# Beyond Games: Systems Software for Your 



## Personal Computer



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Ken Skier

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## Beyond Games: Systems Software for Your 6502 Personal Computer

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Library of Congress Cataloging in Publication Data
Skier, Ken.
Beyond games.
Includes index.

1. 6502 (Computer)-Programming. I. Title.

QA76.8.S63S59 001.64'2 80-28512
ISBN 0-07-057860-5

Cover photo by Bob Hamor

Text set in Paladium
by BYTE Publications
Edited by Blaise Liffick
Design and Production Supervision by Ellen Klempner
Copy Edited by Rich Friedman
Printed and bound using 50\#MH Matte

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## Introduction

## Objectives

Sometimes I hear people talk about how smart computers have become. But computers aren't smart: programmers are. Programmers make microprocessors act like calculators, moon landers, or income tax preparers. Programmers must be smart, because by themselves microprocessors can't do much of anything.

Sound programming, then, is fundamental to successful computer use. With this principle in mind, this book has two objectives: first, to introduce newcomers to some of the techniques, terminology, and power of assembly-language programming in general, and of the 6502 in particular; and second, to present a set of software tools to use in developing assembly-language programs for the 6502.

Chapter 1 takes you on a quick tour of your computer's hardware and software; Chapters 2 thru 4 comprise a short course in assembly-language programming for those readers new to the subject. The rest of the book presents source listings, object code, and assembler listings for programs that you may enter into your computer and run.

Programmers have long sought to develop small and fast programs with the unfortunate result that occasionally code has been written that is unreadable (and even unworkable) simply because a programmer wanted to save a few bytes or a few cycles. In certain instances when memory space is particularly tight or execution time is critical, readability is sacrificed for performance. But today the average programmer is not forced to make this choice. Of course, all other things being equal, I, too, value programs that are quick and compact.

But how often are all other things equal?
While developing the programs that appear in this book, I had a number of objectives, most of them more important than the speed or size of a block of code. I designed these programs to be:

Useful: No program is presented simply to demonstrate a particular program-
ming technique. All of the programs in this book were written because I needed certain things done - usually something I didn't want to be bothered with doing myself. The monitor monitors, the disassembler disassembles, and the text editor lets me enter and edit text strings. These programs earn their keep.

Easy to Use: Simply by glancing at the screen you can tell which program is running and what mode it is in. When a program needs information, it asks you for it and allows you to correct mistakes you might make while answering. This software doesn't require you to remember the addresses of programs or of variables. Functions are mapped to individual keys, and you can assign functions to keys in any way that makes sense to you.

Readable: A beginning 6502 programmer should be able to understand the workings of every program in this book. The labels and comments in the listings were carefully chosen to reveal the purpose of each variable, subroutine, and line of code. I am writing first and foremost for you, the reader, not for the 6502.

Portable: The book's software runs on an Apple II, an Atari 400 or 800, an Ohio Scientific (OSI) Challenger I-P, or a PET 2001. With proper initialization of the System Data Block, it should run on any 6502 -based computer equipped with a keyboard and a memory-mapped, character-graphics video display.

Compatible: These routines are very good neighbors. As long as the other software in your system does not use the second 4 K bytes of memory (hexadecimal memory locations 1000 thru 1FFF), there should be no conflict between your software and the software in this book. In particular, most of the software in this book preserves the zero page, so your software may use the zero page as much as you like, and you won't be bothered with having to save and restore it before and after calls to the software presented herein.

Expandable: The programs in this book are highly modular, and you may extend or restructure them to meet your individual needs. System-specific subroutines are called indirectly, so that other subroutines may be substituted for them, and most values are treated as variables, rather than as constants hard-wired into the code. There are no monolithic programs in this book; they're all subroutines and may be combined in many ways to build powerful new structures.

Compact: I know that every personal computer has exactly the same available memory: too little. I also know ways to write a program in ten or twenty percent less space. But if doing so required sacrificing readability, portability, or expandability, I did not do so. In many cases I feared that to save a byte, I might lose a reader's clear understanding of how a program works. I considered that too great a price to pay for a somewhat smaller program.

Fast: Assuming that the above objectives have been met, the software in this book has been developed to operate as quickly as possible. But in any trade-off between speed and the other objectives, speed loses. A fast program that you can't understand holds little value. None of the programs in this book are likely to make you complain about how long you have to wait. I can't tell if I'm waiting an extra millisecond. Can you?

So go ahead. Read. Program. Enjoy!

## Chapter I:

## Your Computer

The software in this book can run on a number of computers because it assumes very little about the host machine. Let's examine these assumptions and in so doing take a quick tour of your computer.

## The 6502 Microprocessor

We'll start with the 6502 microprocessor, the component in your system that actually computes. By itself, the 6502 can't do much. It has three registers (special memory areas for storing the data upon which the program is operating), called A, $X$, and $Y$, which can each hold a number in the range of 0 to 255 . Different registers have different capabilities. For example, if a number is in A (the accumulator), the 6502 can add to it, or subtract from it, any value up to 255 . But if a number is in the $X$ register or the $Y$ register, the 6502 can only increment or decrement that number (ie: add or subtract one from it).

The 6502 can also set one register equal to the value of another register, and it can store the contents of any register anywhere in memory, or load any register from any location in memory. Thus, although the 6502 can only operate on one number at a time, it can operate on many numbers, just by loading registers from various locations in memory, operating on the registers, and then storing the results of those operations back into memory.

## Types of Memory

You may have heard that a computer stores information as a series of ones and
zeros. This is because the computer's memory is simply an elaborate array of switches, and an individual switch can have only two states: closed or open. These two states may also be expressed as on and off, or as one and zero.

Not all memory switches are the same. Some, in what is called ROM (read-only memory), are hard-wired into your computer's circuitry and cannot be changed except by physically replacing the ROM circuits containing those switches. Others, in what is called RAM (random-access memory) or programmable memory, can be changed by the processor. The 6502 can open or close any of the switches, called bits (binary digits), in its programmable memory, and later on read what it "wrote" into that memory. Figure 1.1 shows how the processor has access to read-only memory and programmable memory.


Figure 1.1: How the 6502 interacts with memory. The arrows indicate the flow of data.

A third kind of memory is set by some external device, not by the 6502. Such memory switches are called input ports, and may be connected to keyboards, terminals, burglar alarms - virtually anything that can generate an electrical signal. The 6502 perceives these externally generated signals by reading the appropriate input ports.

Yet another kind of memory switch, called an output port, generates a high or a low voltage on some particular wire depending on whether the 6502 sets a given memory switch to a one or a zero. One or more of these output ports can enable the 6502 to "talk" to the outside world.

Now don't jump up and think I'm going to show you how to synthesize speech in this book. "Talk" is just my way of anthropomorphizing the 6502. It will happen elsewhere in this book, when the 6502 "sees," "remembers," and "knows" what to do. Of course the 6502 doesn't see, remember, or know anything, but I often find it helpful to put myself in its place. That way I can better understand how a program will run, or why a program doesn't run, and I do see, remember, and know things.

But don't take such verbs too literally. The 6502 doesn't talk. It causes signals to be generated that may be sensed by other devices, such as cassette recorders, printers, disk drives - and yes, even speech synthesizers. But not in this book.

Some peripheral devices are actually connected to both an input and an output port. Examples of these devices are cassette tape machines and floppy-disk drives,
which are mass-storage or secondary-storage devices. Figure 1.2 summarizes the processor's access to memory and to peripheral devices.

PERIPHERALS MEMORY PROCESSOR


Figure 1.2: A summary of the 6502 microprocessor's access to data in main memory and through I/O (input and output) ports. The arrows indicate the flow of data.

A video screen connected to your computer looks like memory to the 6502, so the 6502 can read from and write to the screen. The keyboard is scanned by I/O (input/output) ports that are decoded to look like any other programmable memory
address, so the 6502 can look at the keyboard just by looking at a particular place in memory. Thus, the 6502 can interact directly with memory only, but because all I/O devices are mapped to addresses in memory, the 6502 can interact with the user. See figure 1.3.


Figure 1.3: How the 6502 interacts with the user. Arrows indicate the flow of data.

## The Operating System

Thus far we have discussed your machine's hardware. But the Apple, Atari OSI, and PET computers feature more than hardware. For example, all these computers have an operating system (stored in ROM) which includes the I/O software routines that are needed to use the screen and the keyboard. We are not particularly concerned with how these subroutines work, but I assume your system does have such routines.

There are many other subroutines in your computer's operating system. Your system's documentation should tell you what subroutines are available and provide their addresses. All of this means power for you, the programmer. The more you know about your computer, the more you can make it do. Because the software in this book was developed to run on a number of systems, I chose not to use routines available in your machine's ROM, no matter how powerful they might be, unless I could be sure that they would be available in the operating systems of the Apple, the Atari, the OSI, and the PET computers. In other words, the software in this book does not take full advantage of the power in your operating system. But the software you write, which need only run on your system, should exploit to the fullest the power of your computer's ROM routines.

## BASIC

One of the most important features of your computer is the BASIC interpreter in ROM. This interpreter is a program that enables your computer to understand commands given in BASIC. Your system's documentation should tell you what commands are legal in the particular dialect of BASIC implemented on your machine. BASIC is an easy language to learn and you can do a lot with it.

Unfortunately, not every dialect of BASIC is the same. A program written in BASIC that runs on machine A may not run on machine B. BASIC is a common language, but not a standard one. Is there any language that is standard from system to system?

## 6502 Code

The central processor is the computer's heart. The Apple, Atari, OSI, and PET computers all use the 6502 microprocessor. Every microprocessor has a certain instruction set, or group of instructions, which the microprocessor can execute. These instructions are at a much lower level than the BASIC commands with which you may be familiar. For example, in BASIC you can have a single line in a program to PRINT "HELLO." It would take a sequence of many 6502 instructions to perform the same function.

However, a sequence of microprocessor instructions will run on any computer featuring that microprocessor. Thus, if you write a program consisting of 6502 instructions to perform some function, that program should run on any 6502-based computer. It won't run on an 8080 -based computer, a Z80-based computer, or a 6800 -based computer, but it should run on an Apple, a PET, an Atari, an OSI, or any other system built around a 6502.6502 programs can also run much faster than equivalent programs written in BASIC and can be smaller than BASIC programs. The programs presented in this book are all written in 6502 code, and require only half of the memory available on a computer containing 8,000 bytes of programmable memory, thus leaving more than enough room for your own programs.

## Chapter 2:

## Introduction to Assembler

Ever watch a juggler or a good juggling team? The balls, pins, or whatever are in the air in such intricate patterns that you can hardly follow them, let alone duplicate the performance yourself. It's beautiful, but not magic; just an application of some simple rules. I've learned to juggle recently, and although I'm still a rank beginner, I've taught my two hands to keep three balls moving through the air. Yet neither hand knows very much. A hand will toss a ball into the air, and then it will catch a ball. The other hand will toss a ball into the air, and then it will catch a ball. That's all. My hands perform only two operations: toss and catch. Yet with those two primitive operations I can put on a pleasant little performance.

Assembly-language programming is not so different from juggling. Like juggling, programming enables you to put on an impressive or baffling performance. In its simplest terms, juggling is nothing more than taking something from one place and putting it someplace else. The same thing is true of the central processor: the 6502 takes something from one place and puts it someplace else.

In fact, programming the 6502 is easier than juggling in several ways. First, the 6502 is obviously much faster than even the most skillful juggler. In the time it takes me to pick up a ball with one hand and place that ball somewhere else, the 6502 can get something from one place and put it someplace else hundreds of thousands of times. Sleight of hand requires quickness, and the 6502 is quick.

The 6502 even gives me a helping hand. When I try to juggle, I must keep the balls moving with nothing but my two hands. But my home computer has three hands (registers $A, X$, and $Y$ in the 6502) and thousands of pockets ( 8,000 bytes or more of programmable memory).

A byte is 8 bits of data that may be loaded together into a register. A register holds 1 byte. Each location in memory holds 1 byte. The 6502 can affect only 1 byte in one operation. But because the 6502 can perform hundreds of thousands of opera-
tions each second, it can affect hundreds of thousands of bytes each second.

## Binary

In the final analysis, any value is stored within the computer as a series of bits. If we wish, we may specify a byte by its bit pattern: such a representation uses only ones and zeroes, and is called binary. For example, the number 25 in binary is 00011001.

In binary, each bit indicates the presence or absence of some value. Each bit represents twice as much value, or significance, as the bit to its right, so the rightmost bit is the least significant, and the left-most bit is the most significant. Table 2.1 gives the significance of each bit in an 8 -bit byte:

Table 2.1: Bit significance in an 8 -bit byte.

| Bit Number: | b7 | b6 | b5 | b4 | b3 | b2 | b1 | b0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Bit Significance: | 128 | 64 | 32 | 16 | 8 | 4 | 2 | 1 |

The right-most bit (called bit 0 ) tells us whether we have a one in our byte. The bit to its left (bit 1) tells us whether we have a two; the bit to its left tells us whether we have a four...and the leftmost bit (bit 7) tells us whether we have a 128 in our byte.

To determine the bit pattern for a given value - say, 25 - determine first what powers of two must be added to equal your value. For instance, $25=16+8+1$, so 25 in binary is 00011001 .

Twenty-five can be expressed in other ways as well. Rather than specify every number as a pattern of eight ones and zeros, we often express numbers in hexadecimal representation.

## Hexadecimal

Unlike binary, which requires a group of eight characters to represent an 8 -bit value, hexadecimal notation allows us to represent an 8 -bit value with a group of only two characters. These characters are not limited to 0 and 1 , but may include any digit from 0 to 9 , and any letter from " A " to " F ." That gives us a set of sixteen characters, which is just right because we want to represent numbers in base 16.
(Hexadecimal stands for 16: hex for six, and decimal for ten. Six plus ten equals sixteen.)

To represent a byte in hexadecimal notation, divide the 8 -bit byte into two 4 -bit units (sometimes called nybbles). Each of these 4 -bit units has a value of from 0 to 15 (decimal), which we express with a single hexadecimal digit. A decimal 10 is a hexadecimal $\$ \mathrm{~A}$. (The dollar sign indicates that a number is in hexadecimal representation.) Table 2.2 gives the conversions of decimal to hexadecimal for decimal numbers 0 thru 15.

Table 2.2: Hexadecimal character set.

| Hexadecimal Character |  | Decimal Equivalent |
| :---: | :---: | :---: |
| $\$ 0$ | $=$ | 0 |
| $\$ 1$ | $=$ | 1 |
| $\$ 2$ | $=$ | 2 |
| $\$ 3$ | $=$ | 3 |
| $\$ 4$ | $=$ | 4 |
| $\$ 5$ | $=$ | 5 |
| $\$ 6$ | $=$ | 6 |
| $\$ 7$ | $=$ | 7 |
| $\$ 8$ | $=$ | 8 |
| $\$ 9$ | $=$ | 9 |
| $\$ A$ | $=$ | 10 |
| $\$ B$ | $=$ | 11 |
| $\$ C$ | $=$ | 12 |
| $\$ D$ | $=$ | 13 |
| $\$ E$ |  | 15 |

Appendix A1, Hexadecimal Conversion Table, shows the hexadecimal representation of every number from 0 to 255 decimal.

In this book, object code, the only code that the machine can execute directly, will generally be presented in hexadecimal, and a thorough understanding of hexadecimal will help you to interpret instructions and follow some of the 6502's actions. Even the sketchiest understanding of hexadecimal math, however, should be sufficient for you to follow and use the programs in this book.

## ASCII Characters

Instead of a number from 0 to 255, an 8-bit byte can be used to represent an upper or lower case letter of the alphabet, a punctuation mark, or a printer-control character such as a carriage return. A string of such bytes may represent a word, a message, or even a complete document. Appendix A2, ASCII Character Codes, gives the hexadecimal value for any ASCII character. ASCII stands for American Standard Code for Information Interchange, and is the closest thing the industry has to a standard set of character codes. If you want to store the letter " $A$ " in some location in memory, you can see from Appendix A2 that you must store a $\$ 41$ in that location.

Whether a given byte is interpreted as a number, an ASCII character, or something else depends entirely on the program using that byte. Just as beauty is in the eye and mind of the beholder, so is the meaning of a given byte determined by the program that sees and uses it.

## The Instruction Cycle

A microprocessor such as the 6502 can't do anything without being told. It only knows 151 instructions, called opcodes (operation codes). Each opcode is 1 byte long. An opcode may command the 6502 to take something from one register and to put it someplace in memory, to load some register with the contents of some location in memory, or to perform some other equally simple operation. See Appendix A4 for a list of opcodes for the 6502 microprocessor.

What do 6502 s do all day? They work while programmers play. The 6502 gets an opcode, performs the specified operation, gets the next opcode, performs the specified operation, gets the next opcode, performs the...

You get the picture.
How does the 6502 know where to find the next opcode? The 6502 has a 16-bit register called the PC (program counter). The PC holds the address of some location in memory. When the 6502 starts its instruction cycle, it gets the opcode stored at the memory location specified by the PC. Then it performs the operation specified by that opcode. When it has executed that instruction, it makes the PC point to the next opcode and starts on a new instruction cycle by getting the opcode whose address is now in the PC.

Figure 2.1 shows a flowchart for the instruction cycle of the 6502 microprocessor.
"That's it? That's all the 6502 does?" you ask.
That's it. But with the right program in memory, we can make the 6502 dance.


Figure 2.1: The 6502 instruction cycle.

## Machine Language

A machine-language program is nothing more than a series of machinelanguage instructions stored in memory. If the PC in the 6502 can be made to hold the address of the start of your program, then we say that the PC is pointing to your program. When the 6502 starts its instruction cycle, it will fetch the first opcode in your program, and then perform the operation specified by that opcode. At this point, we say that your program is running.

Each machine-language instruction is stored in memory as a 1-byte opcode, which may be followed by 1 or 2 bytes of operand. Thus, a 6502 machine-language program might be "A9 05200204 A2 F5 60."

Just a bunch of numbers! (Hexadecimal numbers, in this case.) But it is exactly these numbers that the machine understands; hence the term, machine language.

## Assemblers

Machine language is easy to read - if you're a machine. But programmers are people. So programming tools called assemblers have been developed, which take more readable assembly-language source code as input and produce listings and object code as output. The listing is the assembler's output intended for a human reader. The object code is a series of 6502 machine-language instructions intended to be stored in memory and executed by the 6502 .

For each chapter in this book that presents a program, there is an appendix at the back of the book containing an assembler listing and a hexdump of the same program. The assembler listing includes both source and object code, making it easy for you to read the program; the hexdump shows you what the object code for that program actually looks like in your computer's memory. Figure 2.2 shows how an assembler is used to produce an assembler listing for the programmer and object code for the processor.


Figure 2.2: From programmer to object code. The assembler takes source code as input and produces an assembler listing and object code as output.

The programs in this book have all been produced on the OSI 6500 Assembler/Editor, running under the OSI 65-D Disk Operating System, on an OSI C-IP machine with 24 K bytes of programmable memory and one 5 -inch floppy disk. It is likely that the source code presented in this book will assemble immediately or with only minor modification on other 6500 assemblers. (Incidentally, the source code in each chapter of this book should fit into the workspace of a computer with much less than 24 K bytes of user memory, if you delete many of the comments. But then, of course, your listings will be a lot less readable.)

But you don't write a listing; an assembler produces a listing. What you write is assembly-language source code.

## Source Code

An assembly-language source program consists of one or more lines of
assembly-language source code. A line of assembly-language source code consists of up to four fields:

## LABEL MNEMONIC OPERAND COMMENT

The mnemonic, required in all cases, is a group of three letters chosen to suggest the function of a given machine-language instruction. For example, the mnemonic $L D A$ stands for $L$ oa $D$ Accumulator. $L D X$ stands for $L$ oa $D X$ register. TXA means Transfer the $\chi$ register to the Accumulator. 6502 mnemonics are not nearly as meaningful as BASIC commands, but they're a big improvement over the machinelanguage opcodes. See Appendix A3 for a list of 6502 mnemonics.

Some operations require an operand field. For example, the operation load accumulator requires an operand, because the line of source code must specify what you wish to load into the accumulator.

The label and comment fields are optional. A label lets you operate on some location in memory by a name that you have assigned to it. Comments are not included in the object code that will be assembled from your program, but they make your source code and your listings much more meaningful to a human reader. When you write a program, even if no one but yourself will ever read it, try to choose your labels and comments so that someone else can understand the purpose of each part of the program. Such careful documentation will save you a lot of time weeks or months down the road, when you might otherwise reread your program and have no idea why you included some unlabeled, uncommented line of source code.

## Loading a Register

Let's write a simple program to load a register with a number - say, to load the accumulator with the number "10." Since we want to load the accumulator, we'll use the LDA instruction. (If we wanted to load the $X$ register, we would use the LDX instruction, and if we wanted to load the $Y$ register, we'd use LDY.) We know what mnemonic to write into our first line of source code. But a glance at Appendix A6, 6502 Opcodes by Mnemonic and Addressing Mode, shows that LDA has many addressing modes. What operand shall we write into this line of source code?

We know that we want to load the accumulator with a " 10 ," and not with any other number, so we can use the immediate addressing mode to load a " 10 " directly into the accumulator. We'll use a " $\#$ " sign to indicate the immediate mode:

## Example I

LDA \#10

Example 1 is a legitimate line of source code containing only two fields: a mnemonic and an operand. The mnemonic, LDA, means "load the accumulator." But load it with what? The operand tells us what to load into the accumulator. The "\#" sign specifies that this operation is to take place in the immediate mode, which means we want to load the accumulator with a constant to be found in this line of source code, rather than with data or a variable to be found in some location in memory. Then the operand specifies the constant to be loaded into the accumulator, in this case " 10 ."

## Constants

A constant is any value that is known by the programmer and "hard-wired" into the code. A constant does not change during the execution of a program. If a value changes during the execution of a program, then it is a variable, and one or more memory locations must be allocated to hold the current value of each variable.

There are several kinds of constants. Any number is a constant. The number " 7 ," for example, is a constant: a seven now will still be a seven this afternoon. A character is another kind of constant: the letter " A " will still be the letter " A " tomorrow. But a variable, such as one called FUEL, will change during the course of a program (such as a lunar lander simulation), so it is not a constant.

In Example 1, note that the " $\#$ " sign is the only punctuation in the operand field. In the absence of special punctuation marks (such as the dollar sign indicating a hexadecimal number and the apostrophe indicating an ASCII character representation), any numbers given in this book are in decimal.

What object code will be assembled from this line of source code? Let's handassemble it and see. Appendix A6 shows us that the opcode for load accumulator, immediate mode, is $\$ A 9$. So the first byte of object code for this instruction will be $\$$ A9. The second byte must specify what the 6502 should load into the accumulator. We want to load register A with a decimal 10 , which is $\$ 0 \mathrm{~A}$. So the object code assembled from Example 1 is: A9 0A.

When these 2 bytes of object code are executed by the 6502, it will result in the accumulator holding a value of $\$ 0 \mathrm{~A}$, or decimal 10 . In effect, we've just told a juggler: put a " 10 " in your right hand.

