached document called "Examples for the Amiga Hardware Manual".

## Chapter 4: Sprite Hardware

Add the examples in the attached document called "Examples for the Amiga Hardware Manual".

On page 4-14, under section 4.3, "Displaying a Sprite", add the following:

#### CAUTION

Should sprite DMA be turned off while a sprite is being displayed (that is, after VSTART but before VSTOP), the system will continue to display the line of sprite data that was most recently fetched. This causes a vertical bar to appear on the screen. It is recommended that sprite DMA only be turned off during vertical blanking or during some portion of the screen where you are *sure* that no sprite is being displayed.

## Chapter 5: Audio Hardware

Add the attached document called "Addendum to Chapter 5: Audio Hardware".

#### Chapter 6: Blitter Hardware

Add the attached document called "Addendum to Chapter 6: Blitter Hardware".

On page 6-14, under section 6.11, "Line Drawing", add the following line to the table:

Register Name	Bit Number	Bit <u>Name</u>	State	Purpose
BLTCON1	5		0	Reserved for flabloden mode

On page 6-8, add the following to section 6.6, "Ascending and Descending Addressing":

If the source and destination data areas overlap in a blitter operation, there is a possibility of writing to a particular location as the destination before it was read as the source.

To prevent this kind of data destruction, you must take care to correctly choose ascending or descending mode, and possibly you should offset the source or destination.

Some types of overlapped blitter moves will always result in damaged data; therefore, the desired move must be done in two steps:

- 1. Move to a safe place.
- 2. Move to desired destination.

Using the blitter Memory Bus Activity Table<sup>1</sup>, you can observe the order of operations and determine the required offset or mode. Pay careful attention to the notes. It helps to draw pictures.

<sup>&</sup>lt;sup>1</sup> This table is located in the separate "Addendum to the Blitter Hardware Chapter" attached to this document.

### Chapter 8: Interface Hardware

On page 8-4, under "Counter Limitations", the last line should read:

\*\* 45-200 = -155. Because the absolute value exceeds 127, the true count must be

## Appendix A

On page 4, in the table called "LINE MODE (line draw)", change the following bits for BLTCON1:

BIT#	BLTCON0	BLTCON1	
15 14 13 13	START3 START2 START1 START0	TEXTURE3 TEXTURE2 TEXTURE1 TEXTURE0	
05	LF5	0	(Reserved for new mode)

On page 5, in the table called "LINE MODE (line draw)", change the line beginning "OVF ..." to: 0 Reserved for new mode

At the end of the same table, add:

The "B" source is used for texturing the drawn lines.

#### Appendix B

On page 2, after "BLTCDAT", "BLTBDAT", and "BLTADAT", change "&" to "%".

## Appendix F

Remove Appendix F and insert the attached new Appendix F. Appendix F contains a new front section, called "Brief address map for 8520's".

## Appendix G

Remove Appendix G and insert the attached new Appendix G. Appendix G has some minor changes.

Since this appendix is subject to change, anyone who is actually designing a device to interface to the Amiga should contact Amiga Third-Party Support.

## Appendix I

This is a new appendix. It describes the interface to the 68000 bus connector.

## Appendix F

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## Appendix I

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## Examples for the

## Amiga Hardware Manual

#### Chapter 2: Coprocessor Hardware

Add the following example to page 2-4, under section 2.3, "The MOVE Instruction".

The following MOVE instructions point bit-plane pointer 1 at \$21000 and bitplane pointer 2 at \$25000.

DC.W	\$00E0,\$0002	;MOVE \$0002 TO ADDRESS \$0E0 (BPL1PTH)
DC.W	\$00E2,\$1000	;MOVE \$1000 TO ADDRESS \$0E2 (BPL1PTL)
DC.W	\$00E4,\$0002	;MOVE \$0002 TO ADDRESS \$0E4 (BPL2PTH)
DC.W	\$00E6,\$5000	;MOVE \$5000 TO ADDRESS \$0E6 (BPL2PTL)

Add the following examples to page 2-5, under section 2.4, "The Wait Instruction".

This first WAIT instruction waits for scan line 150 (96 hex) with the horizontal position masked off.

DC.W \$9601,\$FF00 ;WAIT FOR LINE 150, IGNORE HORIZONTAL COUNTERS

This second WAIT instruction waits for scan line 255 and horizontal position 254. This will never occur, so the copper stops until the next vertical blanking interval begins.

DC.W FFFF,FFFE; WAIT FOR LINE 255, H = 254 (ENDS COPPER LIST).

Add the following example to page 2-10, under section 2.6, "Putting Together a Copper Instruction List".

This is an example of a complete Copper list. It is a Copper list for two bit-planes, one at \$21000, one at \$25000. At the top of the screen, the color registers are set as follows:

COLOR00 = WHITE COLOR01 = RED COLOR02 = GREENCOLOR03 = BLUE

At line 150, the color registers are reloaded as follows:

 $\bullet$ COLOR00 = BLACK COLOR01 = YELLOWCOLOR02 = CYANCOLOR03 = MAGENTA**COPPERLIST:** DC.W \$00E0,\$0002 ;MOVE \$0002 INTO ADDRESS \$0E0 (BPL1PTH) DC.W \$00E2,\$1000 :MOVE \$1000 INTO ADDRESS \$0E2 (BPL1PTL) DC.W \$00E4,\$0002 ;MOVE \$0002 INTO ADDRESS \$0E4 (BPL2PTH) DC.W \$00E6,\$5000 ;MOVE \$5000 INTO ADDRESS \$0E6 (BPL2PTL) ; Load color registers DC.W \$0180,\$0FFF ;MOVE WHITE INTO ADDRESS \$180 (COLOR00) DC.W \$0182,\$0F00 ;MOVE RED INTO ADDRESS \$182 (COLOR01) DC.W \$0184,\$00F0 ;MOVE GREEN INTO ADDRESS \$184 (COLOR02) DC.W ;MOVE BLUE INTO ADDRESS \$186 (COLOR03) \$0186,\$000F : Wait for line 150 DC.W \$9601,\$FF00 ;WAIT FOR LINE 150, IGNORE HORIZ. POSITION ; Reload color registers DC.W \$0180,\$0000 ;MOVE BLACK INTO ADDRESS \$0180 (COLOR00) DC.W \$0182,\$0FF0 ;MOVE YELLOW INTO ADDRESS \$0182 (COLOR01) DC.W \$0184,\$00FF ;MOVE CYAN INTO ADDRESS \$0184 (COLOR02) DC.W \$0186,\$0F0F ;MOVE MAGENTA INTO ADDRESS \$0186 (COLOR03) ; End copper list by waiting for the impossible DC.W **\$FFFF**, **\$FFFE** ; WAIT FOR LINE 255, H = 254 (NEVER HAPPENS)

Add the following to page 2-13, under section 2.8.1, "The SKIP Instruction".

The following SKIP instruction will skip the instruction following it if VP >= 100 (\$64).

DC.W \$6401,\$FF01; IF VP >= 100,SKIP NEXT INSTR (IGNORE HP). DC.W ....this is the instruction that will be skipped....

Replace the example on page 2-14, under section 2.8.2, "Copper Loops and Branches and Comparison Enable", with the following:

: Copper list to interrupt 68000 once every 16 scan lines, ; in the range VP = 80 through VP = 160. ; DC.W \$5001,\$FFFE ;WAIT FOR VP =\$50, HP = 0\$0F01,\$0F00 ;WAIT FOR VP = xxxx1111DC.W ; The following instruction writes to address \$09C, the ; interrupt request register. Writing \$8010 sets the copper ; interrupt bit in the register, which will interrupt the 68000. DC.W \$009C,\$8010 ;MOVE \$8010 TO \$09C (INTERRUPT 68000) DC.W \$A001,\$FF01 ;SKIP NEXT INSTRUCTION IF VP >= 160; The next MOVE instruction doesn't actually do a move. It forces ; the copper to jump to the address in COP2LC. This must have been ; previously set by either the copper or the 68000. If VP >= 160, ; then this instruction will be skipped. ; DC.W \$008A.\$0000 ;MOVE 0 TO COPJMP2 (COP2LC PREVIOUSLY SET) DC.W ... other copper instructions

#### Chapter 3: Playfield Hardware

Add this example to page 3-10, under "Selecting Number of Bit-Planes".

This example shows how to write to the BPLCON0 register to tell the system to use 2 bitplanes.

; ; Writing \$2200 to BPLCON0 sets the f	ollowing conditions:	
; Low Resolution		
; Use two bitplanes		
; Hold-and-modify $mode = OFF$		
; Single playfield mode		
; Composite video color enabled		
; Genlock audio disabled		
; Light pen disabled		
; Interlace disabled		
; External resync disabled		
;		
BPLCON0 EQU \$DFF100	;BPLCON0 ADDRESS	
MOVE.W #\$2200,BPLCON0	;WRITE TO IT	

Add the following example to page 3-10, under section 3.2.3, "Selecting Horizontal and Vertical Resolution".

This example shows how to set the HIRES and LACE bits.

Writing \$A204 to BPLCON0 sets the following conditions:
High resolution
Use two bitplanes
Hold-and-modify mode = OFF
Single playfield mode
Composite video color enabled
Genlock audio disabled
Light pen disabled
Interlace enabled
External resync disabled
BPLCON0 EQU \$DFF100 ;BPLCON0 ADDRESS
MOVE.W #\$A204,BPLCON0 ;WRITE TO IT

Add the following example to page 3-12, under section 3.2.4, "Allocating Memory for Bit-Planes"

This example shows how to set pointers for memory allocation for bit-planes.

	t \$21000 and the other at \$25000, the processor PL2PT to \$25000. Normally, this is the copper's
; BPL1PTH EQU \$DFF0E0 BPL1PTL EQU \$DFF0E2 BPL2PTH EQU \$DFF0E4 BPL2PTL EQU \$DFF0E6	;HIGH THREE BITS OF BITPLANE 1 POINTER REGISTE ;LOW FIFTEEN BITS ;HIGH THREE BITS OF BITPLANE 2 POINTER REGISTE ;LOW FIFTEEN BITS
; a single longword move. The	o write to a bitplane pointer, but we can use 68000 writes the high order word to the lower- ) and the low order word to the higher-numbered
; MOVE.L #\$21000,BPL1 MOVE.L #\$25000,BPL2	,

Add the following examples to page 3-15, under "A One- or Two-Color Playfield"

This example shows how to define a 1-color bit-plane.

; ; This code fills a low resolution bit-plane with the background color (COLOR00) ; by writing all 0's into its memory area. The bitplane starts at \$21000 and ; is 8000 bytes long. ; LEA \$21000,A0 ;POINT AT BITPLANE MOVE.W #2000,D0 ;WRITE 2000 LONGWORDS = 8000 BYTES LOOP: MOVE.L #0,(A0)+ ;WRITE OUT A ZERO SUBQ.W #1,D0 ;DECREMENT COUNTER

;LOOP UNTIL BITPLANE IS FILLED WITH 0'S

This example shows how to define a 2-color bit-plane.