What if we wanted to load the accumulator with the letter " $M$," rather than with a number? We'd still use LDA to load the accumulator, and we'd still use the immediate mode of addressing, specifying in the operand the constant to be loaded into the accumulator. Either of the following two lines of source code will work:

## Example 2

## LDA \#' M

or
LDA \#\$4D

In each line of source code above, the mnemonic and the " $\#$ " sign tell us we're loading the accumulator in the immediate mode - ie: with a constant. The operand following the " $\#$ " sign specifies the constant. An apostrophe indicates that an ASCII character follows, whereas a " $\$$ " sign indicates that a hexadecimal number follows. Appendix A2 shows that an ASCII " M " $=\$ 4 \mathrm{D}$; they are simply two representations of the same bit pattern. So the two lines of source code above are equivalent; they will both assemble into the same object code: A9 4D.

Which of the two lines of source code is more readable? If a constant will be used in a program as an ASCII character, then represent it in your source code as an ASCII character.

## Storing the Register

Now let's say we want to store the contents of the accumulator someplace in memory. Every location in memory has a unique address (just like houses do), ranging from $\$ 0000$ to $\$$ FFFF. Suppose we decide to store the contents of the accumulator at memory location $\$ 020 \mathrm{C}$. We could do it with the following line of source code:

## Example 3

STA \$020C

Example 3 will assemble into these 3 bytes of machine language: 8D 0C 02 .
According to the Appendix A6, the 6502 opcode for "store accumulator, absolute mode" (STA) is $\$ 8 \mathrm{D}$.

When the 6502 fetches the opcode " 8 D, " it knows that it must store the contents of the accumulator at the address specified by the next 2 bytes. This is why it is called absolute mode. Absolute mode is used when specifying an exact memory location in an instruction.

In the example above, that address seems wrong. It looks like the machinelanguage operand is specifying address $\$ 0 \mathrm{C} 02$, because the bytes are in that order: " $0 C$ " followed by " 02 ." But we want to operate an address $\$ 020 \mathrm{C}$. Is something wrong here?

## Low Byte First

You and I might think something is wrong when the address $\$ 020 \mathrm{C}$ is written as an " $0 C^{\prime \prime}$ followed by an " 02 " but you and I are people. We don't think like the 6502. When you and I write a number, we tend to write the most significant digit first and the least significant digit last. But the 6502 doesn't work that way. When the 6502 interprets two sequential bytes as an address, the first byte must contain the less significant part of the address (the "low byte"), and the second byte must contain the more significant part of the address (the "high byte"). All addressing modes that require a 2 -byte operand require that the 2 bytes be in this order: less significant byte first, followed by the more significant byte.

However, not all addressing modes require a 2 -byte operand.

## Zero-Page Addressing

Memory is divided into pages, where a page is a block of 256 contiguous addresses. The page from $\$ 0000$ to $\$ 00 \mathrm{FF}$ is called the zero page, because all addresses in this page have a high byte of zero. The zero-page addressing mode takes advantage of this fact. Source code assembled using the zero-page addressing mode requires only 1 byte in the operand, because the opcode specifies the zero page mode of addressing, and the high byte of the operand is unnecessary because it is understood to be zero. Thus, you can specify an address in the zero page by the absolute or by the zero-page addressing mode, but the zero-page mode will let you do it using one less byte.

If you want to use some location in the zero page to hold a number, you might decide to use location $\$ 00 \mathrm{~F} 4$. We could write:

## Example 4

STA \$00F4
or
STA \$F4

We could then assemble either line of source code using the absolute addressing mode: 8D F4 00. Or we could assemble either line of source code using the zeropage mode: 85 F4.

The opcode " 85 " means "store accumulator, zero page." Where in the zero page? At location $\$ F 4$ in the zero page, the same location whose absolute address is $\$ 00 \mathrm{~F} 4$.

## Symbolic Expressions

Let's say you want to copy the 3 bytes at memory locations $\$ 0200, \$ 0201$, and $\$ 0202$ to $\$ 0300, \$ 0301$, and $\$ 0302$, respectively. We could write these lines of source code:

## Example 5

LDA $\$ 0200$
STA $\$ 0300$
LDA \$0201
STA \$0301
LDA \$0202
STA \$0302

This alternately loads a byte into the accumulator, then stores the contents of the accumulator into another byte in memory. Note that loading a register from a location in memory changes the register, but leaves the contents of the memory location unchanged.

Or we could write the following code, which refers to addresses as symbolic expressions:

## Example 6

| 1 | ORIGIN $=\$ 0200$ |  |
| :--- | :--- | :--- |
| 2 | DEST | $=\$ 0300$ |
| 3 | LDA | ORIGIN |
| 4 | STA | DEST |
| 5 | LDA | ORIGIN +1 |
| 6 | STA | DEST +1 |
| 7 | LDA | ORIGIN +2 |
| 8 | STA | DEST +2 |

In Example 6, lines 1 and 2 are assembler directives, which equate the labels "ORIGIN" and "DEST" with the addresses $\$ 0200$ and $\$ 0300$, respectively. Other lines of source code following these equates may then refer to these addresses by their labels, or refer to any address as a symbolic expression consisting of labels and, optionally, constants and arithmetic operators. The source code above will cause an assembler to generate exactly the same object code as the source code in Example 5, but Example 6, whose operands consist of symbolic expressions, is much more
readable than Example 5, whose operands are given in hexadecimal.

## Some Exercises

1) Write the 6502 instructions necessary to load the accumulator with the value 127, to load the $X$ register with the letter " $r$," and to load the $Y$ register with the contents of address \$BO92.
2) Write the 6502 instructions necessary to copy the byte at address $\$ 0043$ to the address $\$ 0092$.

## Chapter 3:

## Loops and Subroutines

## Indexed Addressing

Although readable, Example 6 is not very efficient, because it requires two lines of source code to move each byte. If we want to move 50 or 100 bytes must we then write 100 or 200 lines of source code?

Indexed addressing comes in quite handily here. Instead of specifying the absolute or zero-page address on which an operation is to be performed, we can specify a base address and an index register. The 6502 will add the value of the specified index registers to the base address, thereby determining the address on which the operation is to be performed. Thus, if we want to move 9 bytes from an origin to a destination, we could do it in the following manner, using the indexed addressing mode with X as the index register:

## Example 7

ORIGIN $=\$ 0200$
DEST $=\$ 0300$

| INIT | LDX \#0 |
| :--- | :--- |
| GET | LDA ORIGIN,X |
| PUT | STA DEST,X |

ADJUST INX

Initialize X register to zero, so we'll start with the first byte in the block.
Get Xth byte in origin block.
Put it into the Xth position in the destination block.
Adjust $X$ for next byte by incrementing (adding 1) to the $X$ register.

| TEST | CPX \#9 | Done 9 bytes yet? |
| :--- | :--- | :--- |
| BRANCH | BNE GET | If not, go back and get next byte... |

We will use Example 7 in the following sections to introduce several new instructions and addressing modes. Example 7 includes six lines of source code to move 9 contiguous bytes of data. If we tried to move 9 bytes of data with the techniques used in Examples 5 and 6, it would have taken eighteen lines of source code. So with indexed addressing, we've saved ourselves twelve lines of code. But how do these lines work? The lines are labeled so we can look at them one-by-one.

The instruction labeled INIT loads the $X$ register in the immediate mode with the value zero. After executing the line INIT, the 6502 has a value of zero in the $X$ register. We don't know anything about what's in the other registers.

GET loads the accumulator with the Xth byte above the address labeled ORIGIN. The first time the 6502 encounters this line, the $X$ register will hold a value of zero, so the 6502 will load the accumulator with the zeroth byte above the address labeled ORIGIN (ie: it will load the accumulator with the contents of the memory location ORIGIN).

In any line of source code, a comma in the operand indicates that the operation to be performed shall use an indexed addressing mode. A comma followed by an " $X$ " indicates that the X register will be the index register for an instruction, whereas a comma followed by a " $Y$ " indicates that the $Y$ register will be the index for an instruction. There are a number of indexed addressing modes. Two of these are absolute indexed and zero-page indexed. The line GET in Example 7 uses the absolute indexed addressing mode if ORIGIN is above the zero page; if ORIGIN is in the zero page then the line labeled GET can be assembled using the zero-page indexed addressing mode. Zero-page indexed addressing, like zero-page addressing, requires only 1 byte in the operand.

In zero-page indexed and in absolute indexed addressing, the operand field specifies a base address. The 6502 will operate on an address it determines by adding to the base address the value of the specified index register ( X or Y ). Only if the specified index register has a value of zero will the 6502 operate on the base address itself; in all other cases the 6502 will operate on some address higher in memory.

So we've loaded the accumulator with the byte at ORIGIN. Now the 6502 reaches the line labeled PUT in Example 7. This line tells the 6502 to store the accumulator in the $X$ th byte above DEST. We haven't done anything to change $X$ since the line INIT set it to zero, so X still holds a value of zero. Therefore, the 6502 will store the contents of the accumulator in the zeroth byte above DEST (ie: in DEST itself).

At this point, we have succeeded in moving 1 byte from ORIGIN to DEST. $X$ is still zero. Now comes the part that makes indexing worthwhile. The line labeled ADJUST is the shortest line of source code we've seen yet, consisting only of the mnemonic INX, which means "increment the $X$ register." Since the $X$ register was zero, when this line is executed the $X$ register will be left holding a value of one.

## Compare Register

In Example 7, the line labeled TEST compares the value in the X register with the number " 9 ." There are three compare instructions for the 6502, one for each register. CMP compares a value with the contents of the accumulator; CPX compares a value with the contents of the X register, and CPY compares a value with the contents of the $Y$ register.

We can use these compare instructions to compare any register with any value in memory, or, in the immediate mode, to compare any register with any constant. Such comparisons enable us to test for given conditions. For example, in Example 7, the line labeled TEST tests to see if we've moved 9 bytes yet. If the $X$ register holds the value " 9, " then we have moved 9 bytes. (Walk through the loop yourself. When you have moved the zeroth through the eighth bytes above ORIGIN to the zeroth through the eighth positions above DEST, then you have moved 9 bytes.)

A compare instruction never changes the contents of a register or of any location in memory. Thus, the $X$ register does not change when the line labeled TEST is executed by the 6502. What may change, however, are some of the 6502's status flags.

## Status Flags

In addition to the 6502's general-purpose registers ( $\mathrm{A}, \mathrm{X}$, and Y ), the 6502 contains a special register $P$, the processor status register. Individual bits in the processor status register are set or cleared each time the 6502 performs certain operations. These bits, or hardware flags, are:

| C | bit 0: Carry Flag |
| :--- | :--- |
| Z | bit 1: Zero Flag |
| I | bit 2: Interrupt Flag |
| D | bit 3: Decimal Flag |
| B | bit 4: <br> Break Flag |
| V | bit 5: Undefined |
| N | bit 6: Overflow Flag |
| N | bit 7: Negative Flag |

In this book, we will not discuss the use of all the flags in the processor status register. In this quick course in assembly-language programming, and in the software subsequently presented in this book, the three flags we will deal with are C , the
carry flag; $Z$, the zero flag; and $N$, the negative flag.
A compare operation (CMP, CPX, or CPY) does not change the value of registers A, X, or Y, but it does affect the carry, zero, and negative flags.

For example, if a register is compared with an equal value, the zero flag, $Z$, will be set; otherwise, $Z$ will be cleared. If an instruction sets bit 7 of a register or an address, the negative flag of the status register will also be set; conversely, if an instruction clears bit 7 of a register or an address, the negative flag will be cleared. Similarly, mathematical and logical operations set or clear the carry flag, which acts as a ninth bit in all arithmetic and logical operations. Table 3.1 summarizes the effects of a compare instruction on the status flags.

Table 3.1: Status flags affected by compare instructions. Note that if you wish to test the status of the carry flag after a compare, you must set it (using the instruction SEC) before the compare. When testing the $N$ flag, think of the inputs as signed 8 -bit values.

Carry Flag* Negative Flag Zero Flag

Compare a register with an equal value and you set $\mathrm{C}, \quad$ clear N, and set Z. Compare a register with a greater value and you Compare a register with a lesser value and you

| set $C$, | clear $N$, and | set $Z$. |
| :---: | :---: | :---: |
| clear $C$, | clear $N$, and | clear $Z$. |
| set $C$, | clear $N$, and | clear $Z$. |

## Conditional Branching

We can have a program take one action or another, depending on the state of a given flag. For example, two instructions, BEQ , (Branch on result $E$ Qual) and BNE (Branch on result Not Equal) cause the 6502 to branch, or jump to a new instruction, based on the state of the zero flag. An instruction which causes the 6502 to branch based on the state of a flag is called a conditional branch instruction. Other conditional branch instructions are based on the state of other status flags and are given in table 3.2.
*If you wish to test the status of the carry flag after a compare, you must set it (using the instruction SEC) before the compare.

Table 3.2: Conditional branch instructions.
$\left.\begin{array}{llll}\text { Flag } & \text { Instruction } & \text { Description } & \text { Opcode } \\ \text { C } & \text { BCC } & \text { Branch if carry clear. } & 90 \\ \text { C } & \text { BCS } & \begin{array}{l}\text { Branch if carry set. } \\ \text { N }\end{array} & \text { BPL }\end{array} \quad \begin{array}{l}\text { Branch if result positive. }\end{array}\right)$

The line labeled TEST in Example 7 compares the $X$ register to the value " 9 ;" this sets or clears the zero flag. The line labeled BRANCH then takes advantage of the state of the zero flag, by branching back to the line labeled GET if the result of that comparison was not equal. But if $Y$ did equal " 9 ," then the result of the comparison would have been equal, and the 6502 would not branch back to GET. Instead, the 6502 would execute the instruction following the line labeled BRANCH.

## Loops

Example 7 shows a program loop. We cause the 6502 to perform a certain operation many times, by initializing and then incrementing a counter, and testing the counter each time through the loop to see if the job is done.

There's a lot of power in loops. What would we have to add or change in Example 7 so that it moves not 9 , but 90 bytes from one place to another? Happily, we wouldn't have to add anything, and we'd only have to change the operand in the line labeled TEST. Instead of comparing the $X$ register with 9, we'd compare it with 90. See Example 8.

## Example 8

Move 90 bytes from origin to destination.
ORIGIN $=\$ 0200$
DEST $=\$ 0300$

| INIT | LDX \#O | Initialize $X$ register to zero, so we'll start <br> with the first byte in the block. |
| :--- | :--- | :--- |
| GET | LDA ORIGIN,X | Get Xth byte in origin block. <br> Put it into the Xth position in the |
| PUT | STA DEST, X | destination block. <br> ADJUST |
| INX | Adjust X for next byte. |  |
| TEST | CPX \#0 | Done 90 bytes yet? |
| BRANCH | BNE GET | If not, get next byte... |

Writing loops lets us write code that is not only compact, but easily tailored to meet the demands of a particular application. We couldn't do that, however, without indexing and branching.

Loops can be tricky, though. What's wrong with this loop?

## Example 9

$$
\begin{aligned}
& \text { ORIGIN }=\$ 0200 \\
& \text { DEST }=\$ 0300
\end{aligned}
$$

| INIT | LDX \#0 | Initialize $X$ register to zero, so we'll start <br> with the first byte in the block. |
| :--- | :--- | :--- |
| GET | LDA ORIGIN,X | Get Xth byte in origin block. <br> Put it into the Xth position in the |
| PUT | STA DEST, $X$ | destination block. |
| TEST | CPX \#9 | Done 9 bytes yet? |
| BRANCH | BNE GET | If not, get next byte... |

Examine Example 9 very carefully. How does it differ from Example 7? It lacks the line labeled ADJUST, which increments the X register. What will happen when the 6502 executes the code in Example 9 ? It will initialize $X$ to zero; it will get a byte from ORIGIN and move it to DEST. Then it will compare the contents of register $X$ to 9 . Register $X$ won't equal 9 , so it will branch back to GET, where it will do exactly what it did the first time through the loop, because X will still equal zero. Until the X register equals 9 , the 6502 will branch back to GET. But nothing in this loop will ever change the value of X! So the 6502 will sit in this loop forever, getting a byte from ORIGIN and putting it in DEST and determining that the X register does not hold a $9 . .$.

Now look at Example 10. Will it cause the 6502 to loop, and if so, will the 6502 ever exit from the loop? Why, or why not?

## Example 10

ORIGIN $=\$ 0200$
DEST $=\$ 0300$

INIT LDX \#0
GET LDA ORIGIN,X
PUT STADEST,X
ADJUST INX
TEST CPX \#9
BRANCH BNE INIT

Initialize $X$ register to zero, so we'll start with the first byte in the block. Get Xth byte in origin block. Put it into the Xth position in the destination block. Adjust $X$ for next byte. Done 9 bytes yet? If not, get next byte...

## Relative Addressing

All conditional branch instructions use the relative addressing mode, and they are the only instructions to use this addressing mode. Like the zero page and zeropage indexed addressing mode, the relative addressing mode requires only a 1 -byte operand. This operand specifies the relative location of the opcode to which the 6502 will branch if the status register satisfies the condition required by the branch instruction. A relative location of 04 means the 6502 should branch to an opcode 4 bytes beyond the next opcode, if the given condition is satisfied. Otherwise, the 6502 will proceed to the next opcode.

Because the operand in a conditional branch instruction is only 1 byte, it is not possible for a conditional branch instruction to cause a branch more than 127 bytes forward or 128 bytes backward from the current value of the program counter. (A branch backward is indicated if the relative address specified is negative; forward if it's positive. A byte is negative if bit 7 is set. A byte is positive if bit 7 is clear. Thus, a value of 00 is considered positive.) However, an instruction called JMP allows the programmer to specify an unconditional branch to any location in memory. Therefore, if we have a short conditional branch followed by an unconditional jump, we may achieve in two instructions a conditional branch to any location in memory.

## Unconditional Branch

Just as BASIC has its GOTO command, which causes an unconditional branch to a specified line in a BASIC program, the 6502 has its JMP instruction, which un-
conditionally branches to a specified address. A program may loop forever by JMP'ing back to its starting point.

Look at Example 11. Unless a line of code within the loop causes the 6502 to branch to a location outside of the loop, the 6502 will sit in this loop forever.

## Example II

Endless Loop:

| START |  | some |
| :---: | :---: | :---: |
|  | xxxxxxxxx | instructions |
|  | x $x \times x \times x \times x \times x$ |  |
|  | JMP START |  |

## Indirect Addressing

A JMP instruction may be written in either the absolute addressing mode or the indirect addressing mode. Absolute addressing is used in Example 11. The operand is the address to which the 6502 should jump. But in the indirect mode (which is always signified by parentheses in the operand field) the operand specifies the address of a pointer. The 6502 will jump to the address specified by the pointer; it will not jump to the pointer itself.

The line of code "JMP (POINTR)" will cause the 6502 to jump to the address specified by the 2 bytes at POINTR and POINTR +1 . Thus, if POINTR $=\$ 0600$, and the 6502 executes the instruction "JMP (POINTR)" when memory location $\$ 0600$ holds $\$ 00$ and $\$ 0601$ holds $\$ 20$, then the 6502 will jump to address $\$ 2000$. (Remember, addresses are always stored in memory with the low byte first.)

## How Branching Works

Incidentally, all branches, whether relative, absolute, or indirect, work by operating on the contents of the PC (program counter). Before any branch instruction is executed, the PC holds the address of the current opcode. A branch instruction changes the PC, so that in the next instruction cycle the 6502 will fetch not the opcode following the current opcode, but the opcode at the location specified by the branch instruction. Then execution will continue normally from the new address.

## Relocatability

Often I implement short unconditional branches as:
CLC
BCC PLACE
rather than as:
JMP PLACE

This is because the first method (relying as it does on relative rather than absolute addressing) will still work even if you relocate the code in which it is contained. Making your code relocatable will save you time and trouble when you try to move your programs around in memory and still want them to work.

To relocate code containing the second example, you'd have to change the operand field because the absolute address of PLACE will have changed. To relocate code containing the first example, you wouldn't have to change a thing.

## Subroutines

Perhaps the two most powerful instructions available to the assembly-language programmer are the JSR (Jump to SubRoutine) and the RTS (ReTurn from Subroutine). These instructions (equivalent to GOSUB and RETURN in BASIC) enable us to organize chunks of code as building blocks called subroutines.

Think of the subroutine as a job. Your computer can do more work for you if it knows how to do more jobs. Once you teach the 6502 how to do a given job, you won't have to tell it twice. Let's say you're writing a program in which the same operation must be performed at various times within a program. In every location within your program where the operation is required, you could include code to perform that operation. On the other hand, you could write code in one place to perform that operation, but write that code as a subroutine, and then call that subroutine whenever necessary from the main, or calling program. A call to a subroutine causes that routine to execute. When finished, it returns to the instruction following the call in the main program.

It only takes one line of code to call a subroutine. JSR SUB will call the subroutine located at the address labeled SUB. After the 6502 fetches and executes the JSR opcode, the next opcode it fetches will be at the address labeled SUB, in this example. So far it looks like an unconditional JMP. The 6502 will fetch and execute opcodes from the addresses following SUB, until it encounters an RTS instruction.

When the 6502 fetches an RTS instruction, it returns to its caller, jumping to the first opcode following the JSR instruction that called the subroutine. In effect, when a line of code calls a subroutine, the 6502 remembers where it is before it jumps to the new location. Then when it encounters an RTS instruction, it knows the address to which it should return because it remembers where it came from. It then continues to fetch opcodes from the point following the JSR instruction. Figure 3.1 illustrates this procedure. Note that the same subroutine may be called from many different points in the same program, and will always return to the opcode following the JSR instruction that called it.


Figure 3.1: Jump to and return from subroutine. When the processor encounters a JSR (jump to subroutine) instruction, the next instruction executed is the first instruction of the subroutine. Here, the subroutine SUB is called from MAIN. The last instruction executed in a subroutine must be an RTS (return from subroutine) instruction. Here, the instruction at label LAST in subroutine SUB returns control to the next instruction following the call to the subroutine in the main program, the instruction labeled NEXT. The subroutine SUB can be called anywhere in the program MAIN when the particular function of SUB is needed.

Subroutines allow you to structure your software. With structured software, you can make changes to many programs just by changing one subroutine. If, for example, all programs that print characters do so by calling a single-character-print subroutine, then any time you improve that subroutine you improve the printing behavior of all your programs. Changing something only once is a tremendous advantage over having to change something in many different (usually undocumented) places within a piece of code. For these reasons, all of the software in this book uses subroutines.

## Dummies

A dummy subroutine is a subroutine consisting of nothing but an RTS instruction. A line of code in a program can call a dummy subroutine and nothing will happen; the 6502 will return immediately, with its registers unchanged.

So why call a dummy subroutine?
A call to a dummy subroutine provides a "hook," which you may use later to call a functional subroutine. While developing a program, I may have many lines of code that call dummy subroutines. Later, when I write the lower-level subroutines, it's easy to change my program so that it calls the functional subroutines rather than the dummy subroutines. Trying to insert a subroutine call to a program lacking such a hook can make you wish for a "memory shoehorn," which might let you squeeze 3 extra bytes of code into the same address space.

## The Stack

In addition to the addressing modes that enable the 6502 to access addressable memory, one addressing mode lets the 6502 access a 256 -byte portion of memory called the stack.

You may think of this stack as a stack of trays in a cafeteria. The only way a tray can be added is to place it on top of the existing stack. Similarly, the only way to get a tray from the stack is to remove one from the top. This is the LIFO (Last-In, First-Out) method. The last tray placed onto the stack must be the first tray removed.

In our case, when an item is placed onto the top of the stack, it is called a push, and when an item is removed from the top of the stack, it is called a pop. The last item onto the stack is said to be at the top of the stack.

For example, let's say we want to place two items onto the stack. (Each item has an 8 -bit value, perhaps a number or an ASCII character; see figure 3.2a.) First we push item 1 onto the stack, as illustrated in figure 3.2b. All positions above item 1 on the stack are said to be empty, the item 1 is on the top of the stack.

Now, push item 2 onto the stack (see figure 3.2c). What happens? Item 2 is now at the top of the stack, not item 1, although item 1 is still on the stack.

Next, to get item 2 back off the stack, we do a pop (see figure 3.2d). This makes item 1 the top of the stack again. Finally, another pop will remove item 1 from the stack, leaving the stack completely empty. Note that we had to pop item 2 from the stack before we could get to item 1 again. This is the LIFO principle.

The instruction PHA lets you push the contents of the accumulator onto the stack. PLA lets you load the accumulator from the top of the stack (a pop). PHP lets you push the processor status register onto the stack. PLP lets you load the processor status register from the stack.


Figure 3.2: Pushing and popping the stack.

The stack is a very convenient "pocket" to use when you want to store one or a few bytes temporarily without using an absolute place in memory. Subroutines may pass information to the calling routines by using the stack, but be careful: if a subroutine pushes data onto the stack, and fails to pop that data from the stack before executing an RTS instruction, then that subroutine will not return to its caller. This happens because when the 6502 executes a JSR instruction, it pushes the return address-that is, the address of the opcode following the JSR instruc-tion-onto the stack. A subroutine can return to its caller only because its return address is on the stack. If its return address is not at the top of the stack when the subroutine executes an RTS, it will not return to its caller. So a subroutine should always restore the stack before trying to return.

## Chapter 4:

## Arithmetic and Logic

## Character Translation

As demonstrated by Examples 7 and 8, indexed addressing is handy for performing a given operation (such as a move) on a contiguous group of bytes. But it also has another important application: table lookup. For example, let's say you and a friend have decided to write notes to one another using a substitution code. For every letter, number, and punctuation mark in a message, you've agreed to substitute a different character. A "W" will be replaced with a "Y;" a semicolon may be replaced with a " 9 ," etc.

You each have the same table showing you what to substitute for each character that may appear in a message. So you write a note to your friend in English, and then, using this table (which might be in the form of a Secret Agent Decoding Ring) you code, or encrypt, your note. You send the note in its encrypted form to your friend. Anyone else looking at the note would just see garbage, but your friend knows that a message can be found in it. So he gets his copy of the character translation table (which may be in his Secret Agent Decoding Ring), and he translates the encrypted message back into English, looking up the characters that correspond to each character in the coded message.

Children often enjoy coding and decoding messages in this way, but I find it about as much fun as filling out forms - which is no fun at all. Unfortunately, programming often involves character translation. Fortunately, I don't have to do it myself. I let my computer perform any necessary character translation by having it do what our two secret agents were doing: look up answers in a table.

## Example 12 <br> Character Translation Subroutine

XLATE TAX
LDA TABLE, $X$ RTS

Use character to be translated as an index into the table. Look up value in table.
Return to caller, bearing translated character in A and original character in X.

## Transfer Register

In Example 12, the subroutine XLATE assumes when it is called that the accumulator holds the byte to be translated. This byte might be a letter, a number, a punctuation mark, a control code, or a graphic character, but however you think of it, it's an 8 -bit value. Line 1 of XLATE transfers that 8 -bit value from the accumulator to the $X$ register, using the register-transfer instruction TAX.

Register-transfer instructions operate only on registers; they do not affect addressable memory. These instructions allow the contents of one register to be copied, or transferred, to another. The results of a transfer leave the source register unchanged, and the destination register holding the same value as the source register. The 6502's register-transfer instructions are:

TAX Transfer accumulator to $X$ register. TAY Transfer accumulator to $Y$ register. TXA Transfer $X$ register to accumulator. TYA Transfer $Y$ register to accumulator.

Register transfers do not affect the status flags.
These instructions let you transfer $A$ to $X$ or $Y$, or to transfer $X$ or $Y$ to $A$. But how would you transfer $X$ to $Y$, or $Y$ to $X$ ? (Hint: it will take two lines of source code, each line an instruction from the list above.)

## Table Lookup

In Example 12, line 2 of XLATE actually performs the character translation by looking up the desired data in a table. The label, TABLE, identifies the base address for a table that we've previously entered into memory. The indexed addressing
mode allows line 2 to get the $X$ th byte above the base address (ie: to get the $X$ th byte of the table). When that line is executed, the table lookup is complete. The 6502 has looked up and now holds in the accumulator the Xth byte in the table. Now all the 6502 must do is return to its caller, bearing the translated character in A and the original character in X. It accomplishes this with the RTS instruction.

Now you can perform this character translation at any point in any program with just one line of source code:

## JSR XLATE

Table lookup gives me great flexibility as a programmer. If a program uses a table lookup and for some reason I want the program to behave differently, I will probably only have to change some values in the table; it's unlikely that I'll have to change the table lookup code itself. If I've set up my table well, I might not have to change anything in the program except the data in the table.

Table lookup is therefore a very fast and flexible means of performing data translation. But the cost of that speed and flexibility can be size. You might be able to solve any problem with the right tables in memory, but not if you can't afford the memory necessary to hold all those tables. It's great when a program can just look up the answers it needs, but sometimes a program will actually have to compute its answers.

## Arithmetic Operations

The 6502 can perform the following 8 -bit arithmetical operations:

Shift<br>Rotate<br>Increment<br>Decrement<br>Add<br>Subtract

To understand how the 6502 operates on a byte, you must think of the bits in that byte. Even if the byte represents a number or a letter, don't think about what you can do to that number or letter. Think about what you can do to the pattern of bits in that byte.

What can you do to those bits?

Shift
You can shift the bits in a byte one position to the left or to the right. An ASL (Arithmetic Shift Left) operates on a byte in this manner: it moves each bit one bit to the left; it moves the leftmost bit (bit 7) into the carry flag, and it sets the rightmost bit (bit 0 ) to zero. See figure 4.1.


Figure 4.1: Effect of the ASL instruction.

For example, if the byte at location TMP has the following bit pattern:
$\begin{array}{lllllllll}\text { address TMP } & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 0\end{array}$
then after the instruction "ASL TMP" is executed, the data would look like:
$\begin{array}{lllllllll}\text { address TMP } & 1 & 0 & 1 & 0 & 1 & 1 & 0 & 0\end{array}$
with the carry flag being set to the previous value of bit 7 , in this case 0 . If the same instruction is again executed, the data becomes:

| address TMP | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

and the carry flag is set to 1.
A LSR (Logical Shift Right) has just the opposite effect of the ASL. All bits are shifted to the right towards the carry flag, introducing zeroes through bit 7 . See figure 4.2.


Figure 4.2: Effect of the LSR instruction.

For example, if the byte at location TMP is as originally given above, then after the instruction "LSR TMP" is executed, the data at TMP becomes:
$\begin{array}{lllllllll}\text { address TMP } & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1\end{array}$
with the carry flag being set to the previous value of bit 0 , in this case zero. If the same instruction is executed again, the data becomes:
$\begin{array}{lllllllll}\text { address TMP } & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1\end{array}$
with the carry flag set to 1 .
Because a number is represented in binary (each bit represents a successive power of two), some arithmetic operations are simple. To divide a byte by two, simply shift it right; to multiply a value in a byte by two, simply shift it left.

## Rotate

You can also rotate the bits in a byte to the left or to the right through the carry flag. Unlike shifting, rotating a byte preserves all the information originally contained by a byte.

Figure 4.3 shows how a ROL (rotate left) instruction works. For instance, let's say the data at address TMP is originally the same as in previous examples:

$$
\begin{array}{lllllllll}
\text { address TMP } & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 0
\end{array}
$$

and let's say that the carry flag is set (ie: it holds a 1).
After a "ROL TMP" instruction is executed, the data becomes:
$\begin{array}{lllllllll}\text { address TMP } & 1 & 0 & 1 & 0 & 1 & 1 & 0 & 1\end{array}$


Figure 4.3: Effect of the ROL instruction.
and the carry bit is set to the previous value of bit 7, namely 1. Notice that bit 0 in TMP now holds the original contents of the carry flag, and the carry flag holds the original contents of bit 7. Otherwise, everything looks just the same as in the ASL operation. After a second execution of the instruction "ROL TMP," the data becomes:
$\begin{array}{lllllllll}\text { address TMP } & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 1\end{array}$
with the carry flag set to 1 .
In a rotate left instruction, bit 0 is always set from the carry flag. (In the ASL instruction, bit 0 is always set to 0 .) If this had been an ASL instruction, what would the bit pattern at TMP be?

Figure 4.4 shows how a ROR (rotate right) instruction works. It is similar to ROL, except that the carry flag is set from bit 0 , and bit 7 is set from the carry flag.


Figure 4.4: Effect of the ROR instruction.

Rotate a byte left nine times and you'll still have the original byte. The same is true if you rotate a byte right nine times. But shift a byte left nine times, or right nine times, and you know what you've got left? Nothing!

## Increment, Decrement

You can increment or decrement a byte in three ways: using the INC and DEC instructions to operate on a byte in memory, using INX and DEX to operate on the $X$ register, or using INY and DEY to operate on the Y register. None of these instructions affects the carry flag. They do affect the zero flag: $Z$ is set if the result of an increment or decrement is zero; otherwise $Z$ is cleared. The negative flag is set if the result of an increment or decrement is a byte with bit 7 set; otherwise N is cleared.

Note that if you increment a register or address holding $\$ \mathrm{FF}$, it will hold zero. And similarly, if you decrement a register or address holding a zero, it will hold \$FF.

You cannot increment or decrement the accumulator, but you can add or subtract a byte from the accumulator.

## Addition

Example 13 shows how to add a byte from the location labeled NUMBER to the accumulator:

## Example 13

CLC
ADC NUMBER

Clear the carry flag.
Add the contents of location
NUMBER to the accumulator.

After these instructions are executed, the accumulator will hold the low 8 bits of the result of the addition. If, following the addition, the carry flag is set, then the result of the addition was greater than 255 ; if the carry flag is clear, then the result was less than 256 , and, therefore, the accumulator is holding the full value of the result. Remember, the carry flag must be cleared before performing the ADC instruction.

## Subtraction

Subtraction is as easy as addition. To subtract a byte from the accumulator, first set the carry flag (using the SEC instruction) and then subtract from the accumulator a constant or the contents of some address, using the instruction SBC (subtract with carry):

SEC
Set the carry flag. SBC OPERND

Subtract from accumulator the value of OPERND.

If the operand is greater than the initial value of the accumulator, the subtract operation will clear the carry flag; otherwise the carry flag will remain set. In either case, the accumulator will bear the 8 -bit result.

Thus, you clear the carry flag before adding and set the carry flag before sub-
tracting. If the carry flag doesn't change state, then the accumulator bears the entire result. But if the addition or subtraction changes the state of the carry flag, then your result is greater then 255 (for an addition) or less than zero (for a subtraction).

## Decimal Mode

The processor status register includes a bit called the decimal flag. If the decimal flag is set, then the 6502 will perform addition and subtraction in decimal mode. If the decimal flag is clear, then the 6502 will perform addition and subtraction in binary mode. Decimal mode means the bytes are treated as BCD (Binary Coded Decimal), meaning that the low 4 bits of a byte represent a value of 0 thru 9 , and the high 4 bits of the byte represent a value of 0 thru 9 . Neither nybble ( 4 bits) may contain a value of A-F. So, each nybble represents a decimal digit.

The instructions SED and CLD set the decimal flag and clear it, respectively. Unless you'll be operating with figures that represent dollars and cents, you won't need to use the decimal mode. All software in this book assumes that the decimal mode is not used.

Decimal 255 is the biggest value that can be represented by a binary-coded byte, but decimal 99 is the biggest value that can be represented by a byte using Binary Coded Decimal.

## Logical Operations

What if you want to set, clear, or change the state of one or more bits in a byte without affecting the other bits in that byte? Input and output operations often demand such "bit-twiddling," which can be performed by the 6502's logical operations ORA, AND, and XOR.

## Setting Bits

The ORA instruction lets you set one or more bits in the accumulator without affecting the state of the other bits. ORA logically OR's the accumulator with a specified byte, or mask, setting bit $n$ in the accumulator if bit $n$ in the accumulator is initially set or if bit n in the mask is set, or if both of these bits are set. A logical OR will leave bit n of the accumulator clear only if bit n is initially clear in both the accumulator and the mask. Table 4.1 shows a truth table for the logical operator OR. A truth table gives all possible combinations of 2 bits that can be operated upon (in this case, ORed) and the results of these combinations.

Table 4.1: Truth table for the logical OR operand.
Bit 1 Bit $2 \quad$ Result

| 0 | OR 0 | $=0$ |
| :--- | :--- | :--- |
| 0 | OR 1 | $=1$ |
| 1 | OR 0 | $=1$ |
| 1 | OR 1 | $=1$ |

For example, suppose we executed the instruction "ORA \#\$80." Here the mask is $\$ 80$, or the bit pattern 10000000 . This instruction would therefore set bit 7 of the accumulator while leaving all other bits unchanged. So, if the accumulator had a value of 00010010 before the above instruction was executed, it would have the value of 10010010 afterwards.

Another example would be "ORA \#3." Since a decimal 3 becomes 00000011 when converted to an 8 -bit binary mask, the above instruction would set bits 0 and 1 in the accumulator, leaving bits 2 thru 7 unchanged.

How would you set the high 4 bits in the accumulator? The low 4 bits?

## Clearing Bits

You can clear one or more bits in the accumulator without affecting the state of the other bits through the use of the AND instruction. AND performs a logical AND on the accumulator and the mask specified by the operand. AND will set bit $n$ of the accumulator only if bit $n$ of the accumulator is set initially and bit $n$ is set in the mask. If bit n is initially clear in the accumulator or if bit n is clear in the mask, then AND will clear bit n in the accumulator. Table 4.2 gives the truth table for the logical AND operation.

Table 4.2: The truth table for the logical AND.
Bit 1 Bit 2 Result

| 0 | AND | 0 |
| :--- | :--- | :--- |
| 0 | $=0$ |  |
| 1 | AND 1 | $=0$ |
| 1 | AND 1 | $=$ |

For instance, the line of source code "AND \#1" will clear all bits except bit 0 in the accumulator; bit 0 will remain unchanged. "AND \# $\$ \mathrm{FO}$ " will clear the low 4 bits of the accumulator, leaving the high 4 bits unchanged. Select the right mask, and you can clear any bit or combination of bits in the accumulator without affecting the other bits in the accumulator.

## Toggle Bits

The exclusive OR operation, XOR, lets you "flip," or toggle, one or more bits in the accumulator (ie: change the state of one or more bits without affecting the state of other bits). XOR will set bit $n$ of the accumulator if bit $n$ is set in the accumulator but not in the mask, or if bit n is set in the mask but not in the accumulator. If bit n has the same state in both the accumulator and in the mask, then XOR will clear bit n in the accumulator. Table 4.3 shows the truth table for this operation.

Table 4.3: The truth table for the exclusive $O R$ (XOR).
Bit 1 Bit 2 Result

| 0 | XOR | 0 | $=$ | 0 |
| :--- | :--- | :--- | :--- | :--- |
| 0 | XOR | 1 | $=$ | 1 |
| 1 | XOR | 0 | $=$ | 1 |
| 1 | XOR | 1 | $=$ | 0 |

To toggle bit n in the accumulator, simply XOR the accumulator with a mask which has bit $n$ set but all other bits clear. Bit $n$ will change state in the accumulator, but all other bits in the accumulator will remain unchanged.

The logical operators, combined with the 6502's relative branch instructions, make it possible for a program to take one action or another depending on the state of a given bit in memory. Let's say you want a piece of code that will take one action (Action A) if a byte, called FLAG, has bit 6 set; yet take another action (Action B) if that bit is clear. The code of Example 14 shows one way to ignore all other bits in FLAG, and still preserve FLAG.

## Example 14

| LDA FLAG | Get flag byte. |
| :--- | :--- |
| AND \#\$40 | Clear all bits but bit 6. |

PLAN.B

Take Action A, since bit 6 was set in flag.

Take Action B, since bit 6 was clear in flag.

What good are flags? Let me give an example. The flag on a rural mailbox may be either raised or lowered to indicate that mail is or is not awaiting pickup. Raising and lowering those flags requires a little bit of effort (no pun intended), but it enables the mail carrier to complete the route much more quickly than would be possible if every mailbox had to be checked every time around. Presumably, this provides better service for everyone on the route.

That mail carrier's routine is a very sophisticated piece of programming. If we think of the mail carrier as a person following a program, then we can see some of the power and flexibility that come from the use of flags.

The mail carrier's program has two parts: What must be done at the post office and What must be done on the route. At the post office, the mail carrier sorts the mail, bundles letters for the same address and puts the bundles for a given route into a mail sack in some order. This sorting at the post office means the mail carrier on the route can make his or her rounds more quickly, because no further sorting and searching is required. (We won't go into sorting and searching in this book; that's a volume in itself. For a helpful reference see Donald E Knuth's Searching and Sorting.)

Now comes the second part of the mail carrier's program: What must be done on the route. The mail carrier picks up the mail sack and leaves the post office. Driving down country roads, the mail carrier sees a mailbox ahead. Do I have any mail for the people at this address? If so, the mail carrier's mental program says, I'll slow down and deliver it. But what if I don't have any mail now for these people? Do I just keep driving? Do I go to the next address?

Not if I want to keep my job.
The mail carrier looks a little more closely at the mailbox. Is the flag up or down? If it's down, I can just drive by, but if the flag is up I must stop and pick up the outgoing mail.

A flag is just a single bit of information, but by interpreting and responding to the state of flags, even a simple program can respond to many changing conditions. If your computer has 8,000 bytes of programmable memory, that means it has 64,000 bits of memory. Conceivably, you could use most of those bits as flags, perhaps simulating the patterns of outgoing mail in a community of more than 50,000 households.

But you didn't buy a computer to play post office. And you know enough now to follow the programs presented in the following chapters. These programs will in-
clude examples of all the instructions and programming techniques presented in this very fast course in assembly-language programming. The programs in the following chapters will also give you some tools to use in developing your own programs.
(Incidentally, there is one 6502 instruction which doesn't do anything at all. The instruction NOP performs NO operation. Why would you want to perform no operation? Occasionally, it's handy to replace an unwanted instruction with a dummy instruction. When you want to disable some code, simply replace the unwanted code with NOP's. A NOP is represented in memory by \$EA.)

## Chapter 5:

## Screen Utilities

Now let's consider how to display something on the video screen. On the Apple, Atari, OSI, and PET computers, the video-display circuitry scans a particular bank of memory, called the display memory. Every address in the display memory represents, or is mapped to, a different screen location (hence the term memorymapped display). For each character in the display memory, the display circuitry puts a particular image, or graphic, on the screen (hence the term character graphics). To display a character in a given screen location, you need only store that character in the one address within display memory that corresponds to the desired screen location.

To know which address corresponds to a given screen location you must consult a display-memory map. Appendices B1 thru B4 describe how display memory is mapped on the Apple, Atari, OSI, and PET computers. Note that two different systems may have two different addresses for the same screen location. Also note how burdensome it can be to look up the addresses of even a few screen locations just to display a few characters on the video screen.

Rather than address the screen in an absolute manner, we'd like to be able to do so indirectly. Ideally, we'd like a software-controlled "hand" that we can move about the screen. Then we could pick up the character under the hand, or place a new character under the hand, without being concerned with the absolute address of the screen location under the hand at the moment. Such a hand can be implemented quite easily as a zero-page pointer.

## Pointers

A pointer is just a pair of contiguous bytes in memory. Since 1 byte contains 8 bits, a pointer contains 16 bits, which means a pointer can specify any one of more than 65,000 (specifically: $2^{16}$ ) different addresses.

A pointer can specify, or point to, only one address at a time. The low byte of a pointer contains the 8 LSB (least-significant bits) of the address it specifies, and the high byte of the pointer contains the 8 MSB (most-significant bits) of the address it specifies.

Let's say we want a pointer at location $\$ 1000$. We must allocate 2 bytes for the pointer, which means it will occupy the bytes at $\$ 1000$ and $\$ 1001$. $\$ 1000$ will hold the low byte, and $\$ 1001$ will hold the high byte. If we want this pointer to specify address $\$ A B C D$, then we may set it as follows:

| POINTR = \$1000 |  | This assembler directive equates the label <br> POINTR with the value $\$ 1000$. (It's POINTR <br> and not POINTER only because the assembler <br> used in preparing this book chokes on labels |
| :--- | :--- | :--- | :--- |
| longer than six characters - a common, if |  |  |
| arbitrary, limitation.) |  |  |

Now POINTR points to \$ABCD.
Although a pointer may be anywhere in memory, it becomes especially powerful when it's in the zero page (the address space from 0000 to $\$ 00 \mathrm{FF}$ ). The 6502 's indirect addressing modes allow a zero-page pointer to specify the address on which certain operations may be performed. A zero-page pointer must be located in the zero page, but it may point to any location in memory. For example, a zero-page pointer may be used to specify the address in which data will be loaded or stored. Since display memory looks like any other random-access memory to the processor, we may implement our television hand as a zero-page pointer.

## TV.PTR

We want a zero-page pointer that can point to particular screen locations. Let's call it TV.POINTER, or TV.PTR for short. Whenever we examine or modify the screen, we'll do it through the TV.PTR.

Because the TV.PTR must be in the zero page, let's place it at $\$ 0000$, meaning it will occupy the bytes at $\$ 0000$ and at $\$ 0001$. We can do that with the following assembler directive:
TV.PTR = \$0

## TV.PUT

The TV.PTR always specifies the current location on the screen. Thus, to display a graphic at the current location on the screen, we need only load the accumulator with the 8 -bit code for that graphic and then execute the following two lines of code:

| LDY \#0 | A0 | 00 |
| :--- | :--- | :--- |
| STA (TV.PTR), Y | 91 | 00 |

The two lines of above code are sufficient to display a given graphic in the current screen location. But what if you want to display a given character in the current screen location? The ASCII code for a character is not necessarily the same as your system's display code for that character's graphic. To display an " A " in the current screen location, we cannot simply load the accumulator with an ASCII "A" (which is $\$ 41$ ) and then execute the two lines of above code, because the graphic " A " may have a different display code on your system. Instead of displaying an "A," we might display something else. Of the four computers considered in this book, only the Ohio Scientific Challenger I-P has a one-to-one correspondence between any character's ASCII code and that character's graphic code. The Atari, the PET, and the Apple computers lack such a one-to-one correspondence.

How then can we display a given ASCII character in the current screen location? We can do it by assuming that there exists a subroutine called FIXCHR, which will "fix" any given ASCII code, by translating it to its corresponding graphic or display code. FIXCHR will be different for each system, so we won't go into its details here (see the appendix pertaining to your computer for a description and listing of FIXCHR for your system). At this point we will assume only that FIXCHR exists, and that if we call it with an ASCII character in the accumulator, it will return with the corresponding display code in the accumulator.

We already know how to display a given graphic in the current screen location. With FIXCHR we now know how to display any given ASCII character in the current screen location. And since displaying any given ASCII character in the current screen location is something we're likely to do more than once, let's make it a subroutine. We'll call that subroutine TV.PUT since it will let us put a given ASCII
character up on the TV screen:

TV.PUT JSR FIXCHR
LDY \#O
STA (TV.PTR), Y RTS

Convert ASCII character to your system's display code for that character. Put that graphic in the current screen location. Return to caller.