BNE LOOP

; This code is identical to the last example, except the bitplane is filled ; with \$FF00FF00 instead of all 0's. This will produce two colors.

LEA \$21000,A0 ;POINT AT BITPLANE MOVE.W #2000,D0 ;WRITE 2000 LONGWORDS = 8000 BYTES LOOP: MOVE.L #\$FF00FF00,(A0)+ ;WRITE OUT \$FF00FF00 SUBQ.W #1,D0 ;DECREMENT COUNTER BNE LOOP ;LOOP UNTIL BITPLANE IS FULL

Add this example to page 3-18, under "Setting the Display Window Starting Position".

Setting DIWSTRT for the basic playfield:

; This code sets DIWSTRT for a basic playfield. We write \$2C for the vertical ; position and \$81 for the horizontal position.

DIWSTRT EQU \$DFF08E ;DISPLAY WINDOW START REGISTER ADDRESS

MOVE.W #\$2C81,DIWSTRT ;WRITE IT OUT

Add this example to page 3-19, under "Setting the Display Window Stopping Position".

Setting DIWSTOP for the basic playfield:

; This code sets DIWSTOP for a basic playfield. We write \$F4 for ; the vertical position and \$C1 for the horizontal position.

DIWSTOP EQU \$DFF090 ;DISPLAY WINDOW STOP REGISTER ADDRESS ; MOVE.W #\$F4C1,DIWSTOP ;WRITE IT OUT

Add the following examples to page 3-19, under Section 3.2.7, "Telling the System How to Fetch and Display Data".

Setting data fetch for the basic playfield:

; This code sets the data fetch start and stop values for a basic ; playfield. We write \$0038 to DDFSTRT and \$00D0 into DDFSTOP. ; DDFSTRT EQU \$DFF092 DDFSTOP EQU \$DFF094 ; MOVE.W #\$0038,DDFSTRT ;WRITE TO DDFSTRT MOVE.W #\$00D0,DDFSTOP ;WRITE TO DDFSTOP

Setting modulo for the basic playfield:

; This code sets the modulo to 0 for a low resolution playfield with one ; bitplane. For this example, the bitplane is one of the odd-numbered ; ones. We would write to BPL2MOD for even-numbered bitplanes. ; BPL1MOD EQU \$DFF108 ;MODULO FOR ODD BITPLANES ; MOVE.W #0,BPL1MOD ;SET MODULO TO 0

Add the following examples to page 3-26, under a new section, 3.2.11, "Complete Example Playfields".

; This first example sets up a 320 x 200 pixel playfield with ; one bitplane. ; The bitplane lives at \$21000. ; We also set up a copper list at \$20000. CUSTOM EQU \$DFF000 BPLCON0 EQU \$100 BPLCON1 EQU \$102 BPLCON2 EQU \$104 BPL1MOD EQU \$108 DDFSTRT EQU \$092 DDFSTOP EQU \$094 DIWSTRT EQU \$08E DIWSTOP EQU \$090 VPOSR EQU \$004 COLOR00 EQU \$180

COLOR01 EQU \$182

```
COLOR02 EQU $184
COLOR03 EQU $186
DMACON EQU $096
                         COPPER LOCATION REGISTER 1 (HIGH 3 BITS)
COP1LCH EQU $080
                             ;A0 POINTS AT CUSTOM CHIPS
    LEA CUSTOM, A0
    MOVE.W #$1200,BPLCON0(A0)
                                 ONE BITPLANE, ENABLE COMPOSITE COLOR
    MOVE.W #0,BPLCON1(A0)
                                 SET HORIZONTAL SCROLL VALUE TO 0
                                 SET MODULO TO 0 FOR ALL ODD BITPLANES
    MOVE.W #0,BPL1MOD(A0)
    MOVE.W #$0038,DDFSTRT(A0)
                                 SET DATA FETCH START TO $38
    MOVE.W #$00D0,DDFSTOP(A0)
                                 SET DATA FETCH STOP STOP TO $D0
                                 SET DIWSTRT TO $2C81
    MOVE.W #$2C81,DIWSTRT(A0)
    MOVE.W #$F4C1,DIWSTOP(A0)
                                 SET DISPLAY WINDOW STOP TO $F4C1
    MOVE.W #$0F00,COLOR00(A0)
                                 SET BACKGROUND COLOR TO RED
    MOVE.W #$0FF0,COLOR01(A0)
                                  SET COLOR REGISTER 1 TO YELLOW
; FILL BITPLANE WITH $FF00FF00 TO PRODUCE STRIPES
                             POINT AT BEGINNING OF BITPLANE
    MOVE.L #$21000,A1
    MOVE.L #$FF00FF00,D0
                             ;WE WILL WRITE $FF00FF00 LONG WORDS
    MOVE.W #2000,D1
                              2000 \text{ LONG WORDS} = 8000 \text{ BYTES}
        MOVE.L D0(A1)+
                              WRITE A LONG WORD
LOOP:
                              DECREMENT COUNTER
    SUBQ.W #1,D1
                              :LOOP UNTIL BITPLANE IS FILLED
    BNE LOOP
; Set up copper list at $20000
                              POINT AT COPPER LIST DESTINATION
    MOVE.L #$20000,A1
    LEA COPPERL, A2
                              POINT A2 AT COPPER LIST DATA
                              ;MOVE A WORD
CLOOP: MOVE.L (A2),(A1)+
             #$FFFFFFE,(A2)+ ;CHECK FOR LAST LONGWORD OF COPPER LIST
    CMPI.L
    BNE CLOOP
                              LOOP UNTIL ENTIRE COPPER LIST IS MOVED
; Point copper at copper list
    MOVE.L #$20000,COP1LCH(A0)
                                  WRITE TO COPPER LOCATION REGISTER
                                  FORCE COPPER TO $20000
    MOVE.W COPJMP1(A0),D0
; Start DMA
     MOVE.W #$8380,DMACON(A0)
                                  ;ENABLE BITPLANE AND COPPER DMA
                          ;GO DO NEXT TASK
     BRA ....
```

; This is the data for the copper list.

## COPPERL:

DC.W	\$00E0,\$0002	;MOVE \$0002 TO ADDRESS \$0E0 (BPL1PTH)
DC.W	\$00E2,\$1000	;MOVE \$1000 TO ADDRESS \$0E2 (BPL1PTL)
DC.W	<b>\$FFFF,\$FFFE</b>	END OF COPPER LIST
200		

;

```
This second example of a basic playfield sets up a hi resolution, interlaced
display with one bitplane.
The equates are the same as the previous example so they aren't repeated here.
                              ;ADDRESS OF CUSTOM CHIPS
    LEA CUSTOM, A0
    MOVE.W #$9204,BPLCON0(A0)
                                  HIRES, 1 BITPLANE, INTERLACE
    MOVE.W #0,BPLCON1(A0)
                                  ;HORIZONTAL SCROLL VALUE = 0
                                  MODULO = 80 FOR ODD BITPLANES
    MOVE.W #80,BPL1MOD(A0)
                                  DITTO FOR EVEN BITPLANES
    MOVE.W #80,BPL2MOD(A0)
    MOVE.W #$003C,DDFSTRT(A0)
                                  SET DATA FETCH START FOR HI RES
                                  SET DATA FETCH STOP
    MOVE.W #$00D4,DDFSTOP(A0)
                                  SET DISPLAY WINDOW START
    MOVE.W #$2C81,DIWSTRT(A0)
    MOVE.W #$F4C1,DIWSTOP(A0)
                                   SET DISPLAY WINDOW STOP
; Set up color registers
    MOVE.W #$000F,COLOR00(A0)
                                   ; BACKGROUND COLOR = BLUE
                                   :FOREGROUND COLOR = WHITE
    MOVE.W #$0FFF,COLOR01(A0)
; Set up bitplane at $20000.
;
                              POINT A1 AT BITPLANE
    LEA $20000,A1
                              :A2 POINTS AT CHARACTER DATA
    LEA CHARLIST,A2
                              ;WRITE 400 LINES OF DATA
    MOVE.W #400,D1
                              WRITE 20 LONG WORDS PER LINE
    MOVE.W #20,D0
L1:
                              WRITE A LONG WORD
    MOVE.L (A2),(A1)+
    SUBQ.W #1,D0
                              ;DECREMENT COUNTER
                              LOOP UNTIL LINE IS FULL
    BNE L1
;
    MOVE.W #20,D0
                              RESET LONG WORD COUNTER
                              POINT AT NEXT WORD IN CHAR LIST
    ADDQ.L #4,A2
                              ;END OF CHAR LIST?
             \#$FFFFFFFF,(A2)
    CMPI.L
```

```
BNE L2
                                YES, RESET A2 TO BEGINNING OF LIST
    LEA CHARLIST, A2
L2:
                                DECREMENT LINE COUNTER
    SUBQ.W #1,D1
                                ;LOOP UNTIL ALL LINES ARE FULL
    BNE L1
; Start DMA
     MOVE.W #$8300,DMACON(A0)
                                     ENABLE BITPLANE DMA ONLY, NO COPPER.
: Since this example has no copper list, we sit in a loop waiting for vertical
; blank. When it comes, we check the LOF (long frame) bit in VPOSR. If it
; is a 0, then this is a short frame so we point the bitplane pointers to
; $200050; if LOF = 1, then this is a long frame so we point to $20000. This
; keeps the long and short frames in the right relationship to each other.
VLOOP:
                                      READ INTERRUPT REQUESTS
     MOVE.W INTREQR(A0),D0
              #$0020,D0
                                      MASK OFF ALL BUT VERTICAL BLANK
     AND.W
                                      LOOP UNTIL VERTICAL BLANK COMES
     BEQ VLOOP
                                      RESET VERTICAL INTERRUPT
     MOVE.W #$0020,INTREQ(A0)
                                      READ LOF BIT INTO DO BIT 15
     MOVE.W VPOSR(A0),D0
                                      ; IF LOF = 0, JUMP
     BPL VL1
     MOVE.L #$20000,BPL1PTH(A0)
                                      ;LOF = 1, POINT TO 20000
                                      BACK TO TOP
     BRA VLOOP
VL1:
     MOVE.L #$20050,BPL1PTH(A0)
                                      LOF = 0, POINT TO $20050
                                      ;BACK TO TOP
     BRA VLOOP
 Character list
     DC.L $18FC3DF0,$3C6666D8,$3C66C0CC,$667CC0CC
     DC.L $7E66C0CC,$C36666D8,$C3FC3DF0,$0000000
     DC.L $FFFFFFF
```

Add the following example to page 3-50, under "Specifying the Amount of Delay".

Setting the delay for horizontal scrolling:

This code sets the horizontal scroll delay to 7 for both playfields.

BPLCON1 EQU \$DFF102 ;HORIZONTAL SCROLL REGISTER ;

MOVE.W #\$77,BPLCON1

Add the following example to page 3-52, under section 3.6.2, "Hold and Modify Mode".

Hold and modify example:

```
: This code generates a six bit-plane display with hold-and-modify mode
turned on. We load all 32 color registers with black to prove that
; the colors are being generated by hold-and-modify.
; The equates are the usual, so we won't repeat them here.
; First we set up the control registers.
                                POINT A0 AT CUSTOM CHIPS
    LEA CUSTOM, A0
    MOVE.W #$6A00,BPLCON0(A0)
                                    SIX BITPLANES, HOLD AND MODIFY MODE
    MOVE.W \#0,BPLCON1(A0)
                                    ;HORIZONTAL SCROLL = 0
                                    ;MODULO FOR ODD BITPLANES = 0
    MOVE.W #0,BPL1MOD(A0)
    MOVE.W #0,BPL2MOD(A0)
                                    DITTO FOR EVEN BITPLANES
    MOVE.W #$0038,DDFSTRT(A0)
                                    SET DATA FETCH START
    MOVE.W #$00D0,DDFSTOP(A0)
                                    SET DATA FETCH STOP
                                    SET DISPLAY WINDOW START
    MOVE.W #$2C81,DIWSTRT(A0)
                                    SET DISPLAY WINDOW STOP
    MOVE.W #$F4C1,DIWSTOP(A0)
; Set all color registers = black to prove hold and modify mode is working.
                                ;INITIALIZE COUNTER
    MOVE.W #32,D0
    LEA CUSTOM+COLOR00,A1 ;POINT A1 AT FIRST COLOR REGISTER
CREGLOOP:
                                WRITE BLACK TO A COLOR REGISTER
    MOVE.W \#$0000,(A1)+
                                :DECREMENT COUNTER
    SUBQ.W #1,D0
    BNE CREGLOOP
                                LOOP UNTIL ALL COLOR REGISTERS SET.
; Fill six bitplanes with an easily recognizable pattern.
    MOVE.W #2000,D0
                                ;2000 LONGWORDS PER BITPLANE
    MOVE.L #$21000,A1
                                ;POINT A1 AT BIT PLANE 1
```

```
POINT A2 AT BIT PLANE 2
    MOVE.L #$23000,A2
                               POINT A3 AT BIT PLANE 3
    MOVE.L #$25000,A3
    MOVE.L #$27000,A4
                               POINT A4 AT BIT PLANE 4
                               POINT A5 AT BIT PLANE 5
    MOVE.L #$29000,A5
                               POINT A6 AT BIT PLANE 6
    MOVE.L #$2B000,A6
FPLLOOP:
                               ;FILL BIT PLANE 1 WITH $55555555
    MOVE.L #$55555555,(A1)+
                               FILL BIT PLANE 2 WITH $33333333
    MOVE.L #$33333333(A2)+
    MOVE.L \#$0F0F0F0F,(A3)+
                               FILL BIT PLANE 3 WITH $0F0F0F0F
    MOVE.L \#$00FF00FF,(A4)+
                               FILL BIT PLANE 4 WITH $00FF00FF
    MOVE.L #$FFFFFFF,(A5)+ ;FILL BIT PLANE 5 WITH $FFFFFFF
    MOVE.L #$0000000,(A6)+
                               ;FILL BIT PLANE 6 WITH $00000000
                               :DECREMENT COUNTER
    SUBQ.W #1,D0
                               ;LOOP UNTIL ALL BIT PLANES ARE FULL.
    BNE FPLLOOP
; Set up a copper list at $20000.
                               POINT A1 AT COPPER LIST DESTINATION
    MOVE.L #$20000,A1
                               POINT A2 AT COPPER LIST IMAGE
    LEA COPPERL, A2
CLOOP: MOVE.L (A2),(A1)+
                                MOVE A LONG WORD
             #$FFFFFFE,(A2)+ ;CHECK FOR END OF COPPER LIST
    CMPI.L
                                LOOP UNTIL ENTIRE COPPER LIST MOVED.
    BNE CLOOP
; Point copper at copper list
    MOVE.L #$20000,COP1LCH(A0)
                                    LOAD COPPER JUMP REGISTER
    MOVE.W COPJMP1(A0),D0
                                    FORCE LOAD INTO COPPER P.C.
; Start DMA
     MOVE.W #$8380,DMACON(A0)
                                    ENABLE BITPLANE AND COPPER DMA
     BRA .....next stuff to do.....
; Copper list for six bit planes. Bit plane 1 is at $21000, 2 is at $23000,
; 3 is at $25000, 4 is at $27000, 5 is at $29000, 6 is at $28000.
COPPERL:
     DC.W
                                ;BIT PLANE 1 POINTER = $21000
              $00E0,$0002
     DC.W
              $00E2,$1000
                                BIT PLANE 2 POINTER = $23000
     DC.W
              $00E4,$0002
     DC.W
              $00E6,$3000
                                ;BIT PLANE 3 POINTER = $25000
     DC.W
              $00E8,$0002
     DC.W
              $00EA,$5000
                                BIT PLANE 4 POINTER = $27000
     DC.W
              $00EC,$0002
```

DC.W	\$00EE,\$7000	
DC.W	\$00F0,\$0002	BIT PLANE 5 POINTER = \$29000
DC.W	<b>\$00F2,\$9000</b>	
DC.W	<b>\$00F4,\$0002</b>	;BIT PLANE 6 POINTER = \$2B000
DC.W	\$00F6,\$B000	
DC.W	\$FFFF,\$FFFE	;WAIT FOR THE IMPOSSIBLE, I.E., QUIT

## Chapter 4: Sprite Hardware"

Add the following example of a sprite data structure to page 4-9, under "Building the Data Structure".

; The following data is the data structure for the spaceship sprite. It

; will be located at V = 65 and H = 128 on the screen.

SPRITE:

RITE:		
DC.W	\$6D60,\$7200	;VSTART, HSTART, VSTOP
DC.W	\$0990,\$07E0	FIRST PAIR OF DESCRIPTOR WORDS
DC.W	\$13C8,\$0FF0	
DC.W	\$23C4,\$1FF8	
DC.W	\$13C8,\$0FF0	
DC.W	\$0990, <b>\$07</b> E0	
DC.W	\$0000,\$0000	;END OF SPRITE DATA

Add the following example to page 4-14, under section 4.3.1, "Selecting the Sprite DMA Channel and Setting the Pointers".

Initializing sprite data pointers:

; in this example the processor initializes the data pointers for sprite 0. ; Normally, this is done by the copper. The sprite is at address \$20000. ; SPR0PTH EQU \$DFF120 ;SPRITE 0 POINTER HIGH ORDER WORD SPR0PTL EQU \$DFF122 ;LOW ORDER WORD ; MOVE.L #\$20000,SPR0PTH ;WRITE \$20000 TO SPRITE 0 POINTER Add the following example to page 4-14 under Section 4.3, "Displaying a Sprite".

How to display a sprite:

```
This example displays the spaceship sprite at location V=65, H=128.
; The equates are the usual, so they're not repeated here.
; First, we set up a single bit plane.
                              POINT A0 AT CUSTOM CHIPS
    LEA CUSTOM,A0
                                   ;1 BIT PLANE, COLOR IS ON
    MOVE.W #$1200,BPLCON0(A0)
    MOVE.W #$0000,BPL1MOD(A0)
                                   :MODULO = 0
    MOVE.W #$0000,BPLCON1(A0)
                                   ;HORIZONTAL SCROLL VALUE = 0
    MOVE.W #$0024,BPLCON2(A0)
                                   SPRITES HAVE PRIORITY OVER PLAYFIELDS
    MOVE.W #$0038,DDFSTRT(A0)
                                   SET DATA FETCH START
    MOVE.W #$00D0,DDFSTOP(A0)
                                   SET DATA FETCH STOP
    MOVE.W #$2C81,DIWSTRT(A0)
                                   SET DISPLAY WINDOW START
                                   SET DISPLAY WINDOW STOP
    MOVE.W #$F4C1,DIWSTOP(A0)
: Set up color registers
    MOVE.W #$0008,COLOR00(A0)
                                   BACKGROUND COLOR = DARK BLUE
    MOVE.W #$0000,COLOR01(A0)
                                   ;FOREGROUND COLOR = BLACK
    MOVE.W #$0FF0,COLOR17(A0)
                                   ;COLOR 17 = YELLOW
    MOVE.W #$00FF,COLOR18(A0)
                                   :COLOR 18 = CYAN
    MOVE.W #$0F0F,COLOR19(A0)
                                   ; COLOR 19 = MAGENTA
; Move copper list to $20000
                              POINT A1 AT COPPER LIST DESTINATION
    MOVE.L #$20000,A1
    LEA COPPERL, A2
                              POINT A2 AT COPPER LIST SOURCE
                              ;MOVE A LONG WORD
CLOOP: MOVE.L (A2),(A1)+
             #$FFFFFFE,(A2)+ ;CHECK FOR END OF LIST
    CMP.L
    BNE CLOOP
                              LOOP UNTIL ENTIRE LIST IS MOVED
; Move sprite to $25000
                              POINT A1 AT SPRITE DESTINATION
    MOVE.L #$25000,A1
    LEA SPRITE,A2
                              POINT A2 AT SPRITE SOURCE
SPRLOOP:
                              ;MOVE A LONG WORD
     MOVE.L (A2)(A1)+
             #$0000000,(A2)+
                              CHECK FOR END OF SPRITE
     CMP.L
```

#### BNE SPRLOOP

#### ;LOOP UNTIL ENTIRE SPRITE IS MOVED

: Now we write a dummy sprite to \$30000, since all eight sprites are activated ; at the same time and we're only going to use one. The remaining sprites will point to this dummy sprite data. :WRITE IT MOVE.L #\$0000000,\$30000 ; Point copper at copper list MOVE.L #\$20000,CUSTOM+COP1LC ; Fill bitplane with \$FFFFFFF MOVE.L #\$21000,A1 POINT A1 AT BIT PLANE MOVE.W #2000,D0 :2000 LONG WORDS = 8000 BYTESMOVE.L **#**\$FFFFFFF,(A1)+ ;MOVE A LONG WORD OF \$FFFFFFFF FLOOP: ;DECREMENT COUNTER SUBQ.W #1,D0 **BNE FLOOP** LOOP UNTIL BITPLANE IS FULL : Start DMA MOVE.W CUSTOM+COPJMP1,D0 ;FORCE LOAD INTO COPPER ; PROGRAM COUNTER MOVE.W #\$83A0,(CUSTOM+DMACON) ;BITPLANE, COPPER, AND SPRITE DMA BRA ....next stuff to do... ; This is a copper list for one bit plane, and 8 sprites. The bit plane lives ; at \$21000. Sprite 0 lives at \$25000; all others live at \$30000 (the dummy ; sprite). COPPERL: DC.W \$00E0,\$0002 ;BIT PLANE 1 POINTER = \$21000\$00E2,\$1000 DC.W DC.W \$0120,\$0002 :SPRITE 0 POINTER = \$25000 \$0122,\$5000 DC.W SPRITE 1 POINTER = \$30000DC.W \$0124,\$0003 DC.W \$0126,\$0000 SPRITE 2 POINTER = \$30000DC.W \$0128,\$0003 DC.W \$012A,\$0000 DC.W SPRITE 3 POINTER = \$30000\$012C,\$0003