## The Screen Location

However, these examples of modifying and examining screen locations through the TV.PTR will work only if the TV.PTR is actually pointing at a screen location. Therefore, before executing code such as the examples given above, we must be sure the TV.PTR points to a screen location.

There are several ways to do this. If you want to write code that will run on only one machine (or on several machines whose display memory is mapped the same way), then you can use the immediate mode to set the TV.PTR to a given address on the screen. Let's say you want to set the TV.PTR to point to the third column of the fourth row (counting right and down from an origin in the upper-left corner). If you have an Ohio Scientific Challenger I-P, then you can consult your system's documentation and determine that address \$D062 in display memory corresponds to your desired screen location. \$D0 is the high byte of this screen location; $\$ 62$ is the low byte of this screen location. Thus, you can set TV.PTR with the following lines of code:

| LDA \#\$62 | A9 | 62 | Set |
| :--- | :--- | :--- | :--- |
| STA TV.PTR | 85 | 00 | low byte. |
| LDA \#SDO | A9 | D0 | Set |
| STA TV.PTR +1 | 85 | 01 | high byte. |
| $\cdot$ |  |  |  |
| $\cdot$ |  |  |  |
| $\cdot$ |  |  |  |
| $\cdot$ |  |  |  |

This code is fast and relocatable. But it's not very convenient to have to look up a display address every time we write code that displays something on the screen. It
would be much more convenient if we could address the screen as a series of $X$ and $Y$ coordinates. Why not have a subroutine that sets the TV.PTR for us, provided we supply it with the desired X and Y coordinates?

## TVTOXY

TVTOXY is a subroutine that sets the value of the TV.PTR to the display address whose $X$ and $Y$ coordinates are given by the $X$ and $Y$ registers. (Note that we count the columns and rows from zero.) To make the TV.PTR point to the third column from the left in the fifth row from the top, a calling program need only include the following code:

LDY \#2 The leftmost column is column zero, so the third column is column two.
LDY \#4 The topmost row is row zero, so the fifth row is row four.
JSR TVTOXY
Set TV.PTR to screen location whose $X$ and $Y$ coordinates are given by the $X$ and $Y$ registers.

How will TVTOXY work? We could have TVTOXY do just what we were doing: look up the desired address in a table. A computer can look up data in a table very quickly, but the speed may not be worth it if the table requires a lot of memory. If we don't mind waiting a little longer for TVTOXY to do its job, we can have TVTOXY calculate the desired value of TV.PTR, rather than look it up in a table. But how can you calculate the address of a given $X$ and $Y$ location on the screen?

You can't do it without data. But you don't need a large amount of data to determine the address of a given $X, Y$ location in screen memory; you need only have access to the following facts:

HOME The address of the character in the upper-left corner of the screen (ie: the lowest address in screen memory).
ROWINC ROW INCrement: the address difference from one row to the next.

Knowing the values of HOME and ROWINC for a given system, you can calculate the address corresponding to any $X, Y$ location:

| HOME <br> +X Register <br> $+(Y$ Register $) \times$ ROWINC |
| :--- |
| TV.PTR | | Address of character in upper-left corner <br> +X coordinate <br> $+($ Y coordinate $) \times$ ROWINC |
| :--- |
| Address of screen location at column $X$, row $Y$. |

Run through this calculation for several screen locations and compare the results with the addresses you look up in the display-memory map for your system. (Remember that we count columns and rows from zero, not from one.) Now if TVTOXY can run through this calculation for us, we'll never have to look at a display-memory map again; we can write all our display code in terms of cartesian coordinates.

But we shouldn't be satisfied with TVTOXY if it only runs through the above calculation. After all, what happens if TVTOXY is called and the Y register holds a very large number? If the $Y$ register is greater than the number of rows on the screen, then the above calculation will set the TV.PTR to an address outside of display memory. We don't want that. Maybe a calling program will have a bug and call TVTOXY with an illegal value in $X$ or in Y. If TVTOXY doesn't catch the error, the calling program may end up storing characters in memory that is not display. memory. It might end up over-writing part of itself, which would almost certainly invite long and arduous debugging.

I hate debugging. I know I'm going to make mistakes, but I'd like my software to catch at least some bugs before they run amuck. So let's have TVTOXY check the legality of $X$ and $Y$ before blindly calculating the value of TV.PTR.

How can TVTOXY check the legality of $X$ and $Y$ ? How big can $X$ or $Y$ get before it's too big? We need some more data:

TVCOLS The number of columns on the display screen, counting from zero.
TVROWS The number of rows on the display screen, counting from zero.

Now TVTOXY requires the following four facts about the host computer:

HOME
ROWINC
TVROWS
TVCOLS

If we store these facts about the host system in a particular block of memory, then TVTOXY need only consult that block of memory to learn all it needs to know about the screen. TVTOXY can then work as follows:

## TVTOXY

| TVTOXY | SEC | Is $X$ out of range? |
| :---: | :---: | :---: |
|  | CPX TVCOLS |  |
|  | BCC X.OK | If not, leave it alone. |
|  |  | If $X$ is out of range, give |
|  | LDX TVCOLS | it its maximum legal value. Now $X$ is legal. |
| X.OK | SEC | Is $Y$ out of range? |
|  | CPY TVROWS |  |
|  | BCC Y.OK | If not, leave it alone. |
|  |  | If $Y$ is out of range, give |
|  | LDY TVROWS | it its maximum legal value. Now $Y$ is legal. |
| Y.OK | LDA HOME | Set TV.PTR $=$ HOME. |
|  | STA TV.PTR |  |
|  | LDA HOME +1 |  |
|  | STA TV.PTR + 1 |  |
|  | TXA | Add $X$ to TV.PTR. |
|  | CLC |  |
|  | ADC TV.PTR |  |
|  | BCC COLSET |  |
|  | INC TV.PTR+1 |  |
|  | CLC |  |


| COLSET | CPY \#O | Add Y*ROWINC to TV.PTR. |
| :--- | :--- | :--- |
| LOOP | BEQ EXIT |  |
|  | CLC |  |
|  | ADC ROWINC |  |
|  | BCC NEXT |  |

INC TV.PTR+1
NEXT
DEY
BNE LOOP
EXIT
STA TV.PTR
RTS
Return to caller.

## TVDOWN, TVSKIP, TVPLUS

Using TVTOXY, we can set TV.PTR to a screen location with any desired X,Y coordinates. But it would also be convenient to be able to modify TV.PTR relative to its current value. For example, after placing a character on the screen, we might want to make TV.PTR point to the next screen location to the right, or perhaps to the screen location directly below the current screen location. We might even want to make TV.PTR skip over several screen locations to make it point to "the nth screen location from here," where "here" is the current screen location. For these occasions, the subroutines TVDOWN, TVSKIP, and TVPLUS come in handy.

## TVDOWN, TVSKIP, TVPLUS

| TVDOWN | LDA ROWINC <br> CLC <br> BCC TVPLUS | Move TV.PTR down by one row. |
| :--- | :--- | :--- |
| TVSKIP | LDA \#1 | Unconditionally branch. <br> Skip one screen location by increment- |
| ing TV.PTR. |  |  |

RTS Return to caller.

Note that the routines TVDOWN and TVSKIP make use of the routine TVPLUS, which assumes that the accumulator has been set to the number of locations to be skipped. For TVDOWN and TVSKIP, the accumulator is set to ROWINC and 1, respectively.

Right now TVPLUS might not seem long enough to be worth making into a
subroutine. Any program that calls TVPLUS could perform the addition itself, at a cost of only a few bytes, and at a saving of several machine cycles in the process. However, we may make TVPLUS more sophisticated later on.

For example, we could enhance TVPLUS so it performs error checking automatically, to ensure that TV.PTR will never point to an address outside of screen memory. Such error checking would be very burdensome for every calling program to perform, but if and when we insert it into TVPLUS, every caller will automatically get the benefit of that modification.

## VUCHAR

With TV.PUT we can display an ASCII character in the current screen location, and with TVSKIP we can advance to the next screen location. So why not combine the two, creating a subroutine that displays in the current screen location the graphic for a given ASCII character, and then automatically advances TV.PTR so it points to the next screen location? This would make it easy for a calling program to display a string of characters in successive screen positions. Since this subroutine will let the user view a character, let's call it VUCHAR:

VUCHAR JSR TV.PUT

JSR TVSKIP RTS

Display, in the current screen location, the graphic for the character whose ASCII code is in the accumulator. Advance to the next screen location.

We could even squeeze VUCHAR into the code presented above for TVDOWN, TVSKIP, and TVPLUS, by inserting one new line of source code immediately above TVSKIP. (See Appendix C1, the assembler listing for the Screen Utilities, which also includes some error checking within TVPLUS.)

## VUBYTE

With the screen utilities presented thus far, we can display a character on the screen in the current location, but we don't have a utility to display a byte in hexadecimal representation. Let's make one.

We'll call this utility VUBYTE, since it will let the user view a given byte. With VUBYTE, a calling program must take only three steps to display a byte in hexadecimal representation anywhere on the screen:

1) Set a zero-page pointer (TV.PTR) to point to the screen location where the byte should be displayed; 2) load the accumulator with the byte to be displayed; and then 3) call VUBYTE.


Figure 5.1: Flowchart of the routine VUBYTE, which displays a byte in hexadecimal representation on the video screen.

VUBYTE will display the given byte as two ASCII characters in the current position on the screen, and when VUBYTE returns, TV.PTR will be pointing to the screen location immediately following the two screen locations occupied by the displayed characters.

VUBYTE need only determine the ASCII character for the hexadecimal value of the 4 MSB (most-significant bits), store that ASCII character in the screen location pointed to by TV.PTR, then display the ASCII character for the hexadecimal value of the accumulator's 4 LSB (least-significant bits) in the next screen location. See figure 5.1 for a flowchart outlining this.

VUBYTE seems to be asking for a utility subroutine to return the ASCII character for a given 4 -bit value. Let's call this subroutine ASCII. ASCII will return the ASCII character for the hexadecimal value represented by the 4 least-significant bits in the accumulator. It will ignore the 4 most-significant bits in the accumulator.

If we assume that ASCII exists, then we can write VUBYTE:

## VUBYTE <br> PHA <br> LSR A <br> LSR A <br> LSR A <br> LSR A <br> JSR ASCII

JSR VUCHAR

PLA
JSR ASCII

JSR VUCHAR

RTS

## VUBYTE

Save accumulator.
Move 4 MSB
into positions occupied by 4 LSB.

Determine ASCII for accumulator's 4 LSB (which were its 4 MSB ).

Display the ASCII character in the current screen location and advance to next. screen location.

Restore original value of accumulator.
Determine ASCII for accumulator's 4 LSB (which were its 4 LSB).

Display this ASCII character just to the right of the other ASCII character and advance to next screen location.

Return to caller.

Of course, ASCII doesn't exist yet. So let's write it, and then VUBYTE should be complete.

| ASCII |  |  |
| :---: | :---: | :---: |
| ASCII | AND \#\$0F | Clear the 4 MSB in accumulator. |
|  | CMP \#SOA BMI DECIML | Is accumulator greater than 9 ? |
|  | ADC \#6 | If so, it must be A thru F. Add $\$ 36$ to accumulator to convert it to corresponding ASCII character. (We'll add $\$ 36$ by adding $\$ 6$ and then adding $\$ 30$.) |
| DECIML | ADC \#\$30 | If accumulator is 0 thru 9 , add $\$ 30$ to it to convert it to corresponding ASCII character. |
|  | RTS | Return to caller, bearing the ASCII character corresponding to the hexadecimal value initially in the 4 LSB of the accumulator. |

## TVHOME, CENTER

Now we can display a character or a byte at the current screen location, and we can set the current screen location to any given $X, Y$ coordinates or modify it relative to its current value. It would also be handy if we could set the TV.PTR to certain fixed locations: locations that more than one calling program might need as points or origin. For example, a calling program might need to set the TV.PTR to the HOME location (position 0,0), or to the CENTER of the screen:

## TVHOME, CENTER

$\begin{array}{ll}\text { TVHOME } & \text { LDX \#O } \\ & \text { LDY \#O } \\ & \text { JSR TVTOXY }\end{array}$
RTS

Set TV.PTR to the leftmost column of the top row of the screen.

Then return to caller.
\(\left.$$
\begin{array}{ll}\text { CENTER } & \begin{array}{l}\text { LDA TVROWS } \\
\text { LSR A } \\
\text { TAY }\end{array} \\
\begin{array}{ll}\text { LDA TVCOLS } \\
\text { LSR A } \\
\text { TAX }\end{array} & \begin{array}{l}\text { Load A with total rows. } \\
\text { Divide it by two. } \\
\text { Y now holds the number of the central } \\
\text { row on the screen. }\end{array}
$$ <br>
Load A with total columns. <br>
Divide it by two. <br>
X now holds the number of the central <br>

column on the screen.\end{array}\right\}\)| Now X and Y registers hold X, Y coor- |
| :--- |
| dinates of center of screen. |

## TVPUSH, TV.POP

The screen utilities presented thus far enable us to set or modify the current position on the screen. We might also want to save the current position on the screen and then restore that position later. We can do this by pushing TV.PTR onto the stack and then pulling it from the stack:

## TVPUSH

## TVPUSH

| PLA | Pull return address from stack. |
| :---: | :---: |
| TAX | Save it in X... |
| PLA |  |
| TAY | ...and in Y. |
| LDA TV.PTR +1 | Get TV.PTR |
| PHA | and save |
| LDA TV.PTR | it on |
| PHA | the stack. |
| TYA | Place return |
| PHA | address back... |
| TXA |  |
| PHA | ... on stack. |
| RTS | Then return to caller. |

## TVPOP

| TV.POP | PLA | Pull return address from stack. |
| :---: | :---: | :---: |
|  | TAX | Save it in X... |
|  | PLA |  |
|  | TAY | ...and in Y . |
|  | PLA | Restore... |
|  | STA TV.PTR | ...TV.PTR |
|  | PLA | ...from |
|  | STA TV.PTR +1 | ...stack. |
|  | TYA | Place return |
|  | PHA | address back... |
|  | TXA |  |
|  | PHA | ... on stack. |
|  | RTS | Then return to caller. |

Now a calling program can save its current screen position with one line of source code: "JSR TVPUSH." That calling program can then modify TV.PTR and later restore it to its saved value with one line of source code: "JSR TV.POP."

## CLEAR SCREEN

Now that we can set TV.PTR to any $X, Y$ location on the screen, and display any byte or character in the current location, let's write some code to clear all or part of the screen. One subroutine, CLR.TV, will clear all of the video screen for us while preserving the zero page. A second routine, CLR.XY, will start from the current screen location and clear a rectangle, whose $X, Y$ dimensions are given by the $X, Y$ registers. Thus, a calling program can call CLR.TV to clear the whole screen; or a calling program can clear any rectangular portion of the screen, leaving the rest of the screen unchanged, just by making TV.PTR point to the upper left-hand corner of the rectangle to be cleared, and then calling CLR.XY with the X and Y registers holding, respectively, the width and height of the rectangle to be cleared.

CLR.TV JSR TVPUSH
JSR TVHOME

Save the zero-page bytes that will be changed.
Set the screen location to upper-left corner of the screen.
$\left.\begin{array}{lll} & \begin{array}{l}\text { LDX TVCOLS } \\ \text { LDY TVROWS } \\ \text { JSR CLR.XY }\end{array} & \begin{array}{l}\text { Load X,Y registers with } \\ \text { X,Y } \\ \text { Climensions of the screen. } \\ \text { screen location. }\end{array} \\ \text { JSR TV.POP } \\ \text { Restore zero-page bytes that were } \\ \text { changed. } \\ \text { Return to caller, with screen clear and } \\ \text { with zero page preserved. }\end{array}\right\}$

There are many more screen utilities you could develop, but the utilities presented in this chapter are a good basic set. Now programs can call the following subroutines to perform the following functions:

| ASCII: | Return ASCII character for 4 LSB in A. |
| :--- | :--- |
| CENTER: | Set current screen position to center of screen. |
| CLR.TV: | Clear the entire video display, preserving TV.PTR. |
| CLR.XY: | Clear a rectangle of the screen, with X,Y dimensions specified |
| by the X,Y registers. |  |
| TVDOWN: | Move current screen position down by one row. |

TVHOME: Set current screen position to the upper-left corner of the screen.
TVPLUS: Add A to TV.PTR.

TV.POP: TVPUSH:
TV.PUT: TVSKIP: TVTOXY:

VUBYTE:
VUCHAR:

Restore previously saved screen position from stack.
Save current screen location on stack.
Display ASCII character in A at current screen location. Advance to next screen location. Set current screen position to $X, Y$ coordinates given by $X, Y$ registers.

Display A, in hexadecimal form, at current screen location. Advance current screen location past the displayed byte. Display A as an ASCII character in current screen location; then advance to next screen location.

With these screen utilities, a calling program can drive the screen display without ever dealing directly with screen memory or even with the zero page. The calling program need not concern itself with anything other than the current position on the screen, which can be dealt with as a concept, rather than as a particular address hard-wired into the code.

## Chapter 6:

## The Visible Monitor

## Hand Assembling Object Code

An assembler is a wonderful software tool, but what if you don't have one? Is it possible to write 6502 code without an assembler?

You bet!
Not only is it possible to write machine code by hand, but all of the software in this book was originally assembled and entered into the computer by hand. In fact, I hand assembled my code long after I had purchased a cassette-based assembler, because I could hand assemble a small subroutine faster than I could load in the entire assembler.

Hand assembling code imposes a certain discipline on the programmer. Because branch addresses must be calculated by counting forward or backward in hexadecimal, I tried to keep my subroutines very small. (How far can you count backward in hexadecimal?) I wrote programs as many nested subroutines, which I could assemble and test individually, rather than as monolithic, in-line code. This is a good policy even for programmers who have access to an assembler, but it is essential for any programmer who must hand assemble code.

Yet once you've written a program consisting of machine-language instructions, how can you enter it into memory? You can read your program on paper, but how can you present it to the 6502?

A program called a machine-language monitor allows you to examine and modify memory. It also allows you to execute a program stored in memory. The Apple and Ohio Scientific computers each feature a machine-language monitor in ROM (read-only memory). The Atari computers feature a machine-language monitor in a plug-in program cartridge. Your system's documentation should tell you how to use the features of your monitor, but let's take a closer look at one
monitor in particular, the Ohio Scientific 65 V monitor. Because it is stored in readonly memory in the OSI Challenger I-P, I will refer to it as the OSI ROM monitor.

## A Minimal Machine-Language Monitor

You can invoke the OSI ROM monitor quite easily by pressing the BREAK key and then the " M " key. The monitor clears the video screen and presents the display shown in figure 6.1.


Figure 6.1: Ohio Scientific ROM (read-only memory) monitor display.

The display consists of two fields of hexadecimal characters: an address field and a data field. Figure 6.1 indicates that $\$ A 9$ is the current value of address $\$ 0000$.

The OSI ROM monitor has two modes: address mode and data mode. When the monitor is in address mode, you can display the contents of any address simply by typing the address on the keyboard. Each new hexadecimal character will roll into the address field from the right. To display address \$FEOD, you simply type the keys F, E, 0 , and then D.

To change the contents of an address, you must enter the data mode. When the

OSI ROM monitor is in the data mode, hexadecimal characters from the keyboard will roll into the data field on the screen. For your convenience, when the monitor is in the data mode you can step forward through memory (ie: increment the displayed address) by depresssing the RETURN key. Unfortunately, this convenience is not available in address mode, and neither mode allows you to step backward through memory (ie: to decrement the address field).

Beware: the OSI ROM monitor can mislead you. If the monitor is in the data mode and you type a hexadecimal character on the keyboard, that character will roll into the data field on the screen. Presumably that hexadecimal character also rolls into the memory location displayed on the screen. Yet, this might not be the case. In fact, the OSI ROM monitor displays the data you intended to store in an address, rather than the actual contents of that address. If you try to store data in a read-only memory address, for example, the OSI ROM monitor will confirm that you've stored the intended data in the displayed address, yet if you actually inspect that address (by entering address mode and typing in the address), you'll see that you changed nothing. This makes sense - you can't write to read-only memory. But the OSI ROM monitor leads you to think that you can.

The OSI ROM monitor can be confusing in other ways. For example, the display does not tell you whether you're in data mode or address mode; you've got to remember at all times which mode you last told the monitor to use. Furthermore, to escape from address mode you must use one key, while to escape from data mode you must use another key. Therefore you must always remember two escape codes as well as the current mode of the monitor.

Furthermore, the OSI ROM monitor does not make it very easy for you to enter ASCII data into memory. To enter an ASCII message into memory, you must consult an ASCII table (such as Appendix A2 in this book), look up the hexadecimal representation of each character in your message, and then enter each of those ASCII characters via two hexadecimal keystrokes. Then, once you've got an ASCII message in memory, the OSI ROM monitor won't let you read it as English text; you'll have to view that message as a series of bytes in hexadecimal format, and then look up, again in Appendix A2 or its equivalent, the ASCII characters defined by those bytes. That won't encourage you to include a lot of messages in your software - even though meaningful prompts and error messages can make your software much easier to maintain and use.

Finally, it is worth examining the way the OSI ROM monitor executes programs in memory. When you type " $G$ " on the Ohio Scientific Challenger I-P, the OSI ROM monitor executes a JMP (unconditional jump) to the displayed address. That transfers control to the code selected, but it does so in such a way that the code must end with another unconditional jump if control is to return to the OSI ROM monitor. This forces you to write programs that end with a JMP, rather than subroutines that end with an RTS.

Programs that end with a JMP are not used easily as building blocks for other programs, whereas subroutines are incorporated quite easily into software structures of ever-greater power. So wouldn't it be nice if a machine-language monitor
executed a JSR to the displayed address? This would call the displayed address as a subroutine, encouraging users to write software as subroutines, rather than as code that jumps from place to place. Such a monitor might actually encourage good programming habits, inviting the user to program in a structured manner, rather than daring the user to do so. In this chapter we'll develop such a monitor.

## Objectives

If you've spent any time using a minimal machine-language monitor, you've probably thought of some ways to improve it. Based on my own experience, I knew that I wanted a monitor to be:

## 1) Accurate

The data field should display the actual contents of the displayed address, not the intended contents of that address.

## 2) Convenient

It should be possible to step forward or backward through memory, in any mode. It should also be possible to enter ASCII characters into memory directly from the keyboard, without having to look up their hexadecimal representations first, and it should be possible to display such characters as ASCII characters, rather than as bytes presented as pairs of hexadecimal digits.

## 3) Encourage Structured Programming

The monitor should call the displayed address as a subroutine, rather than jump to the displayed address. This will encourage the user to write subroutines, rather than monolithic programs that jump from place to place.

## 4) Simplify Debugging

The monitor should load the 6502 registers with user-defined data before calling the displayed address. Thus a user can initially test a subroutine with different values in the registers. Then, when the called subroutine returns, the monitor should display the new contents of the 6502 registers. Thus, by seeing how it changes or preserves the values of the 6502 registers, the user could judge the performance of the subroutine.

Because my objective was to make the 6502 registers visible to the user by displaying the 6502 registers before and after any subroutine call, I've chosen to call this monitor the Visible Monitor. Figure 6.2 shows its display format.