SPRITE 4 POINTER = \$30000

DC.W

DC.W

\$012E,\$0000

\$0130,\$0003

DC.W	<b>\$0132,\$0000</b>	
DC.W	\$0134,\$0003	;SPRITE 5 POINTER = $30000$
DC.W	<b>\$0136,<b>\$0000</b></b>	
DC.W	<b>\$0138,<b>\$00</b>03</b>	;SPRITE 6 POINTER = $30000$
DC.W	\$013 <b>A</b> ,\$0000	
DC.W	<b>\$013C,\$0003</b>	;SPRITE 7 POINTER = $30000$
DC.W	\$013E,\$0000	
DC.W	\$FFFF,\$FFFE	END OF COPPER LIST
, · Sprite data fo	r spaceship sprite.	It appears on the screen at $V=65$ and
; H=128.	I Spaceship Spires	
, <b>.</b>		
; H=128.	\$6D60,\$7200	;VSTART, HSTART, VSTOP
; H=128. ; SPRITE:		
; H=128. ; SPRITE: DC.W	\$6D60,\$7200	;VSTART, HSTART, VSTOP
; H=128. ; SPRITE: DC.W DC.W	\$6D60,\$7200 \$0990,\$07E0	;VSTART, HSTART, VSTOP
; H=128. ; SPRITE: DC.W DC.W DC.W	\$6D60,\$7200 \$0990,\$07E0 \$13C8,\$0FF0	;VSTART, HSTART, VSTOP
; H=128. ; SPRITE: DC.W DC.W DC.W DC.W	\$6D60,\$7200 \$0990,\$07E0 \$13C8,\$0FF0 \$23C4,\$1FF8	;VSTART, HSTART, VSTOP

Add the following example to page 4-15 under Section 4.4, "Moving a Sprite".

; This is an example of moving a sprite. We bounce the spaceship around : on the screen, making it change direction whenever it reaches an edge. : The sprite position data, containing VSTART and HSTART, lives in memory at \$25000. VSTOP is located at \$25002. We write to these locations to move the sprite. ; Once each frame, we increment (or decrement) VSTART by 1 and HSTART by 2. ; Then we calculate VSTOP which will be the new VSTART + 6. MOVE.B #151,D0 INITIALIZE HORIZONTAL COUNT MOVE.B #194,D1 ;INITIALIZE VERTICAL COUNT MOVE.B #64,D2 INITIALIZE HORIZONTAL POSITION MOVE.B #44,D3 INITIALIZE VERTICAL POSITION MOVE.B #1.D4 INITIALIZE HORIZONTAL INCREMENT VALUE: MOVE.B #1,D5 ;INITIALIZE VERTICAL INCREMENT VALUE ; Here we wait for the vertical blanking bit in INTREQR to turn on. This ; ensures a glitch free display. VLOOP: MOVE.W CUSTOM+INTREQR,D6 ;READ INTERRUPT REQUEST WORD AND.W #\$0020,D6 MASK OFF ALL BUT VERTICAL BLANK BIT BEQ VLOOP ;LOOP UNTIL BIT IS A 1 MOVE.W **#\$0020**,CUSTOM+INTREQ VERTICAL BIT IS ON, SO RESET IT. ; INCREMENT HORIZONTAL VALUE ADD.B D4.D2 #,D0 DECREMENT HORIZONTAL COUNTER SUBQ.B BNE L1 MOVE.B #151,D0 COUNT EXHAUSTED, RESET TO 151 EOR.B #\$FE,D4 NEGATE THE INCREMENT VALUE L1: MOVE.B D2,\$25001 WRITE NEW HSTART VALUE TO SPRITE ADD.B D5.D3 INCREMENT VERTICAL VALUE ;DECREMENT VERTICAL COUNTER SUBQ.B #1,D1 BNE L2 COUNT EXHAUSTED, RESET TO 194 MOVE.B #194.D1 EOR.B #\$FE,D5 ;NEGATE THE INCREMENT VALUE L2: MOVE.B D3,\$25000 WRITE NEW VSTART VALUE TO SPRITE MOVE.B D3,D6 ;MUST NOW CALCULATE NEW VSTOP ADD.B VSTOP ALWAYS VSTART+6 FOR SPACESHIP #6,D6 MOVE.B D6,\$25002 WRITE NEW VSTOP TO SPRITE **BRA VLOOP** :LOOP FOREVER.

Add the following example to page 4-18 under Section 4.6, "Reusing Sprite DMA Channels".

; This example displays the spaceship sprite and then redisplays it as a ; different object. I'm not going to tell you what it is, that's a secret. ; Only the sprite data list is affected by this, so that's all we'll show ; here. However, it looks best with the color registers set as follows: LEA CUSTOM.A0 MOVE.W **#**\$0F00,COLOR17(A0) ; COLOR 17 = REDMOVE.W **#**\$0FF0,COLOR18(A0) ; COLOR 18 = YELLOWMOVE.W **#**\$0FFF,COLOR19(A0) ; COLOR 19 = WHITE; And now for the sprite data. SPRITE: DC.W \$6D60,\$7200 DC.W \$0990,\$07E0 DC.W \$13C8,\$0FF0 DC.W \$23C4,\$1FF8 DC.W \$13C8,\$0FF0 DC.W \$0990,\$07E0 DC.W \$8080,\$8D00 VSTART, HSTART, VSTOP FOR NEW SPRITE DC.W \$1818,\$0000 DC.W \$7E7E,\$0000 DC.W \$7FFE,\$0000 DC.W **\$FFFF,\$2000** DC.W **\$FFFF,\$2000** DC.W **\$FFFF,\$3000** DC.W **\$FFFF,\$3000** DC.W \$7FFE,\$1800 DC.W \$7FFE,\$0C00 DC.W \$3FFC,\$0000 DC.W \$0FF0,\$0000 DC.W \$03C0,\$0000 \$0180,\$0000 DC.W DC.W \$0000,\$0000 ;END OF SPRITE DATA

Add the following example to page 4-23 under section 4.8, "Attached Sprites".

; The following data structure is for the six color spaceship made with two

; attached sprites.

SPRITE0:

RIIEO.		
DC.W	\$6D60,\$7200	; VSTART = 65, HSTART = $128$
DC.W	<b>\$0C30,\$0000</b>	FIRST COLOR DESCRIPTOR WORD
DC.W	\$1818,\$0420	
DC.W	\$342C,\$0E70	
DC.W	\$1818,\$0420	
$\mathbf{DC}.\mathbf{W}$	<b>\$0C30,\$0000</b>	
DC.W	\$0000,\$0000	;END OF SPRITE 0
RITE1:		
DC.W	\$6D60,\$7280	;SAVE AS SPRITE 0 EXCEPT ATTACH BIT ON.
DC.W	\$07E0,\$0000	FIRST DESCRIPTOR WORD FOR SPRITE 1
DC.W	<b>\$0FF0,\$0000</b>	

,	
SPRITE1:	
DC \	λ

DC.W DC.W

DC.W

DC.W

\$1FF8,\$0000

**\$0FF0,\$0000** \$07E0,\$0000

\$0000,\$0000

;END OF SPRITE 1

## Addendum to Chapter 5: Audio Hardware

## Audio State Machine

There is one audio state machine for each channel. See the attached state diagram. The machine has eight states and is clocked at the system clock frequency of 3.58 MHz. Three of the states are basically unused and just transfer back to the idle (000) state. One of the paths out of the idle state is designed for interrupt-driven operation (processor provides data), and the other path is designed for DMA-driven operation (Agnus provides data).

In interrupt-driven operation, transfer to the main loop (010,011 states) is immediate upon data written by the processor. In the 010 state the upper byte is output, and in the 011 state the lower byte is output. Transitions  $010\rightarrow011\rightarrow010$ ... occur whenever the period counter counts down to one. The period counter is reloaded at these transitions. As long as the interrupt is cleared by the processor in time, the machine remains in the main loop. Otherwise, it enters the idle state. Interrupts are generated on every word (011 $\rightarrow$ 010) transition.

In DMA-driven operation, transit to the 001 state occurs and DMA requests are sent to Agnus as soon as DMA is turned on. Because of pipelining in Agnus, the first data word must be thrown away. State 101 is entered as soon as this word arrives. A request for the next data word has already gone out. When the data arrives, state 010 is entered and the main loop continues until the DMA is turned off. The length counter counts down once with each word that comes in. When it finishes, a DMA restart request goes to Agnus along with the regular DMA request. This tells Agnus to reset the pointer to the beginning of the table of data. Also, the length counter is reloaded and an interrupt request goes out soon after the length counter finishes (counts to one). The request goes out just as the last word of the waveform starts its output.

DMA requests and restart requests are transferred to Agnus once each horizontal line, and the data comes back about 14 clock cycles (of 280ns) later.

In attach mode, things run a little differently. Attach volume has requests the same as normal operation (on the  $011\rightarrow010$ ) transition). In attach period, a set of requests occurs on the  $010\rightarrow011$  transition. With both attach period and attach volume high, requests occur on both transitions.

One of the consequences of the way this machine works is that if the sampling rate is set much higher than the normal max sampling rate (~29 kHz), the two samples in the buffer register will be repeated. If the filter on the Amiga is bypassed and the volume is set to max (40 hex), this feature can be used to make modulated carriers up to 1.79 MHz. The modulation is placed in the memory map, with plus values in the even bytes, and minus values in the odd bytes.

The following list shows the symbols used in the state diagram. Capitals are external signals; small letters are local signals.

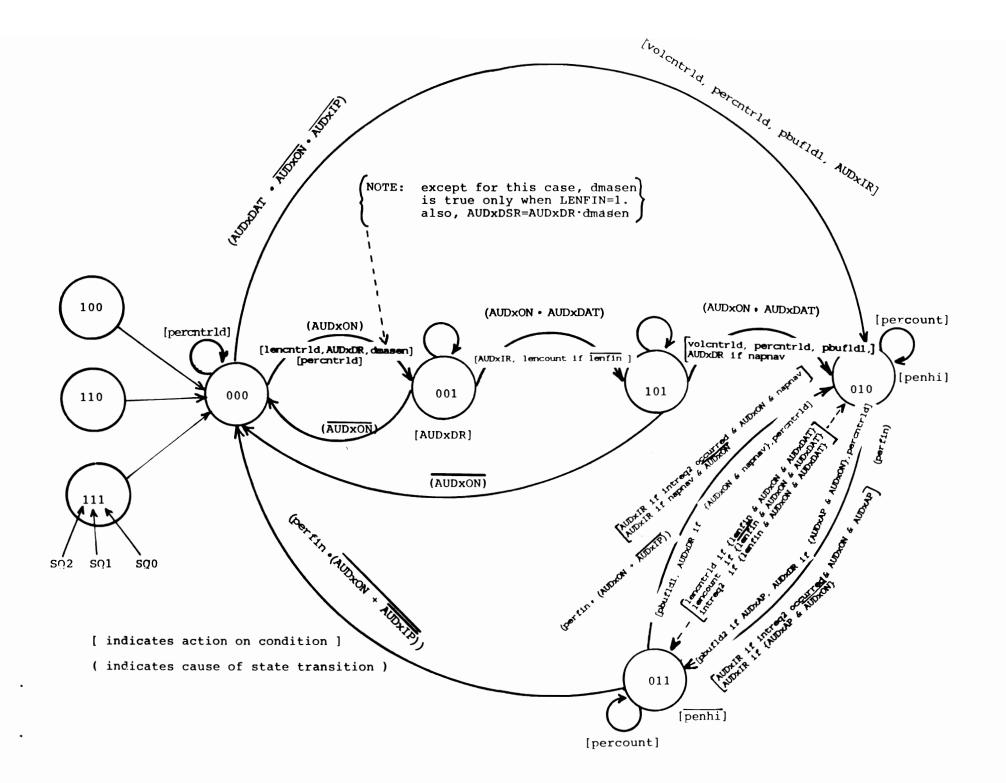
AUDxÓN	DMA on "x" indicates channel number (signal from DMACON).
AUDxIP	Audio interrupt pending (input to channel from interrupt circuitry)
AUDxIR	Audio interrupt request (output from channel to interrupt circuitry)
intreq1	Interrupt request that combines with intreq2 to form AUDxIR
intreq2	Prepare for interrupt request. Request comes out after the next $011 \rightarrow 010$ transition in normal operation.
AUDxDAT	Audio data load signal. Loads 16 bits of data to audio channel.
AUDxDR	Audio DMA request to Agnus for one word of data.
AUDxDSR	Audio DMA request to Agnus to reset pointer to start of block
dmasen	Restart request enable.
percntrld	Reload period counter from backup latch typically written by processor with AUDxPER (can also be written by attach mode).
percount	Count period counter down one latch.
perfin	Period counter finished (value $= 1$ ).
lencntrld	Reload length counter from backup latch.
lencount	Count length counter down one notch.
lenfin	Length counter finished (value $= 1$ ).
volcntrld	Reload volume counter from backup latch.
pbufld1	Load output buffer from holding latch written to by AUDxDAT.
pbufld $2$	Like pbufld1, but only during $010 \rightarrow 011$ with attach period.
AUDxAV	Attach volume. Send data to volume latch of next channel instead of to $D \rightarrow A$ converter.

.

AUDxAPAttach period. Send data to period latch of next channel instead of to the<br/>D $\rightarrow$ A converter.penhiEnable the high 8 bits of data to go to the D $\rightarrow$ A converter.napnav/AUDxAV \* /AUDxAP + AUDxAV—no attach stuff or else attach volume.

Condition for normal DMA and interrupt requests.

sq2,1,0 The name of the state flip-flops, MSB to LSB.



## Addendum to Chapter 6: Blitter Hardware

## Blitter Operations and System DMA

This section explains how the operation of the blitter affects the performance of the rest of the system. The section covers the following topics:

- o Blitter direct memory access (DMA) priority
- o DMA time slot allocation
- o Bus sharing between the 68000 and the bit-plane operations of the blitter and Copper.
- o Effects of different playfield display sizes on sprite display
- o Effects of blitter operation on the 68000's access to memory

## 1. Blitter DMA Priority

The blitter performs its various data fetch, modify, and store operations through a DMA sequence. It shares memory access with other devices in the system. In a system like the Amiga that uses a lot of DMA there must be some control over the priority of the accessing devices. A device's priority indicates its importance relative to other devices. The list below shows the order of priority for DMA operations, from highest priority to lowest.

- o Disk DMA-handles communications with the disk
- o Audio DMA-produces the sound
- o Bit-plane DMA produces the static display
- o Sprite DMA-produces the dynamic display
- o Copper-display-synchronized coprocessor
- o Blitter-data copying and line drawing device

o 68000 microprocessor — central processor

The first four devices in the list all have the same priority. If a disk DMA cycle is missed, some disk data is lost. If an audio cycle is missed, noise is added to the audio output. In displays (bitplane or sprite), there may be flashes or other interruptions on-screen. None of these situations are desirable and are avoided by making all of these devices priority 1. Under certain circumstances, however, the bit-plane DMA will take priority over sprite DMA.

Each of the first four devices in the list is specifically allocated a group of time slots during each horizontal scan of the video beam. If the device does not specifically request to use one of its allocated time slots, the slot is open for other uses.

The Copper has the next priority because it is designed as a display-synchronized coprocessor. It has to perform its operations at the same time during each display frame to remain synchronized with the display beam sweeping across the screen.

At the bottom of the list come the blitter and the 68000, in that order. The blitter is optimized for data copying, modifying, and line drawing operations. It performs these kinds of operations much faster than the 68000 could perform them. It is, therefore, given a higher priority than the processor so that it can do its job in the most efficient manner.

## 2. DMA Time Slot Allocation

During a horizontal scan line (about 63 microseconds), there are 227.5 "color-clocks". A color clock is the basic timing interval for memory access within this system and amounts to approximately 280 ns.

This interval of 227.5 color clocks includes both display time and non-display time on a horizontal line. Of this total time, there are 226 possible memory access cycles at 280 ns each to be allocated to the various devices that need memory access.

Each of the first four time slots in the list below is assigned at one of the odd-numbered slots available. This assignment was made to allow maximum bus utilization by the 68000 (as shown in the section titled "Bit-Plane/Processor Bus Sharing"). The even-numbered slots can be used by the Copper, blitter, or 68000.

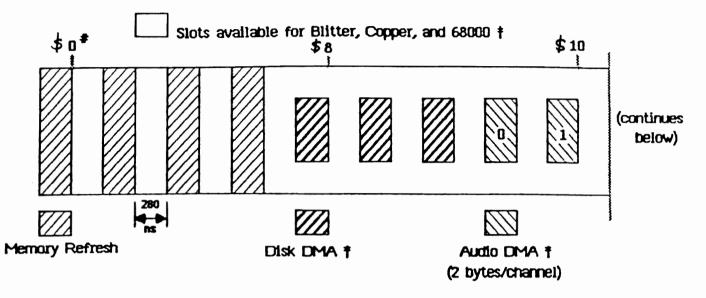
Here is the time-slot allocation per horizontal line:

o 4 cycles for memory refresh

- o 3 cycles for disk DMA
- o 4 cycles for audio DMA (2 bytes per channel)
- o 16 cycles for sprite DMA (2 words per channel)
- o 80 cycles for bit-plane DMA (even and/or odd according to the display size used)

The figures on the following pages show one complete horizontal scan line and how the clock cycles are allocated.

# DMA Time Slot Allocation / Horizontal line



† These operations only take slots if the associated operation is being performed

Note: Copper Data Move instructions require 4 slots. Copper Wait instructions require 6 slots.

# This cycle 0 appears to exclude one of the memory refresh cycles. This is not the case.

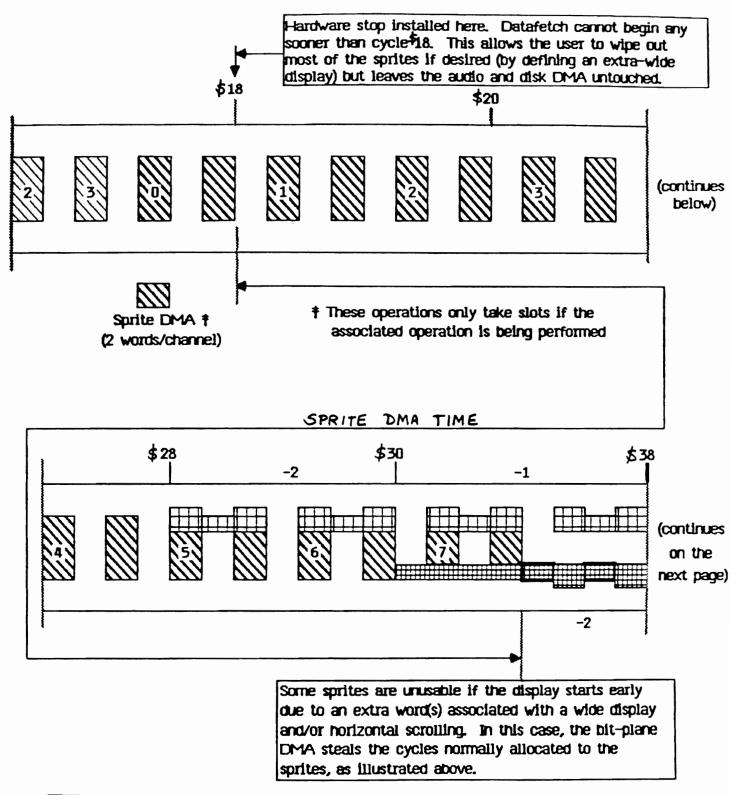
Actual system hardware demands certain specific values for data fetch start and display start. Therefore this timing chart has been "adjusted" to match those requirements.

\$ Indicates a hex number.

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# DMA Time Slot Allocation / Horizontal line (cont'd)

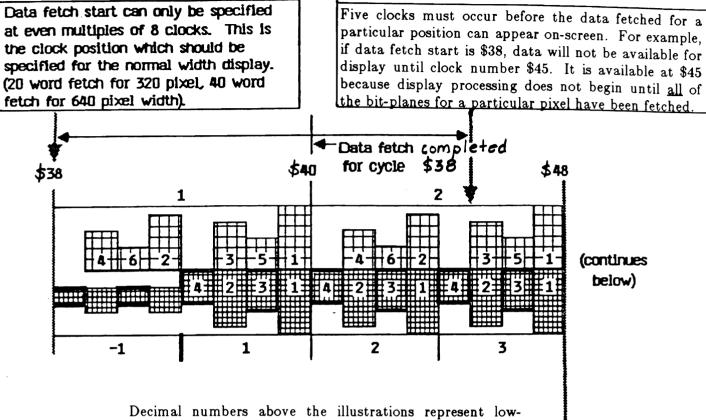




320 mode Bit-Plane DMA, by plane †

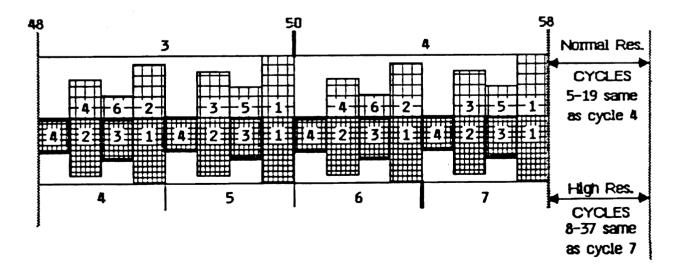
640 mode Bit-Plane DMA, by plane †

# DMA Time Slot Allocation / Horizontal line (cont'd)



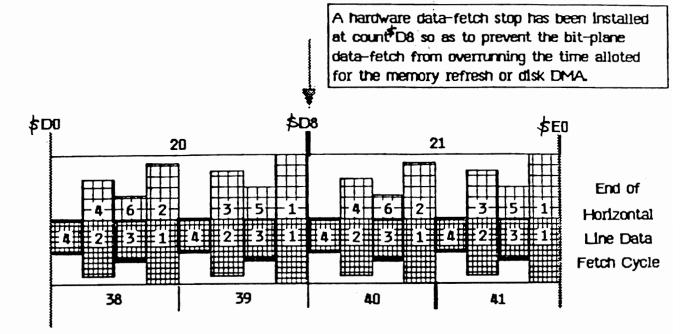
Decimal numbers above the illustrations represent lowresolution cycles. Decimal numbers below the illustrations represent high-resolution cycles. Negative numbers indicate the start of data fetch for displays that are larger than normal.