Figure 6.2: Visible Monitor Display with fields numbered.

## VISIBLE MONITOR DISPLAY

## The Visible Monitor Display

Notice that the display in figure 6.2 has seven fields, not two as in the OSI ROM monitor display. The first two fields (fields 0 and 1) are the same as the two fields in the OSI ROM monitor - that is, they display an address and a hexadecimal representation of the contents of that address. Field 2 is a graphic representation of the contents of the displayed address. If that address holds an ASCII character, then the graphic will be the letter, number, or punctuation mark specified by the byte. Otherwise, that graphic will probably be a special graphic character from your computer's nonstandard (ie: nonASCII) character set.

Fields 3 thru 6 represent four of the 6502 registers: A (the Accumulator), X (the X Register), Y (the Y Register), and P (the Processor Status Register). When you type

G to execute a program, the 6502 registers will be loaded with the displayed values before the program is called; when control returns to the monitor, the contents of the 6502 registers at that time will be displayed on the screen.

In addition to the seven fields mentioned above, the Visible Monitor's display includes an arrow pointing up at one of the fields. In order to modify a field, you must make the arrow point to that field. To move the arrow from one field to another, I've chosen to use the GREATER THAN ( $>$ ) and LESS THAN ( $<$ ) keys. Touching the GREATER THAN key will move the arrow one field to the right, and depressing the LESS THAN key will move the arrow one field to the left. (If my computer had a cursor pad, I would use the cursor-left and the cursor-right keys to move the arrow from field to field, but it doesn't have a cursor pad, so GREATER THAN and LESS THAN have to fill the bill. You may assign the field-movement functions to any keys on your system, but GREATER THAN and LESS THAN are reasonable choices, because they look like arrows pointing right and left, respectively.)

I've chosen to use the space bar to step forward through memory and the return key to step backward through memory, but you may choose other keys if you prefer (eg: the " + " and " - " keys). The space bar seems reasonable to me for stepping forward through memory, because on a typewriter I press the space bar to bring the next character into view; RETURN seems reasonable for stepping backward through memory because RETURN is almost synonymous with "back up," and that's what I want it for: to back up through memory. With such a display and key functions, we ought to have a very handy monitor.

## Data

Before we develop the structure and code of the Visible Monitor, let's decide what variables and pointers it must have.

The Visible Monitor must have some way of knowing what address to display in field 0 . It can do this by maintaining a pointer to the currently selected address. Because it will specify the currently selected address, let's call this pointer SELECT. Then, when the user presses the spacebar, the Visible Monitor need only increment the SELECT pointer. When the user presses RETURN, the Visible Monitor need only decrement the SELECT pointer. That will enable the user to step forward and backward through memory.

The user will also want to modify the 6502 register images. Since there are four register images shown in figure 6.2 let's have 4 bytes, one for each register image. If we keep them in contiguous memory, we can refer to the block of register images as REGISTERS, or simply as REGS (since REGISTERS is longer than six characters, the maximum label length acceptable to the assembler used in the preparation of this book).

Finally, the Visible Monitor must keep track of the current field. Since there can
only be one current field at a time, we can have a variable called FIELD, whose value tells us the number of the current field. Then, when the user wants to select the next field, the Visible Monitor need only increment FIELD, and when the user wants to move the arrow to the previous field, the Visible Monitor need only decrement FIELD. If FIELD gets out of bounds (any value that is not 0 thru 6), then the Visible Monitor should assign an appropriate value to FIELD. The following code declares these variables in the form acceptable to an OSI 6500 Assembler:

| Variables |  |  |
| :---: | :---: | :---: |
| SELECT | .WORD 0 | This points to the currently selected byte. |
| REG.A | .BYTE 0 | REG.A holds the image of Register A (the Accumulator). |
| REG.X | .BYTE 0 | REG. X holds the image of Register $X$. |
| REG.Y | .BYTE 0 | REG.Y holds the image of Register Y. |
| REG.P | .BYTE 0 | REG.P holds the image of the Processor Status Register. |
| FIELD | .BYTE 0 | FIELD holds the number of the current field. |
| REGS $=$ REG.A |  |  |

## Structure

I want to keep the Visible Monitor highly modular, so it can be easily extended and modified. I have therefore chosen to develop the Visible Monitor according to the structure shown in figure 6.3. Clearly, the Visible Monitor loops. It places the monitor display on the screen. It then updates the information in that display by getting a keystroke from the user and performing an action based on that keystroke. It does this over and over.


Figure 6.3: A simple structure for interactive display programs.

With this flowchart as a guide, we can now write the source code for the top level of the Visible Monitor:

## VISMON

| VISMON | PHP | Save caller's status flags. |
| :--- | :--- | :--- |
| LOOP | JSR DSPLAY | Put monitor display on screen. |
|  | JSR UPDATE | Get user request and handle it. |
|  | CLC |  |
|  | BCC LOOP | Loop back to display... |

This is only the top level of the Visible Monitor; it won't work without two subroutines: DSPLAY and UPDATE. So it looks as if we've traded the task of writing one subroutine for the task of writing two. But by structuring the monitor in this way, we make the monitor much easier to develop, document, and debug.

Which subroutine should we write first? Let's start with the DSPLAY module, since the display is visible to the user, and the Visible Monitor must meet the user's needs. Once we know how to drive the display, we can write the UPDATE routine.

## Monitor Display

Figure 6.2 shows the display we want to present on the video screen. As you can see, this display consists of three lines of characters: the label line, the data line, and the arrow line. The label line labels four of the fields in the data line, using the characters A, X, Y, and P. The data line displays an address, the contents of that address (both in hexadecimal representation and in the form of a graphic), and then displays the values of the four registers in the 6502 . Underneath the data line, the arrow line provides one arrow pointing up at one of the fields in the data line.

Since the display is defined totally in terms of the label line, the data line, and the arrow line, we are ready now to diagram the top level of monitor display. See figure 6.4.

With the flowchart in figure 6.4 as a guide, we can now write source code for the top level of the DSPLAY subroutine:


Figure 6.4: Routine to display the monitor information.

## DSPLAY

| DSPLAY | JSR CLRMON | Clear monitor's portion of screen. |
| :--- | :--- | :--- |
|  | JSR LINE. 1 | Display the Label Line. |
|  | JSR LINE. 2 | Display the Data Line. |
|  | JSR LINE. 3 | Display the Arrow Line. |
|  | RTS | Return to caller. |

Now instead of one subroutine (DSPLAY), it looks as if we must write four subroutines: CLRMON, LINE.1, LINE.2, and LINE.3. But as the subroutines grow in number, they shrink in difficulty.

Before we put up any of the monitor's display, let's clear that portion of the screen used by the monitor's display. Then we can be sure we won't have any garbage cluttering up the monitor display.

Since we already have a utility to clear $X$ columns and $Y$ rows from the current location on the screen, CLRMON can just set TV.PTR to the upper-left corner of the screen, load $X$ and $Y$ with appropriate values, and then call CLR.XY. Here's source code:

LDX \#2
LDY \#2
JSR TVTOXY
LDX \#25 We'll clear 25 columns
LDY \#3
JSR CLR.XY
RTS

Set TV.PTR to column 2, row 2 of screen. and 3 rows. Here we clear them. Return to caller.

## Display Label Line

The subroutine LINE. 1 must put the label line onto the screen. We'll store the character string " $A X Y P$ " somewhere in memory, at a location we may refer to as LABELS. Then LINE. 1 need only copy 10 bytes from LABELS to the appropriate location on the screen. That will display the LABEL line for us:

## LINE.I

| LINE. 1 | LDX \#13 | X-coordinate of Label " A ". |
| :---: | :---: | :---: |
|  | LDY \#2 | Y-coordinate of Label " A ". |
|  | JSR TVTOXY | Place TV.PTR at coordinates given by $X, Y$ registers. |
|  | LDY \#0 | Put labels on the screen: |
|  | STY LBLCOL | Initialize label column counter. |
| LBLOOP | LDA LABELS, Y | Get a character and |
|  | JSR VUCHAR | put its graphic on the screen. |
|  | INC LBLCOL | Prepare for next character. |
|  | LDY LBLCOL | Use label column as an index. |
|  | CPY \#10 | Done last character? |
|  | BNE LBLOOP | If not, do next one. |
|  | RTS | Return to caller. |
| LABELS | . BYTE 'A X' | These are the characters |
|  | .BYTE 'Y P' | to be copied to the screen. |
| LBLCOL | .BYTE 0 | This is a counter. |

## Display Data Line

Displaying the data line will be more difficult than displaying the label line, for two reasons. First, the data to be displayed will change from time to time, whereas the labels in the label line need never change. Second, most fields in the data line dis-
play data in hexadecimal representation. To display 1 byte as two hexadecimal digits requires more work than is needed to display 1 byte as one ASCII character. However, we have a screen utility (VUBYTE) to do that work for us. In fact, we have enough screen utilities to make even the display of seven fields of data quite straightforward. Following, then, is the display data-line routine:

## LINE. 2

LINE. 2 LDX \#2
LDY \#3
JSR TVTOXY
LDA SELECT +1 JSR VUBYTE LDA SELECT JSR VUBYTE JSR TVSKIP JSR GET.SL

PHA JSR VUBYTE

JSR TVSKIP PLA JSR VUCHAR

JSR TVSKIP

LDX \#0
VUREGS LDA REGS, $X$ JSR VUBYTE JSR TVSKIP INX
CPX \#4
BNE VUREGS RTS

Load $X$ register with X -coordinate for start of data line. Load Y register with Y -coordinate for data line.
Set TV.PTR to point to the start of the data line.
Display high byte of the currently selected address. Display low byte of the currently selected address. Skip one space after address field. Look up value of the currently selected byte.
Save it.
Display it, in hexadecimal format, in field 1.
Skip one space after field 1. Restore value of currently selected byte. Display that byte, in graphic form, in field 2.
Skip one space after field 2.
Display 6502 register images in fields 4 thru 7 :

Look up the register image. Display it in hexadecimal format. Skip one space after hexadecimal field. Get ready for next register...
Done 4 registers yet?
If not, do next one... If all registers displayed, return.

## Get Currently Selected Byte

Note that the subroutine LINE.2, which puts up the second line of the Visible Monitor's display, does not itself "know" the value of the currently selected byte. Rather, it calls a subroutine, GET.SL, which returns the contents of the address pointed to by SELECT. That makes life easy for LINE.2, but how does GET.SL work?

If SELECT were a zero-page pointer, GET. SL could be a very simple subroutine and take advantage of the 6502's indirect addressing mode:

GET.SL
LDY \#O
LDA (SELECT), Y
RTS

Get the zeroth byte above the address pointed to by SELECT. Return to caller.

However, SELECT is not a zero-page pointer; it's up in page $\$ 12$. And the 6502 doesn't have an addressing mode that will let us load a register using any pointer not in the zero page. So how can we see what's in the address pointed to by SELECT?

We can do it in two steps. First, we'll set a zero-page pointer equal in value to the SELECT pointer, so it points to the same address; and then, since we already know how to load the accumulator using a zero-page pointer, we'll load the accumulator using the zero-page pointer that now equals SELECT. Let's call that zeropage pointer GETPTR, since it will allow us to get the selected byte. Using such a strategy, GET.SL can look like this:

GET.SL
LDA SELECT
STA GETPTR
LDA SELECT +1
STA GETPTR +1
LDY \#0
LDA (SELECT),Y
RTS STA GETPTR LDA SELECT +1 STA GETPTR + 1 LDY \#0 LDA (SELECT), Y RTS

Set GETPTR equal to SELECT: first the low byte;
then the high byte.
Get the zeroth byte above the address pointed to by GETPTR.
Return to caller, with A bearing the contents of the address specified by SELECT.

This second attempt at GET.SL will load the accumulator with the currently selected byte, even when SELECT is not in the zero page. However, beware because by setting GETPTR equal to SELECT, GET.SL changes the value of GETPTR. This can be very dangerous. What, for example, if some other program were using GETPTR for something? That other program would be sabotaged by GET.SL's actions. If we let GET. SL change the value of GETPTR, then we must make sure that
no other program ever uses GETPTR.
Such policing is hard work - and almost impossible if you want your software to run on a system in conjunction with software written by anyone else. Since I want the Visible Monitor to share your system's ROM input/output routines, and since I have no way of knowing what zero-page addresses those routines may use, I must refrain from using any of those zero-page bytes myself. When I have to use zeropage bytes - as now, so that GET.SL can use the 6502's indirect addressing mode I must restore any zero-page bytes I've changed.

Therefore, GET.SL must be a four-part subroutine, which will: 1) save GETPTR; 2) set GETPTR equal to SELECT; 3) load the accumulator with the contents of the address pointed to by GETPTR; and finally, 4) restore GETPTR to its original value. This larger, slower, but infinitely safer version of GET.SL looks like this:

## GET.SL LDA GETPTR PHA LDX GETPTR + 1

LDA SELECT STA GETPTR LDA SELECT +1 STA GETPTR + 1

LDY \#0
LDA (GETPTR), Y
TAY
PLA
STA GETPTR
STX GETPTR + 1
TYA
RTS

Save GETPTR on stack and in X register.

Set GETPTR
equal to
SELECT.

Get the contents of the byte pointed to by SELECT, and save it in $Y$ register.

Restore GETPTR from stack and from X register. Restore contents of current byte from temporary storage in Y to A . Return with contents of currently selected byte in accumulator and with the zero page preserved.

## Display Arrow Line

This routine displays an up-arrow directly underneath the current field:

## LINE. 3

| LINE. 3 | LDX \#2 | Set TV.PTR to |
| :---: | :---: | :---: |
|  | LDY \#4 | beginning of |
|  | JSR TVTOXY | arrow line. |
|  | LDY FIELD | Look up current field. |
|  | SEC | If it is out of bounds, |
|  | CPY \#7 | set it to |
|  | BCC FLD.OK | default field |
|  | LDY \#0 | (the address field). |
|  | STY FIELD |  |
| FLD.OK | LDA FIELDS, Y | Look up column number for current field. |
|  | TAY | Use that column number as an index into the row. |
|  | LDA ARROW | Load accumulator with your system's graphic code for up-arrow. |
|  | STA (TV.PTR), Y | Store up-arrow code in the Yth column of the arrow line. |
|  | RTS | Return to caller. |
| FIELDS | .BYTE 3,6,8 <br> BYTE \$0B, $\$ 0 \mathrm{E}$ | This data area shows which column should get an up-arrow to indicate |
|  | .BYTE \$11,\$14 | any one of fields 0 thru 6 . Changing one of these values will cause the up-arrow to appear in a different column when indicating a given field. |

Now that we have all the routines we need for the monitor display, let us look at how they fit together to form a structure. Here is the hierarchy of subroutines in DSPLAY:

When DSPLAY is called, it will clear the top four rows of the screen, display labels, data, the arrow, and then return. How long do you think it will take to do all this? The code may look cumbersome, but the display is quick!

## Monitor Update

The UPDATE routine is the monitor subroutine that executes functions in response to various keys. The basic key functions we want to implement are as follows:

Key
GREATER THAN
LESS THAN
SPACEBAR
RETURN

## Function

Move arrow one field to the right. Move arrow one field to the left. Increment address being displayed. (Step forward through memory.) Decrement address being displayed. (Step backward through memory.)

If the arrow is in fields $1,3,4,5$, or 6 , then, for
keys 0 thru 9, A thru F
Roll a hexadecimal character into the field pointed to by the arrow.

If the arrow is under field 2 (the graphic field) then, for
All keys
Enter the key's character into field 2 (ie: enter the key's character into the displayed address).

Since the video display need not be refreshed (redisplayed within a given time) by the processor, the UPDATE routine need not return within a given amount of time. The UPDATE routine, therefore, can wait indefinitely for a new character from the keyboard, and then take appropriate action.

We can diagram these functions as shown in figure 6.5. You add additional functions to this routine by adding additional code to test the input character. You then call the appropriate function subroutine which you write.


## Get a Key

First we need a way to get a key from the keyboard. I assume that your system has a read-only memory routine to perform this function. Place the address of that routine (see the appropriate appendix for your system) into a pointer called ROMKEY located at address $\$ 1008$. Once you have set the ROMKEY pointer, you can get a key by calling a subroutine labeled GETKEY, which simply transfers control to the ROM routine whose address you placed in ROMKEY:
GETKEY JMP (ROMKEY)

Now that we have a way to get a key from the keyboard, we should be able to write source code for the monitor-update routine:

## Update

| UPDATE | JSR GETKEY | Get a character from the keyboard. |
| :---: | :---: | :---: |
| IF.GRTR | CMP \#' $>$ | Is it the GREATER THAN key? |
|  | BNE IF.LSR | If not, perform next test. |
| NEXT.F | INC FIELD | If so, select the next field. |
|  | LDA FIELD | If arrow was at the right-most field, |
|  | CMP \#7 | place it underneath the left-most |
|  | BNE EXIT. 1 | field. |
|  | LDA \#0 |  |
|  | STA FIELD |  |
| EXIT. 1 | RTS | Then return. |
| IF.LSR | CMP \#'< | Is it the LESS THAN key? |
|  | BNE IF.SP | If not, perform next test. |
| PREV.F | DEC FIELD | If so, select previous field: |
|  | BPL EXIT. 2 | the field to the left of the |
|  | LDA \#6 | current field. If arrow was at |
|  | STA FIELD | left-most field, place it under right-most field. |
| EXIT. 2 | RTS | Then return. |
| IF.SP | CMP \#SPACE | Is it the space bar? |
|  | BNE IF.CR | If not, perform next test. |
| INC.SL | INC SELECT | If so, step forward through |
|  | BNE EXIT. 3 | memory, by incrementing the |
|  | INC SELECT +1 | pointer that specifies the displayed address. |
| EXIT. 3 | RTS | Then return. |
| IF.CR | CMP \#CR | Is it carriage return? |
|  | BNE IFCHAR | If not, perform next test. |


| DEC.SL | LDA SELECT | If so, step backward through |
| :---: | :---: | :---: |
|  | BNE NEXT. 1 | memory by decrementing the |
|  | DEC SELECT+1 | pointer that selects the |
| NEXT. 1 | DEC SELECT | address to be displayed. |
|  | RTS | Then return. |
| IFCHAR | LDX FIELD | Is arrow underneath the |
|  | CPX \#2 | character field (field 2)? |
|  | BNE IF.GO | If not, perform next test. |
|  |  | Put the contents of A into the currently selected address. |
| PUT.SL | TAY | Use $Y$ to hold the character we'll put in the selected address. |
|  | LDA TV.PTR | Save zero-page pointer TV.PTR |
|  | PHA | on stack and in X before we |
|  | LDX TV.PTR+1 | use it to put character in selected address. |
|  | LDA SELECT | Set TV.PTR equal to SELECT, |
|  | STA TV.PTR | so it points to the |
|  | LDA SELECT+1 | currently selected |
|  | STA TV.PTR+1 | address. |
|  | TYA | Restore to A the character we'll put in the selected address. |
|  | LDY \#0 | Store it in the |
|  | STA (TV.PTR), Y | selected address. |
|  | STX TV.PTR+1 | Restore TV.PTR to |
|  | PLA | its original value. |
|  | STA TV.PTR |  |
|  | RTS | Return to caller, with character originally in A now in the selected address and with zero page unchanged. |
|  | RTS | Then return. |
| IF.GO | CMP \#'G | Is it ' $\mathrm{G}^{\prime}$ for GO? |
|  | BNE IF.HEX | If not, perform next test. |
| GO | LDY REG.Y | If so, load the 6502 registers |
|  | LDX REG.X | with their displayed images. |
|  | LDA REG.P |  |
|  | PHA |  |
|  | LDA REG.A |  |
|  | PLP |  |
|  | JSR CALLSL | Call the subroutine at the selected address. |
|  | PHP | When subroutine returns, |
|  | STA REG.A | save register values in register |
|  | STX REG.X | images. |


|  | STY REG.Y |  |
| :---: | :---: | :---: |
|  | PLA |  |
|  | STA REG.P |  |
|  | RTS | Then return to caller. |
| CALLSL | JMP (SELECT) | Call the subroutine at the selected address. |
| IF.HEX | PHA | Save keyboard character. |
|  | JSR BINARY | If accumulator holds ASCII character for 0 thru 9 or A thru F, BINARY returns the binary representation of that hexadecimal digit. Otherwise BINARY returns with $\mathrm{A}=\mathrm{FF}$ and the minus flag set. |
|  | BMI OTHER | If accumulator did not hold a hexadecimal character, perform next test. |
|  | TAY |  |
|  | PLA |  |
|  | TYA |  |
| ROLLIN | LDX FIELD | Roll A into a hexadecimal field. Is arrow underneath the address field (field 0 )? If not, the arrow must be under another hexadecimal field. |
|  | BNE NOTADR |  |
| ADRFLD | LDX \#3 | Since arrow is underneath the address |
| LOOP. 1 | CLC | field, roll accumulator's hexadecimal digit into the address field by rolling it into the pointer that selects the displayed address. |
|  | ASL SELECT |  |
|  | ROL SELECT + 1 |  |
|  | DEX |  |
|  | BPL LOOP. 1 |  |
|  | TYA |  |
|  | ORA SELECT |  |
|  | STA SELECT |  |
|  | RTS | Then return. |
| NOTADR |  | Is arrow underneath field 1 ? <br> If not, it must be underneath a register image. |
|  | BNE REGFLD |  |
| ROL.SL | AND \#\$0F | Roll A's 4 LSB into contents of currently selected byte. Get the contents of the selected address and shift left 4 times. |
|  | PHA |  |
|  | JSR GET.SL |  |
|  | ASL A |  |
|  | ASL A |  |
|  | ASL A |  |
|  | ASL A |  |
|  | AND \#SFO |  |
|  | STA TEMP | Save it in a temporary variable. |


|  | PLA | Get original A's 4 LSB and |
| :---: | :---: | :---: |
|  | ORA TEMP | OR them with shifted contents of selected address. |
|  | JSR PUT.SL RTS | Store the result in the selected address and return. |
| TEMP | .BYTE 0 | This byte holds the temporary variable used by ROL.SL. |
| REGFLD | DEX | The arrow must be underneath a |
|  | DEX | register image - field 3, 4, 5, or 6 . |
|  | DEX |  |
|  | LDY \#3 |  |
| LOOP. 2 | CLC | Roll accumulator's hexadecimal digit |
|  | ASL REGS, X | into appropriate register image... |
|  | DEY |  |
|  | BPL LOOP. 2 |  |
|  | ORA REGS, X |  |
|  | STA REGS, X |  |
|  | RTS | ...Then return. |
| OTHER | PLA | Restore the raw keyboard character that we saved on the stack. |
|  | CMP\#'Q | Is it ' $Q^{\prime}$ for Quit? |
|  | BNE NOT.Q | If not, perform next test. |
|  | PLA | If so, return to |
|  | PLA | the caller of |
|  | PLP |  |
|  | RTS | VISMON. |
| NOT.Q | JSR DUMMY | Replace this call to DUMMY with a call to any other subroutine that extends the functionality of the Visible Monitor. |
| DUMMY | RTS | Return to caller. |

## ASCII to BINARY Conversion

The Visible Monitor's UPDATE subroutine requires a subroutine called BINARY, which will determine if the character in the accumulator is an ASCII 0 thru 9 or A thru F, and, if so, return the binary equivalent. On the other hand, if the accumulator does not contain an ASCII 0 thru 9 or A thru F, BINARY will return an error code, \$FF. Thus:

## If accumulator holds

\$30 (ASCII "0") ..... $\$ 00$
\$31 (ASCII "1") ..... $\$ 01$
\$32 (ASCII " 2 ") ..... $\$ 02$
\$33 (ASCII " 3 ") ..... $\$ 03$
\$34 (ASCII "4") ..... $\$ 04$
\$35 (ASCII "5") ..... $\$ 05$
\$36 (ASCII "6") ..... $\$ 06$
\$37 (ASCII "7") ..... $\$ 07$
\$38 (ASCII "8") ..... $\$ 08$
\$39 (ASCII " 9 ") ..... $\$ 09$
\$41 (ASCII "A") ..... \$0A
\$42 (ASCII "B") ..... \$0B
\$43 (ASCII "C") ..... \$0C
\$44 (ASCII "D") ..... \$0D
\$45 (ASCII "E") ..... \$0E
\$46 (ASCII "F") ..... \$0F
Any other value ..... \$FF

We could solve this problem with a table, BINTAB, for BINary TABle. If BINTAB is at address $\$ 2000$, then $\$ 2000$ would contain a $\$$ FF, as would $\$ 2001$, $\$ 2002$, and all addresses up to $\$ 202 \mathrm{~F}$, because none of the ASCII codes from $\$ 00$ thru $\$ 2 \mathrm{~F}$ represent any of the characters 0 thru 9 or A thru F. On the other hand, address $\$ 2030$ would contain 00 , because $\$ 30$ (its offset into the table) is an ASCII zero, so $\$ 2030$ gets its binary equivalent: $\$ 00$, a binary zero. Similarly, since $\$ 31$ is an ASCII '1,' address $\$ 2031$ would contain a binary '1:' $\$ 01$. $\$ 2032$ would contain a $\$ 02$; $\$ 2033$ would contain a $\$ 03$, and so on up to $\$ 2039$, which would contain a $\$ 09$.