Decimal numbers inside the illustrations represent the bitplane for which the data is being fetched.





# DMA Time Slot Allocation / Horizontal line (cont'd)



## 3. Bit-Plane/Processor Bus Sharing

The memory access cycles are interleaved to allow the 68000 processor to operate at close to its maximum speed. The 68000 spends about half of a complete processor instruction time doing internal operations and the other half accessing memory. Therefore, if the 68000 is given each alternate 280 ns memory cycle, it will appear to the 68000 that it has the memory all of the time and it will run at full speed.

Not all of the 68000 instructions allow this even-cycle allocation to mate perfectly all of the time. If it doesn't, the processor will have to wait until its next memory slot is available before continuing. Most 68000 instructions do not cause cycles to be missed, so it will run at full speed most of the time if there is no blitter DMA interference.

The following figure illustrates the normal cycle of the 68000.

<> average 68000 cycle>										
<> operation portion	<> access portion									
odd cycle number, assigned to other devices	even cycle, possibly used by 68000 if it requested it									

If there are 4 or less low-resolution bit-planes, then the 68000 can be granted each alternate memory cycle (if it is ready to ask for the cycle and is the highest priority item at the time). However, if there are more than 4 bit-planes, the bit-plane DMA will begin to steal cycles from the 68000 during the display time.

The following figure illustrates the time slots (each 280 ns) that will be taken by bit-plane DMA during the display time (160 slots out of 227 for each horizontal line for a 320-pixel, low-resolution display) for a 6-bit-plane display. As you will see from the figure, the bit-plane steals 50% of the open slots that the processor might have used if there were only 4 bit-planes displayed.

Т		-	T+7						
+		*		+		*			
	4	6	2		3	5	1		

- + an open memory slot that the 68000 might use
- \* a slot that cannot be used by the 68000 due to added bit-plane DMA

If you specify 4 high-resolution bit-planes (640 pixels wide), bit-plane DMA needs all of the available memory time slots during the display time just to grab the 40 data words for each line of the 4 bit-planes (40 X 4 = 160 time slots). This effectively locks out the 68000 (as well as the blitter or Copper) from any memory access during the actual period of the display.

The following figure shows how the time slots are allocated for high-resolution bit-planes.

Т		- timing cycle -							
4	2	3	1	4	2	3	1		

## 4. Effects of Different Display Sizes

The normal, full-sized screen consists of 320 pixels for low-resolution mode or 640 pixels for highresolution mode. This means that either 20 or 40 words will be fetched during the horizontal line display time.

When you want to scroll either or both playfields, one extra data word per line must be fetched from the memory. This extra word provides the pixels that will show on-screen as scrolling occurs.

Screen size is adjustable (see the descriptions of DIWSTRT, DIWSTOP, and DDFSTRT in Chapter 3, "Playfield Hardware". Bit-plane DMA takes precedence over sprite DMA. Therefore, larger screens may block out one or more of the highest-numbered sprites, especially with scrolling.

## 5. Effects of Blitter Operation

As mentioned above, the blitter normally has a higher priority than the processor for DMA cycles. There are certain cases, however, where the blitter and the 68000 will share memory cycles. The blitter uses every cycle. If given the chance, it would steal every available memory cycle. Display, disk, and audio DMA take precedence, so the blitter cannot block them from bus access. Depending on the state of the setting of the blitter DMA mode bit, commonly referred to as the "blitter-nasty" bit, the processor may not be so lucky. This bit is called BLTPRI (for "blitter has priority over processor") and is in register DMACONW.

If BLTPRI is a 1, the blitter will hang onto the bus for every available memory cycle. In the nondisplay time (horizontal blanking interval), this could potentially be every single cycle. Even in blitter-nasty mode, there may be cases when you can grant a few cycles to the 68000. Even though the blitter would seem to be acting as a bus-hog, there are data windows that occur during certain blitter operations when the blitter is doing something internally rather than on the bus. These cycles are part of the basic operating sequence of the blitter itself.

If BLTPRI is a 0, the DMA manager will monitor the 68000 cycle requests. If the 68000 is unsatisfied for three consecutive memory cycles, the blitter will release the bus for 1 cycle.

The DMA manager knows when these cycles of the blitter are about to occur and will release them to the 68000 if it is waiting for memory access. Blitter-nasty mode doesn't *prevent* the 68000 from executing instructions, but it surely slows it down.

Table 1 shows all of the possible operating modes of the blitter along with the distribution of the memory access windows within its operation. The table shows three words of a blit (the first, any middle, and the last) and how the bus activity occurs for this sequence. The following conventions are used:

- o A, B, and C stand for the sources.
- o D stands for the destination.
- o A0 is the first memory word fetch; A1 is any middle memory word fetch; A2 is the last memory word fetch.
- An asterisk (\*) following a source name indicates an internal cycle used for that source; the cycle does no data fetch. For these functions, the cycle fetch columns have blank spaces that indicate open memory cycles where some other device can access the bus. Thus functions such as item 3, ABC\*->D can be read as AB->D.

- o Z stands for a no-data-stored case (no destination).
- o (f) stands for fill and (n) stands for no-fill, if the use of fill makes a difference.

Table 1: Blitter Memory Bus Activity

(	Operation Type		N	<i>A</i> emo	ry Cy	ycle [	)ata ]	Fetch	or D	ata S	store	Usag	e	
1.	АВСD	<b>A0</b>	<b>B0</b>	C0		<b>A</b> 1	<b>B</b> 1	C1	D0	A2	B2	C2	D1	D2
<b>2</b> .	АВСZ	<b>A0</b>	<b>B0</b>	<b>C0</b>	<b>A1</b>	<b>B1</b>	C1	A2	<b>B2</b>	C2				
3.	A B C*D(f)	<b>A0</b>	<b>B0</b>			<b>A1</b>	<b>B</b> 1	D0		A2	B2	D1		D2
<b>4</b> .	A B C*D(n)	<b>A0</b>	<b>B0</b>	<b>A</b> 1	<b>B</b> 1	D0		A2	<b>B</b> 2	D1		D2		
<b>5</b> .	A B C*Z	<b>A0</b>	<b>B0</b>		<b>A</b> 1	<b>B1</b>		A2	B2					
<b>6</b> .	A CD	<b>A</b> 0	<b>C0</b>		<b>A1</b>	C1	D0	A2	C2	D1		D2		
7.	A CZ	<b>A0</b>	<b>C0</b>	<b>A1</b>	C1	A2	C2							
8.	A C*D(f)	<b>A0</b>			<b>A1</b>	D0	A2	D1		D2				
9.	A C*D(n)	<b>A0</b>		A1	D0	A2	D1		D2					
10.	A C*Z	<b>A</b> 0		<b>A1</b>		A2								
11.	A* B C D		BO	<b>C0</b>			B1	<b>C1</b>	D0		B2	C2	D1	D2
1 <b>2</b> .	A* B CZ		<b>B</b> 0	<b>C0</b>		B1	C1		<b>B2</b>	C2				
13.	A* B C* D(f)		<b>B0</b>				<b>B</b> 1	D0			B2	D1		D2
14.	A* B C*D(n)		<b>B0</b>			<b>B</b> 1	D0		B2	D1		D2		
15.	A* B C*Z		<b>B</b> 0			<b>B</b> 1			<b>B2</b>					
16.	A* CD		$\mathbf{C0}$			C1	D0		C2	D1		D2		
17.	A* CZ		C0		C1		C2							
18.	A* C*D(f)					D0			D1			D2		
19.	A* C*D(n)				D0		D1		D2					
<b>20</b> .	A* C*Z													
21.	LINE DRAWING													
	MODE		C0	*	D0		C1	*	D1		C2	*	D2	

Notes for Table 1:

- a. Function 9 is a block move only when both shift values are 0. Use function 14 if either shift value is nonzero.
- b. Function 10 is a zero test; function 19 is a constant fill.
- c. Function 21 is for line drawing. The cycles marked with \* are not used by the blitter but are not released to the system. Sorry.

d. Functions 11, 13, 14, and 16 have special restraints when source and data overlap. During some priority overrides, the order of data fetch and store will be reversed. When using these functions with overlapping source and data, the overlap must be greater than two words. Otherwise, you risk overwriting your source data.

### APPENDIX F - Amiga Hardware Manual

This appendix contains information about the 8520 peripheral interface adapters.

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### QUICK REFERENCE

### BRIEF ADDRESS MAP FOR 8520'S

The system hardware selects the 8520's (aka CIA's) when the upper three address bits are 101.

Furthermore, CIAA is selected when A12 is low, A13 high; CIAB is selected when A12 is high, A13 low.

You can use either byte or word addresses to access the 8520's. For byte access (seems to be the usual case), A0 must be 0 for CIAA, 1 for CIAB. For word access, CIAB communicates on data bits 15"8; CIAA communicates on data bits 7"0. (A0 is always 0 for word access, naturally).

Address bits All, Al0, A9, and A8 are used to specify which of the 16 internal registers you want to access. This is indicated by "r" in the address.

All other bits are don't cares. So, CIAA is selected by the following binary address: 101x xxxx xx01 rrrr xxxx xxx0. CIAB address: 101x xxxx xx10 rrrr xxxx xxx1

With future expansion in mind, we have decided on the following addresses: CIAA = BFEr01; CIAB = BFDr00.

### CIAB Address Map

	Register Name	7	6	5	Data b 4		2	1	0
BFD000 BFD100 BFD200 BFD300 BFD400 BFD500 BFD600 BFD700 BFD800 BFD900 BFDA00 BFDB00 BFDD00 BFDD00 BFDD00 BFDE00		/DTR /MTR ddr fo ddr fo CIAB CIAB CIAB CIAB Horizo Horizo not us CIAB CIAB	/RTS /SEL3 or port or port fimer 1 fimer 1 fim	/CD /SEL2 t A (B t B (B A low A high B low B high sync e sync e sync e data upt co	/CTS /SEL1 FD000); FD100); byte byte byte	/DSR /SEL0 1 = 0 1 = 0 ounter ounter	SEL /SIDE output output bits 7 bits 1 bits 2	POUT DIR (set t (set t .5~8	BUSY
BEDE00		CIAB	Contro	l regi	ster B				

Note: CIAB can generate INT6.

CIAA Address Map

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Byte	Register			Data	bits				
Address	Name	7	6 5 	4	3	2	1	0	
BFE001		IR1 /F arallel		Y /TK0	/WPRO	/CHING	/LED	OVL	
BFE101 BFE201	do	ir for	port A	(BFE001)	; 1 = 0	output	(Set	to 0x03)	
BFE301 BFE401	do Cl	ir for j [AA Tim	port B er <b>A</b> lo	(BFE101) w byte	; 1 = 0	output	(Can	be in or	out)
BFE501			er A hi er B lo	gh byte					
BFE601 BFE701	CI	AA Tim	er B hi	gh byte					
BFE801 BFE901				nter bit nter bit					
BFEA01 BFEB01		) Hz ev ot used		nter bit	<b>s 23~1</b>	5			
BFEC01	CI	AA Ser	ial dat	a regist			)		
BFED01 BFEE01				control gister A		er			
BFEF01	CI	AA Con	trol re	gister E	5				

Note: CIAA can generate INT2.

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INTERFACE SIGNALS

Clock Input

The 02 clock is a TTL compatible input used for internal device operation and as a timing reference for communicating with the system data bus.

CS - Chip Select Input

The CS input controls the activity of the 8520. A low level on CS while 02 is high causes the device to respond to signals on the R/W and address (RS) lines. A high on CS prevents these lines from controlling the 8520. The CS line is normally activated (low) at 02 by the appropriate address combination.

R/W - Read/Write Input

The R/W signal is normally supplied by the microprocessor and controls the direction of data transfers of the 8520. A high on R/W indicates a read (data transfer out of the 8520), while a low indicates a write (data transfer into the 8520).

RS3-RS0 - Address Inputs

The address inputs select the internal registers as described by the Register Map.

DB7-DB0 - Data Bus Inputs/Outputs

The eight data bus output pins transfer information between the 8520 and the system data bus. These pins are high impedance inputs unless CS is low and R/W and 02 are high, to read the device. During this read, the data bus output buffers are enabled, driving the data from the selected register onto the system data bus.

IRQ - Interrupt Request Output

IRQ is an open drain output normally connected to the processor interrupt input. An external pull-up resistor holds the signal high, allowing multiple IRQ outputs to be connected together. The IRQ output is normally off (high impedance) and is activated low as indicated in the functional description.

RES - Reset Input

A low on the RES pin resets all internal registers. The port pins are set as inputs and port registers to zero (although a read of the ports will return all highs because of passive pull-ups). the timer control registers are set to zero and the timer latches to all ones. All other registers are reset to zero.

REGISTER MAP

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Each 8520 has 16 registers which you may read or write. Here is the list of registers and the access address of each within the memory space dedicated to the 8520:

				Register	-	
RS3	RS2	RS1	RS0	# (hex)	NAME	MEANING
0	0	0	0	0	PRA	Peripheral Data Register A
0	0	0	1	1	PRB	Peripheral Data Register B
0	0	1	0	2	DDRA	Data Direction Register A
0	0	1	1	3	DDRB	Direction Register B
0	1	0	0	4	TALO	Timer A Low register
0	1	0	1	5	TAHI	Timer A High register
0	1	1	0	6	TBLO	Timer B Low register
0	1	1	1	7	TBHI	Timer B High register
1	0	0	0	8		Event LSB
1	0	0	1	9		Event 8-15
1	0	1	0	Α		Event MSB
1	0	1	1	В		No Connect

Oct 15 17:15 1985 Appendix_F Page 4											
1100CSDRSerial Data Register1101DICRInterrupt Control Register1110ECRAControl Register A111FCRBControl Register B											
SOFTWARE NOTE:											
The operating system kernel has already allocated the use of all 4 of the timers TA and TB in the 8520's. If you are running under control of the system exec, be aware of the following allocation of system resources:											
<pre>8520A, timer A Commodore serial communications (if no serial comm happening, timer becomes available). 8520A, timer B Video beam follower (used when synchronizing the blitter device to the video beam, see the description of QBSBlit in the system software manual). If no beam-sync'ed blits are in process, this timer will be available.</pre>	1										
8520B, timer A Keyboard (used continuously, wheneve system EXEC is in control). 8520B, timer B Virtual timer device (used continuous whenever system EXEC is in control, used for task switching/interrupts).											
REGISTER NAMES											

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The names of the registers within the 8520's are as follows. The address at which each is to be accessed is given here in this list.

Address	for:		
8520-A	8520-B	NAME	EXPLANATION
			(write)/(read mode)
BFE001 BFE101 BFE201 BFE301 BFE401 BFE501 BFE601 BFE701 BFE801 BFE901	BFD000 BFD100 BFD200 BFD300 BFD400 BFD500 BFD600 BFD700 BFD800 BFD900	PRA PRB DDRB DDRA TALO TAHI TBLO TBHI	Peripheral Data Register A Peripheral Data Register B Data Direction Register "A" Data Direction Register "B" TIMER A Low Register TIMER A High Register TIMER B Low Register TIMER B High Register Event LSB Event 8 - 15

BFEA01	BFDA00		Event MSB
BFEB01	BFDB00		No connect
BFEC01	BEDC00	SDR	Serial Data Register
BFED01	BEDD00	ICR	Interrupt Control Register
BFEE01	BFDE00	CRA	Control Register A
BFEF01	BFDF00	CRB	Control Register B

REGISTER FUNCTIONAL DESCRIPTION:

I/O PORTS (PRA, PRB, DDRA, DDRB)

Ports A and B each consist of an 8-bit Peripheral Data Register (PR) and an 8-bit data direction register (DDR). If a bit in the DDR is set to a 1, the corresponding bit position in the PR becomes an output. If a DDR bit is set to a 0, the corresponding PR bit is defined as an input.

When you READ a PR register, you read the actual current state of the I/O pins (PAO-PA7, PBO-PB7, regardless of whether you have set them to be inputs or outputs.

Ports A and B have passive pull-up devices as well as active pull-ups, providing both CMOS and TTL compatibility. Both ports have two TTL load drive capability.

In addition to their normal I/O operations, ports PB6 and PB7 also provide timer output functions.

### HANDSHAKING

Handshaking occurs on data transfers using the PC output pin and the FLAG input pin. PC will go low on the third cycle after a Port B access. This signal can be used to indicate "data ready" at PORT B or "data accepted" from PORT B. Handshaking on 16-bit data transfers (using both ports A and B) is possible by always reading or writing PORT A first. FLAG is a negative edge sensitive input which can be used for receiving the PC output from another 8520, or as a general purpose interrupt input. Any negative transition on FLAG will set the FLAG interrupt bit.

REG	NAME	D <b>7</b>	D6	D5	D <b>4</b>	D <b>3</b>	D <b>2</b>	D1	D0
0	PRA	<b>PA7</b>	PA6	P <b>A5</b>	PA4	PA3	PA2	PA1	P <b>A</b> 0
1	PRB	PB7	PB6	PB5	PB4	PB3	PB2	PB1	PB0
2	DDRA	DPA7	DPA6	DPA5	DPA4	DPA3	DPA2	DPA1	DPA0
3	DDRB	DPB7	DPB6	DPB5	DPB4	DPB3	DPB2	DPB1	DPB0

INTERVAL TIMERS (TIMER A, TIMER B)

Each interval timer consists of a 16-bit read-only Timer Counter and a 16-bit write-only Timer Latch. Data written to the timer is latched into the Timer Latch, while data read from the timer

is the present contents of the Timer Counter.

The latch is also called a prescalar in that it represents the countdown value which must be counted before the timer reaches an underflow (no more counts) condition. This latch (prescalar) value is a divider of the input clocking frequency.

The timers can be used independently, or linked for extended operations. Various timer operating modes allow generation of long time delays, variable width pulses, pulse trains, and variable frequency waveforms. Utilizing the CNT input, the timers can count external pulses or measure frequency, pulse width, and delay times of external signals.

Each timer has an associated control register, providing independent control over each of the following functions:

### START/STOP

A control bit allows the timer to be started or stopped by the microprocessor at any time.

### PB On/Off

A control bit allows the timer output to appear on a PORT B output line (PB6 for timer A and PB7 for timer B). This function over-rides the DDRB control bit and forces the appropriate PB line to become an output.

### Toggle/Pulse

A control bit selects the output applied to PORT B while the PB On/Off bit is ON. On every timer underflow, the output can either toggle or generate a single positive pulse of one cycle duration. The Toggle output is set high whenever the timer is started, and set low by RES.

One-Shot/Continuous

A control bit selects either timer mode. In one-shot mode, the timer will count down from the latched value to zero, generate an interrupt, reload the latched value, then stop. In continuous mode, the timer will count down from the latched value to zero, generate an interrupt, reload the latched value, and repeat the procedure continuously.

In one-shot mode, a write to Timer High (register 5 for Timer A, register 7 for Timer B) will transfer the timer latch to the counter and initiate counting regardless of the start bit.

Force Load

A strobe bit allows the timer latch to be loaded into the timer counter at any time, whether the timer is running or not.

### INPUT MODES

Control bits allow selection of the clock used to decrement the timer. TIMER A can count 02 clock pulses or external pulses applied to the CNT pin. TIMER B can count 02 pulses, external CNT pulses, TIMER A underflow pulses, or TIMER A underflow pulses while the CNT pin is held high.

The timer latch is loaded into the timer on any timer underflow, on a force load, or following a write to the high byte of the prescalar while the timer is stopped. If the timer is running, a write to the high byte will load the timer latch, but not reload the counter.

### BIT NAMES on READ-register

REG	NAME	D <b>7</b>	D6	D5	D <b>4</b>	D3	D2	D1	D0
4	TALO	TAL7	TAL6	TAL5	TAL4	TAL 3	TAL 2	TAL1	TAL0
5	TAHI			TAH5					
б	TBLO			TBL5					
7	TBHI	TBH7	TBH6	TBH5	TBH4	ТВНЗ	TBH2	TBH1	<b>TBHO</b>

### BIT NAMES on WRITE-register

REG	NAME	D <b>7</b>	D6	D5	D <b>4</b>	D3	D <b>2</b>	D1	D0
4	TALO	PAL7	PAL6	PAL5	PAL4	PAL3	PAL2	PAL1	PAL0
5	TAHI	PAH7	PAH6	PAH5	PAH4	РАНЗ	PAH2	PAH1	PAH0
6	TBLO	PBL7	PBL6	PBL5	PBL4	PBL3	PBL2	PBL1	PBLO
7	TBHI	PBH7	PBH6	PBH5	PBH4	PBH3	PBH2	PBH1	PBH0

### TIME OF DAY CLOCK

TOD consists of a 24-bit binary counter. Positive edge transitions on this pin cause the binary counter to increment. The TOD pin has a passive pull-up on it.

A programmable ALARM is provided for generating an interrupt at a desired time. The ALARM registers are located at the same addresses as the corresponding TOD registers. Access to the ALARM is governed by a Control Register bit. The ALARM is write-only; any read of a TOD address will read time regardless of the state of the ALARM access bit.

A specific sequence of events must be followed for proper setting and reading of TOD. TOD is automatically stopped whenever a write to the register occurs. The clock will not start again until after a write to the LSB Event Register. This assures that TOD will always start at the desired time.