Addresses $\$ 203 \mathrm{~A}$ thru $\$ 2040$ would each contain $\$ \mathrm{FF}$, because none of the ASCII codes from $\$ 3 \mathrm{~A}$ thru $\$ 40$ represent any of the characters 0 thru 9 or A thru F. On the other hand, address $\$ 2041$ would contain a $\$ 0 \mathrm{~A}$, because $\$ 41$ is an ASCII ' A ' and $\$ 0 \mathrm{~A}$ is its binary equivalent: a binary 'A.' By the same reasoning, $\$ 2042$ would contain $\$ 0 B$; $\$ 2043$ would contain $\$ 0 C$, and so on up to $\$ 2046$, which would contain $\$ 0 \mathrm{C}$, and so on up to $\$ 2046$, which would contain $\$ 0 \mathrm{~F}$. Addresses $\$ 2047$ thru $\$ 20 \mathrm{FF}$ would contain $\$$ FFs because none of the values $\$ 47$ thru $\$$ FF is an ASCII 0 thru 9 or A thru F .

To use such a table, BINARY need only be a very simple routine:

BINARY
TAY
LDA BINTAB, Y
RTS

Use ASCII character as an index. Look up entry in BINary TABle. Return with it.

This is a typical example of a fast and simple table lookup code. But it requires a 256-byte table. Perhaps slightly more elaborate code can get by with a smaller table, or do away altogether with the need for a table. Such code must calculate, rather than look up, its answers. Let's look closely at the characters we must convert.

Legal inputs will be in the range $\$ 30$ thru $\$ 39$ or the range $\$ 41$ thru $\$ 46$. An input in the range $\$ 30$ thru $\$ 39$ is an ASCII 0 thru 9, and subtracting $\$ 30$ from such an input will convert it to the corresponding binary value. An input in the range $\$ 41$ thru $\$ 46$ is an ASCII A thru F, so subtracting $\$ 36$ will convert it to its corresponding binary value. For example, $\$ 41$ (an ASCII ' $\mathrm{A}^{\prime}$ ) minus $\$ 36$ equals $\$ 0 \mathrm{~A}$ (a binary ' A '). Any value not in either of these ranges is illegal and should cause BINARY to return a \$FF.

Given these input/output relationships, BINARY need only determine whether the character in the accumulator lies in either legal range, and if so perform the appropriate subtraction, or, if the accumulator is not in a legal range, then return a \$FF.

Here's some code for BINARY which makes these judgments, thus eliminating the need for a table:

| BINARY | SEC | Prepare to subtract. |
| :---: | :---: | :---: |
|  | SBC \#\$30 | Subtract \$30 from character. |
|  | BCC BAD | If character was originally less than $\$ 30$, it was bad, so return \$FF. |
|  | CMP \#\$0A | Was character in the range $\$ 30$ thru \$39? |
|  | BCC GOOD | If so, it was a good input, and we've already converted it to binary by subtracting $\$ 30$, so we'll return now with the character's binary equivalent in the accumulator. |
|  | SBC \#7 | Subtract 7. |
|  | CMP \#\$10 | Was character originally in the range $\$ 41$ thru \$46? |
|  | BCS GOOD | If so, it was a good input, and we've already converted it to binary by subtracting \$37, so we'll return now with the character's binary equivalent in the accumulator. |
| BAD | LDA \#\$FF | Indicate a bad input by returning |
|  | RTS | minus, with A holding \$FF. |
| GOOD | LDX \#0 | Indicate a good input by returning |
|  | RTS | plus, with A holding the character's binary equivalent. |

## Visible Monitor Utilities

The Visible Monitor makes the following subroutines available to external callers:

BINARY Determine whether accumulator holds the ASCII representation for a hexadecimal digit. If so, return binary representation for that digit. If not, return an error code (\$FF).

CALLSL DEC.SL GETKEY

GET.SL GO

INC.SL
PUT.SL VISMON

Call the currently selected address as a subroutine. Select previous address, by decrementing SELECT pointer. Get a character from the keyboard by calling machine's read-only memory routine indirectly.
Get byte at currently selected address.
Load registers from displayed images and call displayed address. Upon return, restore register images from registers.
Select next byte (increment SELECT pointer).
Store accumulator at currently selected address.
Let user give the Visible Monitor commands until user presses ' $Q$ ' to quit.

Figure 6.6 illustrates the hierarchy of the various routines of the Visible Monitor, some of which are detailed in later chapters.


Figure 6.6: A hierarchy of the routines of the Visible Monitor.

Use the minimal machine-language monitor on your computer to enter the Visible Monitor into memory; then have your monitor pass control to the Visible Monitor. The Visible Monitor display should appear in the upper portion of your video display. If it's not fully visible, adjust the value HOME in the screen parameters (HOME is the pointer at $\$ 1000$ ). Use the GREATER THAN and LESS THAN character keys to move the arrow from field to field. Place the arrow under field 0 and roll hexadecimal characters into the address. Select an address in the lower portion of screen memory and use the Visible Monitor to place characters on the screen. Enter characters to the screen using both field 1 (the hexadecimal data field), and field 2 (the character field).

Select the address of the TVT routine in your system. Press G to call that subroutine. You should see the character in the accumulator print on the screen. Try exploring other memory locations. Try writing to a read-only memory address. Why doesn't that work? Try writing to the upper portion of the screen. Why doesn't that work?

## Chapter 7:

## Print Utilities

The Visible Monitor is a useful tool for examining and modifying memory, but at the moment it's mute: it can't "talk" to you except through the limited device of the fields in its display. You can use the Visible Monitor's character entry feature to place ASCII characters directly into screen memory, thus putting messages on the screen manually. However, as yet we have no subroutines to direct a complete message, report, or other string of characters to the screen, to a printer, or to any other output device.

Most programs require some means of directing messages to the screen, thus providing the user with the basis for informed interaction, or to a printer, thus providing a record of that interaction. This chapter presents a set of print utilities to perform these functions.

Fortunately, there are subroutines in your computer's operating system to perform character output. The Apple, Atari, OSI and PET computers each feature a routine to print a character on the screen, thus simulating a TVT (TeleVision Typewriter), and they each feature another routine to send a character to the device connected to the serial output port: usually a printer. I don't plan to reinvent those wheels in this chapter. Rather, the chapter's software will funnel all character output through code that calls the appropriate subroutine in your computer's operating system. And since we're going to have code that calls the two standard character output routines, why not provide a hook to a user-written character output routine, as well? Such a feature will make it trivial for you to direct any character output (eg: messages, hexdumps, disassembler listings, etc) to the screen and the printer, or to any special output device you may have on your system, provided that you've written a subroutine to drive that device.

## Selecting Output Devices

It should be possible for any program to direct character output to the screen, and/or to the printer, and/or to the user-written subroutine. Therefore, we'll need subroutines to select and deselect (stop using) each of these devices and to select and deselect all of these devices. Let's call these routines TVT.ON, TVTOFF, PR.ON, PR.OFF, USR.ON, USR.OFF, ALL.ON, and ALLOFF. With these subroutines, a calling program can select or deselect output devices individually or globally.

The line of source code which will select the TVT as an output device follows:

## JSR TVT.ON

This line will deselect the TVT:

## JSR TVTOFF

That's a pretty straightforward calling sequence.
The select and deselect subroutines will operate on three flags: TVT, PRINTR, and USER. The TVT flag will indicate whether the screen is selected as an output device; the PRINTR flag will indicate whether the printer is selected as an output device; and the USER flag will indicate whether the user-provided subroutine is selected as an output device.

For convenience, we'll have a separate byte for each flag and define a flag as "off" when its value is zero, and "on" when its value is nonzero.

Using this definition of a flag, we can select a given device simply by storing a nonzero value in the flag for that device; we can deselect a device simply by storing a zero in the flag for that device.

The definitions for the flags and listings of the select and deselect subroutines follow:

## Device Flags

| OFF $=0$ | When a device flag $=$ zero, that device <br> is not selected. |
| :--- | :--- |
| ON $=\$$ FF | When a device flag $=\$$ FF, that device is <br> selected. |
| ThYTE ON | This flag is zero if TVT is not selected; <br> nonzero otherwise. Initially, the TVT is <br> selected. |


| PRINTR USER | .$B Y T E ~ O F F$ .BYTE OFF | This flag is zero if the PRINTR is not selected; nonzero otherwise. Initially, the printer is not selected. <br> This flag is zero if the user-provided output subroutine is not selected; nonzero otherwise. Initially, the userprovided function is deselected. |
| :---: | :---: | :---: |
| Select and Deselect Subroutines |  |  |
| TVT.ON | LDA \#ON <br> STA TVT <br> RTS | Select TVT as an output device by setting the flag that indicates the "select" state of the TVT. |
| TVTOFF | LDA \#OFF <br> STA TVT <br> RTS | Deselect TVT as an output device by clearing the flag that indicates the "select" state of the TVT. |
| PR.ON | LDA \#ON STA PRINTR RTS | Select printer as an output device by setting the flag that indicates the "select" state of the printer. |
| PR.OFF | LDA \#OFF STA PRINTR RTS | Deselect printer as an output device by clearing the flag that indicates the "select" state of the printer. |
| USR.ON | LDA \#ON <br> STA USER <br> RTS | Select user-written subroutine as an output device by setting the flag that indicates the "select" state of the output routine provided by the user. |
| USROFF | LDA \#OFF STA USER RTS | Deselect user-written subroutine as an output device by clearing the flag that indicates the "select" state of the output routine provided by the user. |
| ALL.ON | JSR TVT.ON JSR PR.ON JSR USR.ON RTS | Select all output devices by selecting each output device individually. |
| ALLOFF | JSR TVTOFF JSR PR.OFF JSR USROFF RTS | Deselect all output devices by deselecting each output device individually. |

## A General Character-Print Routine

Now that a calling routine can select or deselect any combination of output devices, we need a routine that will output a given character to all currently selected output devices. Let's call this routine PR.CHR, because it will PRint a CHaRacter.

All the software in this book that outputs characters will do so by calling PR.CHR; none of that software will call your system's character-output routines directly. That makes the software in this book much easier to maintain. If you ever replace your system's TVT output routine or its printer-output routine with one of your own, you won't have to change the rest of the software in this book. That software will continue to call PR.CHR. However, if many lines of code in many places called your system's character-output routines directly, then replacing a read-only memory output routine with one of your own would require you to change many operands in many places. Who needs to work that hard? Funneling all character output through one routine, PR.CHR, means we can improve our character output in the future without difficulty.

When it is called, PR.CHR will look at the TVT flag. If the TVT flag is set, it will call your system's TVT output routine. Then it will look at the PRINTR flag. If the PRINTR flag is set, it will call your system's routine that sends a character to the serial output port. Finally, it will look at the USER flag. If the USER flag is set, it will call the user-provided character-output routine. Having done all of this, PR.CHR can return. Figure 7.1 is a flowchart for PR.CHR.

Figure 7.1: To print a character to all currently selected output devices (PR.CHR, a general character-output routine).


## Output Vectors

If the character output routines are located at different addresses in different systems, how can PR.CHR know the addresses of the routines it must call? It can't. But it can call those subroutines indirectly, through pointers that you set.

You must set three pointers, or output vectors, so that they point to the character output routines in your system. A pointer called ROMTVT must point to your system's TVT output routine; a pointer called ROMPRT must point to your system's routine that sends a character to the serial output port; and a pointer called USROUT must point to your own, user-written, character-output routine. (If you have not written a special character-output subroutine, USROUT should point to a dummy routine which is nothing but an RTS instruction.) Then, if you ever relocate your TVT output routine, your printer-output routine, or your user-written output routine, you'll only have to change one output vector: ROMTVT, ROMPRT, or USROUT. Everything else in this book can remain the same.

ROMTVT, ROMPRT, and USROUT need not be located anywhere near PR.CHR. That means we can keep all the pointers and data specific to your system in one place. We can store the output vectors with the screen parameters, in a single block of memory called SYSTEM DATA. See Appendix B1, B2, B3, or B4 for your computer.

The source code of the PR.CHR routine follows:

## PR.CHR

| PR.CHR | STA CHAR |
| :--- | :--- |
|  | BEQ EXIT |
|  | LDA TVT |
|  | BEQ IF.PR |
|  | LDA CHAR |
|  | JSR SEND.1 |
| IF.PR | LDA PRINTR |
|  | BEQ IF.USR |
|  | LDA CHAR |
|  | JSR SEND. 2 |
|  | LDA USER |
|  |  |
|  | BEQ EXIT |
|  | LDAS CHAR |
|  | JSR SEND.3 |
|  | RTS |
| CHAR | .BYTE 0 |

Save the character.
If it's a null, return without printing it. Is TVT selected?
If not, test next device.
If so, send character indirectly to system's TVT output routine.
Is printer selected?
If not, test next device. If so, send character indirectly to system's printer driver. Is user-written output subroutine selected?
If not, test next device. If so, send character indirectly to user-written output subroutine. Return to caller.
This byte holds the last character passed to PR.CHR.

## Vectored Subroutine Calls

| SEND. 1 | JMP (ROMTVT) |
| :--- | :--- |
| SEND. 2 | JMP (ROMPRT) |
| SEND. 3 | JMP (USROUT) |

## Specialized Character-Output Routines

Given PR.CHR, a general character-output routine, we can write specific character-output routines to perform several commonly required functions. For example, it's often necessary for a program to print a carriage return and a line feed, thus causing a new line, or to print a space, or to print a byte in hexadecimal format. Let's develop several dedicated subroutines to perform these functions. Since each of these subroutines will call PR.CHR, their output will be directed to all currently selected output devices.

Here are source listings for a few such subroutines: CR.LF, SPACE, and PR.BYT:

## PRINT A CARRIAGE RETURN-LINE FEED

$$
\mathrm{CR}=\$ 0 \mathrm{D} \quad \text { ASCII carriage return character }
$$

$\mathrm{LF}=\$ 0 \mathrm{~A} \quad$ ASCII line feed character.
$\left.\begin{array}{lll}\text { CR.LF } & \begin{array}{l}\text { LDA \#CR } \\ \text { JSR PR.CHR }\end{array} & \begin{array}{l}\text { Send a carriage return and a } \\ \text { line feed to the currently selected } \\ \text { device(s). }\end{array} \\ \text { LDA \#LF }\end{array}\right)$

## PRINT A SPACE

| SPACE | LDA \#\$20 | Load accumulator with ASCII space. <br> Print it to all currently selected output <br> devices. |
| :--- | :--- | :--- |
|  | JSR PR.CHR | Return. |

## PRINT BYTE

| PR.BYT | PHA | Save byte. |
| :--- | :--- | :--- |
|  | LSR A | Determine ASCII for the 4 MSB (most- |

LSR A
LSR A
LSR A
JSR ASCII
JSR PR.CHR
PLA
JSR ASCII
JSR PR. CHR
RTS
significant bits) in the byte:
byte:

Print that ASCII character to the current device(s). Determine ASCII for the 4 LSB (leastsignificant bits) in the byte that was passed to this subroutine. Print that ASCII character to the current device(s).
Return to caller.

## Repetitive Character Output

Since some calling programs might need to output more than one space, a new line, or other character, why not have a few print utilities to perform such repetitive character outputs? In each case, the calling program need only load the $X$ register with the desired repeat count. Then it would call SPACES to print $X$ spaces, CR.LFS to print $X$ new lines, or CHARS to print the character in the accumulator $X$ times. Calling any of these routines with zero in the $X$ register will cause no characters to be printed. To output seven spaces, a calling program would only have to include the following two lines of code:

LDX \#7
JSR SPACES

To output four blank lines, a program would require these two lines of code:

> LDX \#4
> JSR CR.LFS

To output ten asterisks, a program would need these three lines of code:

> LDA \#*
> LDX \#10
> JSR CHARS

In order to support these calling sequences, we'll need three small subroutines, SPACES, CR.LFS, and CHARS:

## Print X Spaces; Print X Characters

| SPACES | LDA \#\$20 | Load accumulator with ASCII space. |
| :--- | :--- | :--- |
| CHARS | STX REPEAT | Initialize the repeat counter. |
| RPLOOP | PHA | Save character to be repeated. |
|  | LDX REPEAT | Has repeat counter timed out yet? |
|  | BEQ RPTEND | If so exit. If not, |
|  | DEC REPEAT | decrement repeat counter. |
|  | JSR PR.CHR | Print character to all currently selected |
|  |  | output devices. |
|  | PLA |  |
|  | CLC | Loop back to repeat |
|  | BCC RPLOOP | character, if necessary. |
| RPTEND | PLA | Clean up stack. |
|  | RTS | Return to caller. |

Print X New Lines

| CR.LFS | STX REPEAT | Initialize repeat counter. |
| :--- | :--- | :--- |
| CRLOOP | LDX REPEAT | Exit if repeat counter has timed out. |
|  | BEQ END.CR |  |
|  | DEC REPEAT | Decrement repeat counter. |
|  | JSR CR.LF | Print a carriage return and line feed. |
|  | CLC | Loop back to see if done yet. |
|  | RCC CRLOOP |  |
| END.CR | RTS | If done, return to caller. |
| REPEAT | :BYTE | This byte is used as a repeat counter by <br>  |
|  |  |  |

## Print a Message

Some calling programs might need to output messages stored at arbitrary places in memory. So let's develop a subroutine, called PR.MSG, to perform this function. PR.MSG will print a message to all currently selected output devices. It must get characters from the message in a sequential manner and pass each character to PR.CHR, thus printing it on all currently selected output devices.

But how can PR.MSG know where the message starts and ends?
We could require that the message be placed in a known location, but then

PR.MSG would lose usefulness as it loses generality. We could require that a pointer in a known location be initialized so that it points to the start of the message. But that would still tie up the fixed 2 bytes occupied by that pointer. Or we could have a register specify the location of a pointer that actually points to the start of the message. Presumably a calling program can find some convenient 2 bytes in the zero page to use as a pointer, even if it must save them before it sets them. The calling program can set this zero-page pointer so that it points to the beginning of the message, and then set the $X$ register so that it points to that zero-page pointer. Having done so, the calling program may call PR.MSG. Using the indexed indirect addressing mode, PR.MSG can then get characters from the message.

When PR.MSG has printed the entire message, it will return to its caller.
How will PR.MSG know when it has reached the end of the message? We can mark the end of each message with a special character: call it ETX, for End of TeXt. And for reasons which will become clear in Chapter 10, A Disassembler, we'll also start each message with another special character: TEX, for TEXt follows.

If we can develop PR.MSG to work from these inputs, then it won't be hard for a calling program to print any particular message in memory. Let's look at the required calling sequence.

A message, starting with a TEX and ending with an ETX, begins at some address. We'll call the high byte of that address MSG.HI and the low bye of that address MSG.LO. Thus, if the message starts at address $\$ 13 A 9$, MSG. $\mathrm{HI}=\$ 13$ and MSG.LO = \$A9.

MSGPTR is some zero-page pointer. It may be anywhere in the zero page. If the calling program does not have to preserve MSGPTR, it can print the message to the screen with the following code:

JSR TVT.ON
LDA \#MSG.LO
STA MSGPTR LDA \#MSG.HI STA MSGPTR +1 LDX \#MSGPTR JSR PR.MSG

Select TVT as an output device. (Any other currently selected output device will echo the screen output.)
Set MSGPTR
so it points
to the start
of the message.
Set $X$ register so it points to MSGPTR.
Print the message to all currently selected output devices.

If the calling program must preserve MSGPTR, it will have to save MSGPTR and MSGPTR +1 before executing the above lines of code and restore MSGPTR and MSGPTR +1 after executing the above lines of code.

That looks like a reasonably convenient calling sequence. So now let's turn our attention to PR.MSG itself and develop it so it meets the demands of its callers.

## Print a Message

| PR.MSG | STX TEMP.X | Save X register, which specifies message pointer. |
| :---: | :---: | :---: |
|  | LDA 1,X | Save message pointer. |
|  | PHA |  |
|  | LDA 0,X | $\div$ |
|  | PHA |  |
| LOOP | LDX TEMP.X | Restore original value of $X$, so it points to message pointer. |
|  | LDA ( $0, \mathrm{X}$ ) | Get next character from message. |
|  | CMP \#ETX | Is it the end of message indicator? |
|  | BEQ MSGEND | If so, handle the end of the message... |
|  | INC 0,X | If not, increment the message pointer |
|  | BNE NEXT | so it points to the next |
|  | INC 1, X | character in the message. |
| NEXT | JSR PR.CHR | Send the character to all currently selected output devices. |
|  | CLC | Get next character |
|  | BCC LOOP | from message. |
| MSGEND | PLA | Restore message pointer. |
|  | STA 0,X |  |
|  | PLA |  |
|  | STA 1,X |  |
|  | RTS | Return to caller, with MSGPTR preserved. |
| TEMP.X | .BYTE 0 | This data cell is used to preserve the initial value of $X$. |

## Print the Following Text

Even more convenient than PR.MSG would be a routine that doesn't require the caller to set any pointer or register in order to indicate the location of a message. But if no pointer or register indicates the start of the message, how can any subroutine know where the message starts?