Since a carry from one stage to the next can occur at any time with respect to a read operation, a latching function is included to keep all Time Of Day information constant during a read sequence. All TOD registers latch on a read of MSB Event and remain latched until after a read of LSB Event. The TOD clock continues to count when the output registers are latched. If only one register is to be read, there is no carry problem and the register can be read "on the fly" provided that any read of MSB Event is followed by a read of LSB Event to disable the latching.

> BIT NAMES for WRITE TIME/ALARM or READ TIME REG NAME - - -----LSB Event E7 E6 E5 E4 E3 8 E2 **E1** E0 9 Event 8-15 E15 E14 E13 E12 E11 E10 E9 **E8** MSB Event E23 E22 E21 E20 E19 E18 E17 E16 Α WRITE

CRB7 = 0CRB7 = 1 ALARM

SERIAL PORT (SDR)

The serial port is a buffered, 8-bit synchronous shift register. A control bit selects input or output mode.

INPUT MODE

In input mode, data on the SP pin is shifted into the shift register on the rising edge of the signal applied to the CNT pin. After 8 CNT pulses, the data in the shift register is dumped into the Serial Data Register and an interrupt is generated.

OUTPUT MODE

In the output mode, TIMER A is used as the baud rate generator. Data is shifted out on the SP pin at 1/2 the underflow rate of TIMER A. The maximum baud rate possible is 02 divided by 4, but the maximum usable baud rate will be determined by line loading and the speed at

which the receiver responds to input data.

To begin transmission, you must first set up TIMER A in continuous mode, and start the timer. Transmission will start following a write to the Serial Data Register. The clock signal derived from TIMER A appears as an output on the CNT pin. The data in the Serial Data Register will be loaded into the shift register, then shifted out to the SP pin when a CNT pulse occurs. Data shifted out becomes valid on the next falling edge of CNT and remains valid until the next falling edge.

After 8 CNT pulses, an interrupt is generated to indicate that more data can be sent. If the Serial Data Register was reloaded with new information prior to this interrupt, the new data will automatically be loaded into the shift register and transmission will continue.

If no further data is to be transmitted after the 8th CNT pulse, CNT will return high and SP will remain at the level of the last data bit transmitted.

SDR data is shifted out MSB first. Serial input data should appear in this same format.

### BIDIRECTIONAL FEATURE

The bidirectional capability of the Serial Port and CNT clock allows many 8520's to be connected to a common serial communications bus on which one 8520 acts as a master, sourcing data and shift clock, while all other 8520 chips act as slaves. Both CNT and SP outputs are open drain to allow such a common bus. Protocol for master/slave selection can be transmitted over the serial bus or via dedicated handshake lines.

REG	NAME	D <b>7</b>	D6	D5	D4	D3	D <b>2</b>	D1	D0	
С	SDR	S7	S6	S5	<b>S4</b>	<b>S</b> 3	S2	S1	S0	

INTERRUPT CONTROL REGISTER (ICR)

There are 5 sources of interrupts on the 8520:

-underflow from TIMER A (timer counts down past 0) -underflow from TIMER B -TOD ALARM -Serial Port Full/Empty -FLAG

A single register provides masking and interrupt information. The Interrupt Control Register consists of a write-only MASK register and a read-only DATA register. Any interrupt will set the corresponding bit in the DATA register. Any interrupt that is enabled by a 1-bit in that position in the MASK will set the IR bit (MSB) of the DATA register,

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and bring the IRQ pin low. In a multi-chip system, the IR bit can be polled to detect which chip has generated an interrupt request.

When you read the DATA register, its contents are cleared (set to 0) and the IRQ line returns to a high state. Since it is cleared on a read, you must assure that your interrupt polling or interrupt service code can preserve and respond to all bits which may have been set in the DATA register at the time it was read. With proper preservation and response, it is easily possible to intermix polled and direct interrupt service methods.

You can set or clear one or more bits of the MASK register without affecting the current state of any of the other bits in the register. This is done by setting the appropriate state of the MSBit, which is called the SET/CLEAR bit. In bits 6-0, you yourself form a mask which specifies which of the bits you wish to affect. Then, using bit 7, you specify HOW the bits in corresponding positions in the mask are to be affected.

- If bit 7 is a 1, then any bit 6-0 in your own mask word which is set to a 1 SETS the corresponding bit in the MASK register. Any bit which you have set to a 0 causes the MASK register bit to remain in its current state.
- If bit 7 is a 0, then any bit 6-0 in your own mask word which is set to a 1 CLEARS the corresponding bit in the MASK register. Again, any 0 bit in your own mask word causes no change in the contents of the corresponding MASK register bit.

If an interrupt is to occur based on a particular condition, then that corresponding MASK bit must be a 1.

Example: Suppose you want to SET the TIMER A interrupt bit (enable the TIMER A interrupt), but want to be sure that all other interrupts are CLEARED. Here is the sequence you can use:

	01111110B,A0	
mov.d	A0,ICR	;MSB is 0, means clear ;any bit whose value is
		;1 in the rest of the byte
movi.b	10000001B,A0	
mov.b	A0,ICR	;MSB is 1, means set ;any bit whose value is ;1 in the rest of the byte ;(do not change any values ; wherein the written value
		; bit i <b>s</b> a zero)
Interrupt Control	Register:	

REG NAME D7 D6 D5 D4 D3 D2 D1 D0

Read

---- ---- ---- ---- ----\_ \_ \_ \_ \_ \_ \_ FLG SP 0 0 ALRM TB TA ICR IR D Write Interrupt Control MASK: D5 D4 D3 D2 D1 D0 REG NAME D7 D6 \_ \_ \_ \_ \_ \_ \_ \_ \_ \_ \_ S/C х FLG SP ALRM TB TA D ICR x

### CONTROL REGISTERS

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There are two control registers in the 8520, CRA and CRB. CRA is associated with TIMER A and CRB is associated with TIMER B. The format of the registers is as follows:

CONTROL REGISTER A:

- BIT NAME FUNCTION
  - 0 START 1 = start TIMER A, 0 = stop TIMER A. This bit is automatically reset (= 0) when underflow occurs during one-shot mode.
  - 1 PBON 1 = TIMER A output on PB6, 0 = PB6 is normal operation.
  - 2 OUTMODE 1 = TOGGLE, 0 = PULSE.
  - 3 RUNMODE 1 = one-shot mode, 0 = continuous mode.
  - 4 LOAD 1 = FORCE LOAD (this is a STROBE input, there is no data storage; bit 4 will always read back a zero and writing a 0 has no effect.)
  - 5 INMODE 1 = TIMER A counts positive CNT transitions, 0 = TIMER A counts 02 pulses.
  - 6 SPMODE 1 = SERIAL PORT=output (CNT is the source of the shift clock) 0 = SERIAL PORT=input (external shift clock is required)

BIT MAP OF REGISTER CRA:

REG #	NAME	TOD IN	SPMODE	INMODE	LOAD	RUNMODE	OUTMODE	PBON	START
E	CRA		0=input 1=output				•	0=PB60FF 1=PB60N	•

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(STROBE) shot |<----- TIMER A Variables ----->| All unused register bits are unaffected by a write and forced to 0 on a read. CONTROL REGISTER B: FUNCTION BIT NAME 1 =start TIMER B, 0 =stop TIMER B. 0 START This bit is automatically reset (= 0) when underflow occurs during one-shot mode. 1 = TIMER B output on PB7, 0 = PB7 is normal operation. PBON 1 1 = TOGGLE, 0 = PULSE.OUTMODE 2 1 = one-shot mode, 0 = continuous mode.RUNMODE 3 1 = FORCE LOAD (this is a STROBE input, there is no 4 LOAD data storage; bit 4 will always read back a zero and writing a 0 has no effect.) Bits CRB6 and CRB5 select one of four possible input 6,5 INMODE modes for TIMER B, as follows: CRB6 CRB5 Mode Selected \_ \_ \_ \_ \_ \_ \_ \_ 0 0 TIMER B counts 02 pulses TIMER B counts positive CNT transitions 0 1 TIMER B counts TIMER A underflow pulses 0 1 TIMER B counts TIMER A underflow pulses 1 1 while CNT pin is held high. 1 = writing to TOD registers sets ALARM, 7 ALARM 0 = writing to TOD registers sets TOD clock. Reading TOD registers always reads TOD clock, regardless of the state of the ALARM bit. BIT MAP OF REGISTER CRB: RUNMODE OUTMODE PBON START REG # NAME ALARM INMODE LOAD 0=pulse 0=PB70FF CRB 0=TOD 00=02 1=FORCE 0=cont. 0=stop 1=ALARM 01=CNT LOAD 1=one-1=toggle 1=PB7ON 1=start shot 10=TIMER A (STROBE) 11=CNT+TIMER A

|<-----TIMER B Variables----->|

All unused register bits are unaffected by a write and forced to 0 on a read.

# PORT SIGNAL ASSIGNMENTS

This part specifies how various signals relate to the available ports of the 8520. This information enables the programmer to relate the port addresses to the outside-world items (or internal control signals) which are to be affected. This part is primarily for the use of the systems programmer and should generally not be utilized by applications programmers. Systems software normally is configured to handle the setting of particular signals, no matter how the physical connections may change. In other words, if you have a version of the system software that matches the rev. level of the machine (normally a true condition), when you ask that a particular bit be set, you don't care which port that bit is connected to. Thus applications programmers should rely on system documentation rather than going directly to the ports. Note also that in this, a multi-tasking operating system, many different tasks may be competing for the use of the system resources. Applications programmers should follow the established rules for resource access in order to assure compatibility of their software with the system.

Address BFERFF	data bits 7-0 (Al2*) (int2)
PA6game port PA5RDY* PA4TK0* PA3WPRO*	1, pin 6 (fire button*) 0, pin 6 (fire button*) disk ready* disk track 00* write protect* disk change* led light (0=bright) memory overlay bit
SPKDAT CNTKCLK	keyboard data
PB7P7 PB6P6 PB5P5 PB4P4 PB3P3 PB2P2 PB1P1 PB0P0	data 7 data 6 data 5 Centronics parallel interface data 4 data 3 data 2 data 1 data 0
PCdrdy* Fack*	centronics control

Address BFDRFE data bits 15-8 (A13\*) (int6)

PA7..com line DTR\*, driven output

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PA6..com line RTS\*, driven output PA5..com line carrier detect\* PA4..com line CTS\* PA3..com line DSR\* PA2.SEL centronics control paper out ---+ PA1..POUT ---+ | PA0.BUSY busy Т commodore -+ SP...BUSY commodore ---+ CNT. . POUT PB7..MTR\* motor PB6..SEL3\* select external 3rd drive select external 2nd drive PB5..SEL2\* select external 1st drive PB4..SEL1\* select internal drive PB3..SEL0\* PB2..SIDE\* side select\* direction PB1..DIR step\* PB0..STEP\* PC...not used F....INDEX\* disk index\* .fi

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### AMIGA AUTO-CONFIGURATION ARCHITECTURE

This document is is process. Please contact Commodore-Amiga for the latest available information on this topic.