It can look on the stack.
Why not have a subroutine, called Print-the-Following, which prints the message that follows the call to Print-the-Following. Since Print-the-Following is longer than six characters, let's shorten its name to "PRINT:", letting the colon in "PRINT:" suggest the phrase "the following." A calling program might then print "HELLO" with the following lines of code:

Whenever the 6502 calls a subroutine, it pushes the address of the subroutine's caller onto the stack. This enables control to return to the caller when the subroutine ends with an RTS, because the 6502 knows it can find its return address on the stack. The subroutine PRINT: can take advantage of this fact by pulling its own return address off the stack, and using it as a pointer to the message that should be printed. When it reaches the end of the message, it can place a new return address on the stack, an address that points to the end of the message. Then PRINT: can execute an RTS. Control will then pass to the 6502 code immediately following the ETX at the end of the message. The source code for PRINT: follows:

## PRINT:

LOOP

ENDIT

PLA
TAX
PLA
TAY
JSR PUSHSL
STX SELECT
STY SELECT +1 JSR INC.SL

JSR INC.SL
JSR GET.SL
CMP \#ETX
BEQ ENDIT
JSR PR.CHR
CLC
BCC LOOP LDX SELECT
LDY SELECT +1

Pull return address from stack and save it in registers $X$ and $Y$.

Save the select pointer, because we're going to use it as a text pointer.
Set SELECT $=$ return address.
Increment SELECT pointer so it points to TEX character. Increment select pointer so it points to the next character in the message. Get character. Is it end of message indicator?
If so, adjust return address and return. If not, print the character to all currently selected devices.
Then loop to get next character...
\(\left.$$
\begin{array}{ll}\text { JSR POP.SL } & \begin{array}{l}\text { Restore select pointer to its original } \\
\text { value. }\end{array}
$$ <br>
TYA \& Push address <br>

of ETX\end{array}\right]\)| PHA | onto the stack. |
| :--- | :--- |
| TXA | Return (to byte immediately following <br> PTX |
| RTS |  |

## Saving and Restoring the SELECT Pointer

Now that a number of subroutines are accessing the contents of memory with the SELECT utilities (GET.SL, PUT.SL, INC.SL and DEC.SL) we should provide yet another pair of SELECT utilities to enable the subroutines to save and restore the SELECT pointer. With such save and restore functions, any subroutine can use the SELECT pointer to access memory, without interfering with the use of the SELECT pointer by other subroutines. PUSHSL will push the SELECT pointer onto the stack and POP.SL will pop the SELECT pointer off the stack. PUSHSL and POP.SL will each preserve $X, Y$, and the zero page.

Save Select Pointer
(Preserving X,Y, and the Zero Page)

| PUSHSL | PLA <br> STA RETURN | Pull return address from stack and <br> store it temporarily in RETURN. |
| :--- | :--- | :--- |
|  | PLA |  |
|  | STA RETURN+1 |  |
|  | LDA SELECT+1 | Push select pointer onto stack. |
|  | PHA |  |
|  | LDA SELECT |  |
| PHA |  |  |
|  | LDA RETURN+1 | Push return address back onto stack. |
| PHA |  |  |
|  | LDA RETURN |  |
|  | PHA | Return to caller. (Caller will find select <br> RTS |
|  |  |  |

## Restore Select Pointer (Preserving $X, Y$, and the Zero Page)

POP.SL

| PLA | Save return address temporarily. |
| :--- | :--- |
| STA RETURN |  |
| PLA |  |
| STA RETURN+1 |  |
| PLA | Restore select pointer from stack. |
| STA SELECT |  |
| PLA |  |
| STA SELECT+1 |  |
| LDA RETURN+1 | Place return address back on stack. |
| PHA |  |
| LDA RETURN |  |
| PHA | Return to caller. |
| RTS | This pointer is used by PUSHSL and <br> WORD 0 |
|  | POP.SL to preserve their return ad- <br> dresses. |

    STA RETURN
    PLA
    STA RETURN+1
    PLA
    STA SELECT
    PLA
    STA SELECT+1
    LDA RETURN +1 Place return address back on stack.
    PHA
    LDA RETURN
    PHA
    RTS
    RETURN .WORD 0

Return to caller.
This pointer is used by PUSHSL and POP.SL to preserve their return addresses.

## Conclusion

With the print utilities presented in this chapter, it should be easy to write the character-output portions of many programs, making it possible for calling programs to select any combination of output devices and to send individual characters, bytes, or complete messages to those devices. The calling programs will be completely insulated from the particular data representations used by the print utilities. The calling programs do not need to know the nature or location of the outputdevice flags or the addresses of the output vectors; they need only know the addresses of the print utilities.

Similarly, although the print utilities use subroutines that operate on the SELECT pointer, the print utilities themselves never access the SELECT pointer directly. They are completely insulated from the nature and location of the SELECT pointer. As long as they know the addresses of the SELECT utilities, the print utilities can get the currently selected byte, select the next or the previous byte, save the SELECT pointer onto the stack, and restore the SELECT pointer from the stack. If at some point we should implement a different representation of "the currently selected byte," we need only change the SELECT utilities; the print utilities, and all other programs which use the SELECT utilities need never change.

Insulating blocks of code from the internal representation of data in other blocks of code makes all the code much easier to maintain. The following print utilities are available to external callers:

CHARS
CR.LF
CR.LFS
PR.BYT
PR.CHR
PR.MSG
PRINT:
SPACE
SPACES

Send the character in the accumulator " X " times to all currently selected output devices. Cause a new line on all currently selected devices. Cause " X " new lines on all currently selected devices.
Print the byte in the accumulator, in hexadecimal representation.
Print the character in the accumulator on all currently selected devices.
Print the message pointed to by a zero-page pointer specified by $X$.
Print the message following the call to "PRINT:". Send a space to all currently selected output devices. Send "X" spaces to all currently selected output devices.

## Exercises

1) Write a printer test program, which sends every possible character from $\$ 00$ to $\$$ FF to the printer.
2) Rewrite the printer test program so that it prints just one character per line.

## Chapter 8:

## Two Hexdump Tools

The Visible Monitor allows you to examine memory, but only 1 byte at a time. You'll quickly feel the need for a software tool that will display or print out the contents of a whole block of memory. This is especially useful if you wish to debug a program. You can't debug a program if you're not sure what's in it. A hexdump tool will show you what you've actually entered into the computer, by displaying the contents of memory in hexadecimal form.

I've developed two kinds of hexdump programs, each for a different type of output device. When I'm working at the keyboard, I want a hexdump routine that dumps from memory to the screen, a line or a group of lines at a time. But for documentation and for program development or debugging away from the keyboard, I want a hexdump routine that dumps to a printer.

Most of the code required to dump from memory will be the same, whether we direct output to the screen or to the printer. However, there are enough differences between the two output devices that it is convenient to have two hexdump programs, one for the screen and one for the printer. Let's call them TVDUMP and PRDUMP.

## TVDUMP

TVDUMP should be very responsive: when you are using the Visible Monitor, a single keystroke should cause one or more lines to be dumped to the screen. But how can TVDUMP know what lines you want to dump? Since the Visible Monitor allows you to select any address by rolling hexadecimal characters into the address field or by stepping forward and backward through memory, we might as well have

TVDUMP dump memory beginning with the currently selected address.
Since we're basing TVDUMP on the Visible Monitor's currently selected address, we can use some of the Visible Monitor's subroutines to operate on that address. GET.SL will get the currently selected byte, and INC.SL will increment the SELECT pointer, thereby selecting the next byte. The print utilities TVT.ON and PR.BYT will let us select the screen as an output device and print the accumulator in hexadecimal representation.

We ought to have TVDUMP provide a dump that will be easily readable, even on the narrow confines of a twenty-five- or forty-column display. That means we can't display a full hexadecimal line ( 16 bytes) on one screen line if we want to have a space between each byte. We can provide hexdumps that split each hexadecimal line into two screen lines. See outputs $A$ and $B$ in figure 8.1.

Output A:

| 0200 | HH | HH | HH | HH | HH | HH | HH | HH | HH |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0208 | HH | HH | HH | HH | HH | HH | HH | HH | HH |
| 0210 |  |  |  |  |  |  |  |  |  |
| 0218 | HH | HH | HH | HH | HH | HH | HH | HH | HH |
| 0 |  |  |  |  |  |  |  |  |  |

-29 columns

Output B:
0200
HH HH HH HH HH HH HH HH 0208
HH HH HH HH HH HH HH HH
0210
HH HH HH HH HH HH HH HH
0218
HH HH HH HH HH HH HH HH
-23 columns

Figure 8.1: Two TVDUMP formats.

One way to provide such a hexdump is shown by the flowchart in figure 8.2. Using this flowchart as a guide, let's develop source code to perform the TVDUMP function:


Figure 8.2: Flowchart of the screen Hexdump Program.

## CONSTANTS

$C R=\$ 0 D$ $\mathrm{LF}=\$ 0 \mathrm{~A}$

GET.SL
INC.SL
PR.BYT
SELECT

Carriage return. Line feed.

## REQUIRED SUBROUTINES

| VARIABLES |  |  |
| :---: | :---: | :---: |
| COUNTR | .BYTE 0 | This byte counts the number of lines dumped by TVDUMP. |
| HEXLNS | .BYTE 4 | Number of hexadecimal lines to be dumped by TVDUMP. (Set this to any number you like. To dump a single hexadecimal line [ 16 bytes] , set HEXLNS = 1.) |
| TVDUMP |  |  |
| TVDUMP | JSR TVT.ON | Select TVT as an output device. <br> (Other devices will echo the dump.) |
|  | LDA HEXLNS | Set COUNTR to the number of lines |
|  | STA COUNTR | to be dumped by TVDUMP. |
|  | LDA SELECT | Set SELECT to beginning |
|  | AND \#\$F8 | of a screen line ( 8 bytes) |
|  | STA SELECT | by zeroing 3 LSB in SELECT. |
|  | LDX \#2 | Skip two lines on the screen. |
|  | JSR CR.LFS |  |
| DUMPLN | JSR PR.ADR | Print the selected address. |
|  | JSR CR.LF | Advance to a new line on the screen. (This call to CR.LF may be replaced |
|  |  | with a call to SPACE on systems with |
|  |  | screens more than 27 columns wide. |
|  |  | This will yield the Output A rather than |

## VARIABLES

## TVDUMP

TVDUMP JSR TVT.ON
LDA HEXLNS
STA COUNTR
LDA SELECT
AND \#\$F8
STA SELECT
LDX \#2
JSR CR.LFS
JSR PR.ADR JSR CR.LF

Select TVT as an output device. (Other devices will echo the dump.)
Set COUNTR to the number of lines to be dumped by TVDUMP. Set SELECT to beginning of a screen line ( 8 bytes) by zeroing 3 LSB in SELECT. Skip two lines on the screen.

Print the selected address. Advance to a new line on the screen. (This call to CR.LF may be replaced with a call to SPACE on systems with screens more than 27 columns wide. This will yield the Output A rather than

|  |  | Output B.) |
| :--- | :--- | :--- |
| DMPBYT | JSR SPACE | Print a space. |
|  | JSR DUMPSL | Dump currently selected byte. |
|  | JSR INC.SL | Select next address by incrementing <br> select pointer. |
|  | LDA SELECT | Is it the beginning of a new |
|  | AND \#O7 | screen line? (3 LSB = 0?) |
|  | BNE DMPBYT | If not, dump next byte.. |
|  | JSR CR.LF | If so, advance to a new line on the |
|  |  | screen. |
|  | LDA SELECT | Does this address mark the beginning of |
|  | AND \#SOF | a new hexadecimal line? |
|  | (4 LSB of SELECT =0?) |  |
|  | BNE IFDONE |  |
|  | JSR CR.LF | If so, skip a line on the screen. |
|  | DEC COUNTR | Dumped last line yet? |
|  | BNE DUMPLN | If not, dump next line. |
|  | JSR TVTOFF | Deselect TVT as an output device. |
|  | RTS | Return to caller. |

## DUMP CURRENTLY SELECTED BYTE

This subroutine gets the currently selected byte (the byte pointed to by SELECT) and prints it in hexadecimal format on all selected devices.

| DUMPSL | JSR GET.SL | Get currently selected byte. |
| :--- | :--- | :--- |
|  | JSR PR.BYT | Print it in hexadecimal format. |
|  | RTS | Return to caller. |

## PRINT ADDRESS

This subroutine prints, on all selected devices, the currently selected address (ie: the value of the SELECT pointer).

| PR.ADR | LDA SELECT +1 | Get the high byte of SELECT... |
| :--- | :--- | :--- |
|  | JSR PR.BYT | (..and print it in hexadecimal format. |
|  | LDA SELECT | Get the low byte of SELECT... |
|  | JSR PR.BYT | ...and print it in hexadecimal format. |
|  | RTS | Then return to caller. |

## PRDUMP

With the subroutine presented thus far in this chapter, we can dump to the screen just by calling TVDUMP. But what if we want to print a hexdump? Is a hexdump program that prints any different from one that dumps to the screen? Can we simply select the printer instead of the TVT and leave the rest of the code the same?

We could. But then we wouldn't be taking full advantage of the printer. TVDUMP produces an output that is easily read within the twenty-five or forty columns of a video display. Most printers can output sixty-four columns or more. We should take advantage of the extra width offered by a printer.

We should also recognize the difference in responsiveness between a screen and a hard-copy device. When I'm using a screen-based hexdump, I don't mind hitting a single key every time I want some lines dumped to the screen. But with a printing hexdump, I don't want to strike a key repeatedly to continue the dump. I don't mind striking a number of keys at the beginning in order to specify the memory to be dumped, but once I've done that I don't want to be bothered again. I want to set it and forget it.

When called, a printing hexdump program should announce itself by clearing the screen and displaying an appropriate title (eg: "PRINTING HEXDUMP"). Then it should ask you to specify the starting address and the ending address of the memory to be dumped.

Once it knows what you want to dump, PRDUMP should print a hexdump of the specified block of memory. For your convenience, PRDUMP should tell you what block of memory it will dump; then it should provide a header for each column of data and indicate the starting address of each line of data. (See the "D" appendices.)

Using the flowchart of figure 8.3 as a guide, we can write source code for the top level of the PRINTING HEXDUMP:

Figure 8.3: To print a Hexdump.

$\left.\left.\begin{array}{lll}\text { PRDUMP } & \begin{array}{l}\text { JSR TITLE } \\ \text { JSR SETADS }\end{array} & \begin{array}{l}\text { Display the title. } \\ \text { Let user set start address and end ad- } \\ \text { dress of memory to be dumped. } \\ \text { (SETADS returns with SELECT = EA, }\end{array} \\ \text { the end address.) }\end{array}\right] \begin{array}{ll}\text { Set SELECT = SA, the starting address. }\end{array}\right\}$

## Get Starting, Ending Address

The printing hexdump program must secure from the user the starting address and the ending address of the memory to be dumped. The subroutine, SETADS, will perform these functions. It will place an appropriate prompt on the screen ("Set Starting Address" or "Set Ending Address") and then allow the user to specify an address.

Putting a prompt on the screen is easy: just select the TVT by calling TVT.ON, call "PRINT:" and follow this call with a TEX (start of text) character, the text of the prompt, and then an ETX (end of text) character. How can we allow the user to specify an address? We could make a subroutine, called GETADR, which gets an address by enabling the user to set some pointer. That sounds mighty familiar - that's what the Visible Monitor does. Conveniently, the Visible Monitor is a subroutine, which returns to its caller when the user presses $Q$ for Quit. Therefore, after putting
the appropriate prompt on the screen, SETADS will call the Visible Monitor. When the Visible Monitor returns, the SELECT pointer will specify the requested address.

## SET STARTING ADDRESS, ENDING ADDRESS

| SETADS | JSR TVT.ON | Select TVT as an output device. All other selected output devices will echo the screen output. |
| :---: | :---: | :---: |
|  | JSR PRINT: | Put prompt on the screen: |
|  | . BYTE TEX |  |
|  | .BYTE CR,LF,LF |  |
|  | .BYTE | 'SET STARTING ADDRESS |
|  | .BYTE | 'AND PRESS " ${ }^{\text {" }}$.' |
|  | .BYTE ETX |  |
|  | JSR VISMON | Call the Visible Monitor, so user can specify a given address. |
|  | JSR SAHERE | Set starting address equal to address set by the user. |
| SET.EA | JSR PRINT: | Put prompt on the screen: |
|  | .BYTE TEX |  |
|  | .BYTE CR,LF,LF |  |
|  | .BYTE | 'SET ENDING ADDRESS' |
|  | .BYTE | 'AND PRESS "Q".' |
|  | .BYTE ETX |  |
|  | JSR VISMON | Call the Visible Monitor, so user can specify a given address. |
|  | SEC | If user tried to set an |
|  | LDA SELECT +1 | ending address less than |
|  | CMP SA+1 | the starting address, |
|  | BCC TOOLOW | make user do it over. |
|  | BNE EAHERE | If SELECT is greater than SA, set $E A=$ SELECT. That will make EA greater than SA. |
|  | LDA SELECT |  |
|  | CMP SA |  |
|  | BCC TOOLOW |  |
| EAHERE | LDA SELECT + 1 | Set EA=SELECT... |
|  | STA EA+1 |  |
|  | LDA SELECT |  |
|  | STA EA |  |
|  | RTS | ... and return. |
| SAHERE | LDA SELECT +1 | Set SA=SELECT... |
|  | STA SA+1 |  |


| TOOLOW | LDA SELECT |  |
| :---: | :---: | :---: |
|  | RTS | ...and return. |
|  | JSR PRINT: | Since user set ending address |
|  | .BYTE STX, <br> .BYTE CR,LF,LR | too low, print error message: |
|  | .BYTE | 'ERROR! ${ }^{\prime}$ |
|  | .BYTE | 'END ADDRESS LESS ' |
|  | .BYTE | 'THAN START ADDRESS, ' |
|  | .BYTE | 'WHICH IS |
|  | .BYTE ETX |  |
|  | JSR PR.SA | Print starting address. ...and let the user set |
|  | JMP SET.EA | the ending address again. |
| SA | .WORD 0 | Pointer to starting address of memory to be dumped. |
| EA | .WORD \$FFFF | Pointer to ending address of memory to be dumped. |

Now that the user can set the starting address and the ending address for a hexdump (or for any other program that must operate on a contiguous block of memory), we should have utilities that print out the starting address, the ending address, or the range of addresses selected by the user. If the user set \$D000 as the starting address and \$D333 as the ending address, we should be able to call one subroutine that prints "\$D000," another that prints "\$D333," and a third that prints "\$D000 - \$D333."

Let's call these subroutines PR.SA, to print the starting address; PR.EA, to print the ending address; and RANGE, to print the range of addresses.

## Print Starting Address

The following subroutine prints the value of SA, the starting address, in hexadecimal format:

| PR.SA | LDA \#'S | Print a dollar sign to <br> indicate hexadecimal. |
| :--- | :--- | :--- |
|  | JSR PR.CHR | LDA SA+1 |
|  | Print high byte of starting address. |  |
|  | JSR PR.BYT |  |
|  | LDA SA | Print low byte of starting address. |
|  | JSR PR.BYT |  |
|  | RTS | Return to caller. |

The following subroutine prints the value of EA, the ending address, in hexadecimal format:

PR.EA
LDA \#' ${ }^{\prime}$
JSR PR.CHR
LDA EA +1
JSR PR.BYT
LDA EA
Print low byte of ending address. JSR PR.BYT
RTS
Return to caller.

## Print Range of Addresses

RANGE
JSR PR.SA
Print starting address.
LDA \#' -
Print a hyphen.
JSR PR.CHR
JSR PR.EA
Print ending address.
RTS
Return to caller.

## HEADER

We want a routine to print an appropriate header for the hexdump. It should accomplish two tasks: identify the block it will dump, and print a hexadecimal digit at the top of every column of hexdump output. Thus, HEADER should produce the output shown between the following lines:

## DUMPING HHHH-HHHH

$$
\begin{array}{llllllllllllllll}
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & A & B & C & D & E
\end{array}
$$

Notice the blank line following the line of hexadecimal characters. This will insure a blank line between the header and the dump itself, making for a more
readable output. (See the hexdumps in the D series of appendices which were produced with PRDUMP.)

Here are a few lines of code to print the first line of the header:

JSR PRINT:
.BYTE TEX,CR,LF
.BYTE 'DUMPING '
.BYTE ETX
JSR RANGE
JSR CR.LF

What about the rest of the header? Since all we want to do is print the hexadecimal digits 0 thru $\$ \mathrm{~F}$, with appropriate spacing between them, the rest of HEADER can just be some code to count from 0 to $\$ F$, convert to ASCII, and print:

## PRINT HEXADECIMAL DIGITS (Version I)

|  | LDX \#7 | Print seven spaces. |
| :--- | :--- | :--- |
|  | JSR SPACES |  |
|  | LDA \#O | Initialize column counter |
| HXLOOP | STA COLUMN | to zero. |
|  | LDA COLUMN | Convert column counter to |
|  | JSR ASCII | an ASCII character and |
|  | JSR PR.CHR | print it. |
|  | LDX \#2 | Space twice after the character. |
|  | JSR SPACES |  |
|  | INC COLUMN | Increment the column counter. |
|  | LDA COLUMN | Loop if counter not greater |
|  | AND \#SFO | than \$0F. |
|  | BEQ HXLOOP |  |
|  | LDX \#2 | Otherwise, skip two lines |
|  | JSR CR.LFS | after the header. |
|  | RTS | Then return. |
| COLUMN | .BYTE 0 | This 1-byte variable is used to count |
|  |  | from 00 to \$OF. |

Version 1 of PRINT HEXADECIMAL DIGITS will work, and in only 49 bytes. But that's 49 bytes of code, which among other things must count and branch, and if for some reason one of those bytes is wrong, Version 1 of PRINT HEXADECIMAL DIGITS will probably go directly into outer space. But we could write PRINT

HEXADECIMAL DIGITS in a much more straightforward manner, which, though somewhat more costly in terms of memory required, will be more readable and less likely to run amuck.

PRINT HEXADECIMAL DIGITS need only call "PRINT:", and follow this call with a text string consisting of the desired hexadecimal digits.

## PRINT HEXADECIMAL DIGITS (Version 2)

JSR PRINT:
.BYTE TEX

| .BYTE | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| .BYTE | $\prime 8$ | 9 | A | B | C | D | E | $\mathrm{F}^{\prime}$ |

.BYTE CR,LF,LF
.BYTE ETX
RTS

Version 2 of PRINT HEXADECIMAL DIGITS requires 60 bytes. But it's more readable than Version 1 of PRINT HEXADECIMAL DIGITS, and it can be modified much more easily: just change the text in the message it prints. You don't have to calculate branch addresses or test the terminal condition in a loop. This is just one example of a programming problem that may be solved in a computation-intensive or a data-intensive manner.

Where other factors are about equal, I prefer data-intensive subroutines, because they're more readable and easier to change. Even in this case, I'm willing to pay the extra 20 bytes for a version of PRINT HEXADECIMAL DIGITS that I don't have to read twice. Hence, PRINT HEXADECIMAL DIGITS Version 2, and not Version 1, will appear in the assembler listings of HEADER in Appendix C5.

## PRLINE

Clearly, most of the work of PRDUMP will be performed by the subroutine PRLINE, which dumps one line of memory to the printer. It will stop when it has dumped 16 bytes (one hexadecimal line) or has dumped through the ending address specified by the user.

As we did for TVDUMP, let's use SELECT as a pointer to the first byte that must be dumped by PRLINE. When PRLINE is called, it must see if the currently selected byte (the byte pointed to by SELECT) is at the start of a hexadecimal line. A byte is at the beginning of a hexadecimal line if the 4 LSB (least-significant bits) of its address are zero. Thus, $\$ 4 \mathrm{ED} 8$ is not the start of a hexadecimal line, but $\$ 4 \mathrm{ED} 0$ is.

If the currently selected byte is not the beginning of a hexadecimal line, PRLINE should space over to the appropriate column for that byte. If the currently selected
byte is at the beginning of a hexadecimal line, PRLINE should print the address of the currently selected byte and space twice.

Once it has spaced over to the proper column, PRLINE need only get the currently selected byte, print it in hexadecimal format, space once, and then do the same for the next byte, until it has dumped the entire line or has dumped the last byte requested by the user.

Figure 8.4 gives a flowchart for the following routine:


## PRLINE

| PRLINE | JSR CR.LF | Advance printhead to a new line. |
| :---: | :---: | :---: |
|  | LDA SELECT | Determine starting |
|  | PHA | column |
|  | AND \#SOF | for this dump. |
|  | STA COLUMN | Now COLUMN holds the number of the column in which we will dump the first byte. |
|  | PLA | Set SELECT pointer to |
|  | AND \#\$F0 | beginning of a hexadecimal line. |
|  | STA SELECT |  |
|  | JSR PR.ADR | Print the selected address. |
|  | LDX \#3 | Space three times - to the |
|  | JSR SPACES | first column. |
|  | LDA COLUMN | Do we dump from the first column? |
|  | BEQ COL.OK | If so, we're at the correct column now. |
| LOOP | LDX \#3 | If not, space three |
|  | JSR SPACES | times for each byte not |
|  | JSR INC.SL | dumped. |
|  | DEC COLUMN |  |
|  | BNE LOOP |  |
| COL.OK | JSR DUMPSL | Dump the currently selected byte. |
|  | JSR SPACE | Space once. |
|  | JSR NEXTSL | Select the next byte in memory, unless we've already dumped through the end address. |
|  | BMI EXIT | (MINUS means we've dumped through the end address.) |
| NOT.EA | LDA SELECT | Dumped entire line? |
|  | AND \#\$0F | (4 LSB of SELECT $=0$ ? |
|  | CMP \#0 | If so, we've dumped the entire line. If not, |
|  | BNE COL.OK | select the next byte and dump it... |
| EXIT | RTS | PRLINE returns MINUS, with $A=\$ F F$, if it dumped through ending address. Otherwise it returns PLUS, with $\mathrm{A}=0$. |

## Select Next Byte

NEXTSL tests to see if SELECT is less than the ending address. If so, it increments SELECT and returns PLUS (with zero in the accumulator). If not, it
preserves SELECT and returns MINUS (with $\$$ FF in the accumulator).

## NEXTSL

NEXTSL SEC
LDA SELECT +1
CMP EA +1
BCC SL.OK
BNE NO.INC

SEC
LDA SELECT
CMP EA
BCS NO.INC
SL.OK JSR INC.SL
LDA \#0
RTS
NO.INC LDA \#\$FF
RTS

Prepare to compare. Is high byte of SELECT less than high byte of end address (EA)? If so, SELECT is less than EA, so it may be incremented. If SELECT is greater than EA, don't increment SELECT. SELECT is in the same page as EA, prepare to compare low bytes: Is low byte of SELECT less than low byte of EA? If not, don't increment it. Since SELECT is less than EA, we may increment it.
Set "incremented" return code and return.
Set "not incremented" return code and return.

Go to Start of Block
GOTOSA sets SELECT $=$ SA, thus selecting the first byte in the block defined by SA and EA:

```
GOTOSA LDA SA
    STA SELECT
    LDA SA+1
    STA SELECT+1
```

Set SELECT
equal to
START ADDRESS
of block.
RTS

Now the two hexdump tools are complete. You may invoke either tool directly from the Visible Monitor by displaying the start address of the given hexdump tool and pressing "G." This will work fine for PRDUMP: you'll get a chance to set the starting address and the ending address that you want to dump, and then you'll see the dump on both the printer and the screen. If you start TVDUMP with a " G " from the Visible Monitor, you'll only get a dump of TVDUMP itself. You won't be able to use TVDUMP to dump any other location in memory. Why? Because TVDUMP dumps from the displayed address, and to start any program with a " $G$ " from the Visible Monitor, you must first display the starting address of that program. Prob-
ably you'd like to be able to use TVDUMP to dump other areas in memory. To do so, you must assign a Visible Monitor key (eg: " H ") to the subroutine TVDUMP, so that the Visible Monitor will call TVDUMP whenever you press that key. See Chapter 12, Extending the Visible Monitor.

## Chapter 9:

## A Table-Driven Disassembler

With the Visible Monitor you can enter object code into your computer. With hexdump tools you can dump that object code to the screen or to a printer. However, you still can't be sure you've entered the instructions you intended to enter unless you refer back and forth from your hexdump to Appendix A4, The 6502 Opcode List. You must verify that every opcode you entered is for the instruction and the addressing mode that you had intended. You must count forward or backward in hexadecimal to make sure that the operands in your branch instructions are correct. If you entered one opcode or operand incorrectly, then even though your handwritten program may be correct, the version in your computer's memory will be wrong.

A disassembler (the opposite of an assembler) can make your life a lot easier by displaying or printing the mnemonics represented by the opcodes you entered into your computer, and by showing you the actual addresses and addressing modes represented by your operands. The disassembler can't know that address 0000 has the label "TV.PTR," but it can let you know that a given instruction operates on address 0000.

A disassembled line includes the following fields:

| Field <br> Number | Field <br> Description |
| :--- | :--- |
| 1. | Mnemonic. |
| 2. | Operand. |
| 3. | Address of opcode. |
| 4. | Opcode in hexadecimal. |

5. First byte of operand (if present) in hexadecimal.
6. $\quad$ Second byte of operand (if present) in hexadecimal.

Here's a disassembled line, with each of the fields numbered:

| 1 | 2 | 3 | 4 | 5 | 6 | (Field Numbers) |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| JSR | 0400 | $08 A C$ | 20 | 00 | 04 | (Disassembled Line) |

As with hexdump tools, I find it convenient to have two disassemblers: one for the screen and one for the printer. The screen-oriented disassembler should direct a certain number of disassembled lines to the screen whenever it is called. On the other hand, the printing disassembler should get a starting address and an ending address from the user and print a continuous disassembly of that portion of memory. As before, when I direct output to a printer I want to set it and forget it.

Whether we disassemble to the screen or to a printer, we will disassemble one line at a time. How can a program disassémble a line? The same way a person does. You look at an opcode in memory and then consult a table such as Appendix A4 to determine the operation represented by that opcode. Each operation has two attributes, a mnemonic and an addressing mode. The procedure is simple. Write the mnemonic; then, from the addressing mode determine whether this opcode takes no operand, a 1-byte operand, or a 2-byte operand. If it takes an operand, look at the next byte or two in memory and then write the operand for the mnemonic.

Thus, if you wish to disassemble object code from some place in memory, and you find an $\$ 8 \mathrm{D}$ at that location, you can determine from Appendix A6 that \$8D represents "store accumulator, absolute mode." Therefore, you'll write: "STA," which is the mnemonic for store the accumulator.

The absolute mode requires a 2-byte operand, so you'll look at the 2 bytes following the $\$ 8 \mathrm{D}$. If $\$ 36$ follows the $\$ 8 \mathrm{D}$ and is itself followed by \$D0, then the disassembled line will look like this:

STA \$D036

That's a lot easier to read than the original 3 bytes of object code:

## DISASSEMBLY

| JSR | 0400 | $1 E 00$ | 20 | 00 | 04 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| JSR | $04 A 0$ | 1E03 | 20 | A0 | 04 |
| LDA | $(0021), Y$ | 1E06 | B1 | 21 |  |
| CLC |  | 1E08 | 18 |  |  |
| BCC | $1 E 00$ | 1E09 | 90 | F5 |  |

## HEXDUMP



Figure 9.1: Disassembly and hexdump of the same object code.

TO DISASSEMBLE ONE LINE:


Figure 9.2: Algorithm for disassembling one line of code.

That looks pretty simple. We can use the SELECT pointer to indicate the current byte within memory, and we'll assume that lower-level subroutines exist or will exist to do the jobs required by DSLINE, which disassembles one line. With those assumptions, we can write source code for DSLINE:

## DISASSEMBLE ONE LINE



## Print Mnemonic

We need a subroutine called MNEMON which prints the three-letter mnemonic for a given opcode. How can MNEMON do this? How do we do it? We look it up in a table such as Appendix A4. We could have a similar table in memory and then have MNEMON sequentially look up from the table the three characters comprising the desired mnemonic. That would require a 3-byte mnemonic for each of 256 possible opcodes: a 758-byte table. That's a lot of memory! Perhaps if we organize our data better we'll need less memory.

For example, why include the same mnemonic more than once in the table? Eight different opcodes use the mnemonic LDA; why should I use up 24 bytes to store "LDA" eight times? We could have a table of mnemonic names, which is nothing more than an alphabetical list of the three-letter mnemonics. There are only fifty-six different mnemonics; if we add one pseudo-mnemonic, "BAD," to mean that a given opcode is not valid, then we still have only fifty-seven mnemonics. The table of mnemonic names will therefore require only 171 bytes.

If you have a given opcode, how can you know which mnemonic in the table of mnemonic names corresponds to your opcode? A mnemonic code is some number that uniquely identifies a given mnemonic. Let's assume that we have a table of mnemonic codes which gives the mnemonic code for each possible opcode.

Now you can look up in the table of mnemonic codes the mnemonic code corresponding to a given opcode, and then use the mnemonic code as an index to the table of mnemonic names. The three sequential characters located in the table of mnemonic names will comprise the mnemonic for your original opcode.

This method requires not one but two tables. The two together, however, require considerably less memory than our first table did. The table of mnemonic codes will be 256 -bytes long, since it must have an entry for every possible opcode, including invalid ones. The table of mnemonic names, on the other hand, will be only 171-bytes long, so the two tables together require only 427 bytes. That's 331 bytes or 43 percent less memory than our first table required.

Space saved in tables may not be worth it if large or complicated code is required as an index to those tables, but in this case the code is quite simple:
\(\left.$$
\begin{array}{lll}\text { MNEMON } & \begin{array}{l}\text { LDX \#3 } \\
\text { STX LETTER }\end{array} & \begin{array}{l}\text { There are three letters in a mnemonic. } \\
\text { We'll keep track of the letters by count- } \\
\text { ing down to zero. }\end{array}
$$ <br>
Prepare to use the opcode as an index. <br>
Look up the mnemonic code for that op- <br>
code. (MCODES is the table of <br>
mnemonic codes.) <br>
Prepare to use that mnemonic code as <br>

an index.\end{array}\right]\)| TAX | LDA MCODES, |
| :--- | :--- |

As you can see, MNEMON requires only 30 bytes of code in machine language: 2 bytes to hold variables and 427 bytes for the two tables (MNAMESS and MCODES). The entire subroutine requires 459 bytes, but since most of those bytes are data in tables, comparatively little can go wrong with the program. If the wrong bytes are keyed into the table of mnemonic names, then the disassembler will print one or more incorrect characters in a mnemonic. But MNEMON won't crash! Bad
data in means bad data out, but at least MNEMON will run, and a running program is a lot easier to correct than one that crashes and burns.

So again we have a data-intensive, rather than a computation-intensive, subroutine. The tables required by MNEMON are included in Appendix C8.

## Print Operand

Now we come to the tricky part: printing the right operand given an opcode at some location in memory. When I disassemble object code by hand, I write the operand in two steps: first I determine the addressing mode of the given opcode, and then, if that addressing mode takes an operand, I write down the proper operand in the proper form. Proper form means including a comma and an $X$ or a $Y$ for every indexed instruction, including parentheses in the proper places for indirect instructions, and printing out all addresses high byte first, since that makes it easier to read an address.

OPERND (the subroutine that prints an operand for a given opcode in a given location in memory) will therefore determine the addressing mode for a given opcode, and then call an appropriate subroutine to handle that addressing mode:

## OPERND

| OPERND | TAX | Look up addressing mode code for |
| :--- | :--- | :--- |
| LDA MODES,X | this opcode. <br> TAX |  |
|  | TAX now indicates the addressing mode. |  |

MODES is a table giving the addressing mode for each opcode.
Note that OPERND can work only if we have a routine called MODE.X which somehow transfers control to the subroutine that handles addressing mode "X." How can MODE.X do this? One way is to have a table of pointers, in which the Xth pointer points to the subroutine that handles addressing mode "X." MODE.X must then transfer control to the Xth subroutine in this table. It would be nice if the 6502 offered an indexed JSR instruction, which would call the subroutine whose address is the Xth entry in the table. Unfortunately, the 6502 doesn't offer an indexed JSR instruction, so we'll have to simulate one in software.

Fortunately, the 6502 does offer an indirect JMP. If a pointer, called SUBPTR, can be made to point to a given subroutine, then the instruction JMP (SUBPTR) will transfer control to that subroutine. Therefore, MODE.X need only set SUBPTR equal to the $X$ th pointer in a table of subroutine pointers, and with the instruction

JMP (SUBPTR), it can transfer control to the Xth subroutine in the table.

## HANDLE ADDRESSING MODE "X"

| MODE.X | LDA SUBS, X | Get low byte of Xth pointer in the table of subroutine pointers. |
| :---: | :---: | :---: |
|  | STA SUBPTR | Set low byte of subroutine pointer. |
|  | INX | Adjust index to get next byte. |
|  | LDA SUBS, X | Get high byte of Xth pointer in the table of subroutine pointers. |
|  | STA SUBPTR +1 | Set high byte of subroutine pointer. |
|  | JMP (SUBPTR) | Jump to the subroutine specified by the subroutine pointer. That subroutine will then return to the caller of MODE.X, not to MODE.X itself. |
| SUBS |  | This is a table of pointers, in which the Xth pointer points to the subroutine tha handles addressing mode X . |

## Disassembler Utilities

Given MODE.X, OPERND can call the right subroutine to handle any give addressing mode. Now all we need are thirteen different subroutines, one for each $c$ the 6502's different addressing modes.

Before writing those subroutines, however, let's think for a moment about whe they must do, and see if we can't write a few utility subroutines to perform thos functions. With a proper set of utilities, the addressing mode subroutines themselvt need only call the right utilities in the right order.

The following set of utilities seems reasonable:

- ONEBYT:
Print a 1-byte operand.
- TWOBYT:
Print a 2-byte operand.
- RPAREN:
- LPAREN:
- XINDEX:
- YINDEX:
Print a right parenthesis.
Print a left parenthesis.
Print a comma and then the letter "X."
Print a comma and then the letter "Y."

| ONEBYT | JSR INC.SL | Advance to byte following opcode. |
| :--- | :--- | :--- |
|  | JSR DUMPSL | Print it in hexadecimal. |
|  | RTS | Return to caller. |

## Print a 2-Byte Operand: TWOBYT

A 2-byte operand always specifies an address with the low byte first. To print a 2-byte operand high byte first, we must first print the second byte in the operand and then print the first byte in the operand; each, of course, in hexadecimal format.

| TWOBYT | JSR INC.SL | Advance to first byte of operand. |
| :--- | :--- | :--- |
|  | LDA GET.SL | Load that byte into accumulator. |
|  | PHA | Save it. |

ONEBYT and TWOBYT each leave SELECT pointing at the last byte of the operand.

## Print Right, Left Parenthesis: RPAREN, LPAREN

RPAREN prints a right parenthesis to all currently selected devices. LPAREN prints a left parenthesis to all currently selected devices.

| RPAREN | LDA \#') | Load accumulator with ASCII code for <br> right parenthesis. |
| :--- | :--- | :--- |
|  | BNE SENDIT | Send it to all currently selected devices. <br> LPad accumulator with ASCII code for |
| LPAREN | LDA \#'( | Left parenthesis. |
| SENDIT | JSR PR.CHR | Send it to all currently selected devices. <br> Return to caller. |

## Index with Register X: XINDEX

XINDEX prints a comma and then the letter " $\mathrm{X}:$ "

| XINDEX | LDA \#', | Load accumulator with ASCII code for a <br> comma; then print it to |
| :--- | :--- | :--- |
|  | JSR PR.CHR | all currently selected devices. |

## Index with Register Y: YINDEX

YINDEX prints a comma and then the letter "Y:"

| YINDEX | LDA \#', | Load accumulator with ASCII code for a <br> comma; then print it to all |
| :--- | :--- | :--- |
|  | JSR PR.CHR | currently selected devices. |

So much for the disassembler utilities. Now with a single subroutine call we can print a 1-byte or a 2-byte operand (and, of course, we can print a no-byte operand), and we can print any of the frequently used characters and character combinations. Okay, let's write some addressing mode subroutines:

## Addressing Mode Subroutines

Because the 6502 has thirteen different addressing modes, we'll need thirteen different addressing mode subroutines:

Subroutine Addressing Mode
ABSLUT Absolute

| ABS.X | Absolute, $X$ |
| :--- | :--- |
| ABS. $Y$ | Absolute, $Y$ |
| ACC | Accumulator |
| IMPLID | Implied |
| IMMEDT | Immediate |
| INDRCT | Indirect |
| IND.X | Indirect,X |
| IND.Y | Indirect,Y |
| RELATV | Relative |
| ZEROPG | Zero Page |
| ZERO.X | Zero Page,X |
| ZERO.Y | Zero Page,Y |

The main job for each subroutine will be to print the operand in the proper form. Although a given addressing mode will always have the same number of characters in its operand, unfortunately, different addressing modes may have operands of different lengths. For example, implied addressing mode has no characters in its operand, whereas indirect indexed addressing requires eight characters in its operand, if leading zeros are included.

But no matter how many characters appear in an operand, we want to make sure that field 3 (the address field) always begins at the same column. Therefore, every addressing-mode subroutine will return with A holding the number of characters in the operand, with $X$ holding the number of bytes in the operand, and with SELECT pointing at the last byte in the operand (or at the opcode, if it was a 1-byte instruction). Then FINISH can print an appropriate number of spaces before printing fields 3 thru 6.

## Absolute Mode: ABSLUT

To print the operand for an instruction in the absolute mode, we need only print a 2-byte operand. Thus, 8D B2 04 will disassemble as:

$$
\text { STA 04B2 8D B2 } 04
$$

| ABSLUT | JSR TWOBYT |  |
| :--- | :--- | :--- |
|  | LDX \#2 | X holds number of bytes in operand. |
|  | LDA \#4 | A holds number of characters in <br> operand. |

## Absolute, X Mode: ABS.X

To print the operand for an instruction in the absolute, X mode, we must print a 2-byte operand, a comma, and then an " $X$ :"

## LDA D09A,X BD 9A D0

ABS.X

JSR ABSLUT JSR XINDEX LDX \#2
LDA \#6
RTS

Print the 2-byte operand.
Print the comma and the "X." $X$ holds number of bytes in operand.
A holds number of characters in operand.
Return to caller.

## Abolute, Y Mode: ABS.Y

To print the operand for an instruction in the absolute, Y mode, we must print a 2-byte operand, a comma, and then a " Y :"

## ORA 02FE,Y 19 FE 02

| ABS.Y | JSR ABSLUT | Print the 2-byte operand. |
| :--- | :--- | :--- |
|  | JSR YINDEX | Print the comma and the "Y." |
|  | LDX \#2 | X holds number of bytes in operand. |
|  | LDA \#6 | A holds number of characters in |
|  | RTS | operand. |
|  | Return to caller. |  |

Accumulator Mode: ACC
To print the operand for an instruction in the accumulator mode, we need only print the letter " A :"

| ACC | LDA \#'A | Load accumulator with ASCII code for <br> the lette A. |
| :--- | :--- | :--- |
|  | JSR PR.CHR | Print it on all currently selected devices. |
| LDX \#0 | X holds number of bytes in operand. |  |
| LDA \#1 | A holds number of characters in |  |
|  | operand. |  |
| RTS | Return to caller. |  |

## Implied Mode: IMPLID

Implied mode has no operand, so just return:

## CLC 18

| IMPLID | LDX \#0 | X holds number of bytes in operand. |
| :--- | :--- | :--- |
|  | LDA \#0 | A holds number of characters in <br> operand. |

## Immediate Mode: IMMEDT

Immediate mode requires a 1-byte operand, which we'll print in hexadecimal format. Thus, it should disassemble the two consecutive bytes "A9 41" as follows:

LDA \#\$41 A9 41

| IMMEDT | LDA \#'\# | Print a '\#' sign. |
| :--- | :--- | :--- |
|  | JSR PR.CHR |  |
|  | LDA \#'\$ | Print a dollar sign. |
|  | JSR PR.CHR |  |
|  | JSR ONEBYT | Print 1-byte operand in hexadecimal for- |
|  |  | mat. |
|  | LDX \#1 | X holds number of bytes in operand. |
|  | LDA \#4 | A holds number of characters in |
|  | operand. |  |
|  | RTS | Return to caller. |

## Indirect Mode: INDRCT

To print the operand for an instruction in the indirect mode, we need only print an absolute operand within parentheses. Thus, the three consecutive bytes "6C 00 04" will disassemble as:

JMP (0400) 6C 0004

| INDRCT | JSR LPAREN | Print left parenthesis. |
| :--- | :--- | :--- |
|  | JSR ABSEUT | Print the 2-byte operand. |
|  | JSR RPAREN | Print the right parenthesis. |
|  | LDX \#2 | X holds number of bytes in operand. |
|  | LDA \#6 | A holds number of characters in |
|  | RTS | operand. |
|  | Return to caller. |  |

## Indirect, X Mode: IND.X

To print the operand for an instruction in the indirect, X addressing mode, we need to print a left parenthesis, a zero-page address, a comma, the letter " X ," and then a right parenthesis. Thus, the two consecutive bytes "A1 3C" will disassemble as:

## LDA (3C,X) A1 3C

| IND.X | JSR LPAREN | Print a left parenthesis. |
| :---: | :---: | :---: |
|  | JSR ZERO.X | Print a zero-page address, a comma, and the letter "X." |
|  | JSR RPAREN | Print a right parenthesis. |
|  | LDX \#1 | X holds number of bytes in operand. |
|  | LDA \#8 | A holds number of characters in operand. |
|  | RTS | Return to caller. |

## Indirect, Y Mode: IND.Y

To print the operand for an instruction in the indirect, Y mode, we must print a left parenthesis, a zero-page address, a right parenthesis, a comma, and then the letter "Y." Thus, the two consecutive bytes "B1 AF" will disassemble as:

LDA (AF), Y B1 AF

IND.Y

JSR LPAREN
JSR ZEROPG
JSR RPAREN
JSR YINDEX
LDX \#1
LDA \#8
RTS

Print a left parenthesis.
Print a zero-page address.
Print a right parenthesis.
Print a comma and then the letter " Y ."
$X$ holds number of bytes in operand.
A holds number of characters in operand.
Return to caller.

## Relative Mode: RELATV

Relative mode can be tricky. A relative branch instruction specifies a forward branch if its operand is plus (in the range of 00 to $\$ 7 \mathrm{~F}$ ), but it specifies a backward branch if its operand is minus (in the range of $\$ 80$ to $\$ F F$ ). Therefore, in order to determine the address specified by a relative branch instruction, we must first determine whether the operand is plus or minus, so we can determine whether we're branching forward or backward. Then we must add or subtract the least-significant 7 bits of the operand to or from the address immediately following the operand of the branch instruction; the result of that calculation will be the actual address specified by the branch instruction.

| RELATV | JSR INC.SL | Select next byte in memory. |
| :--- | :--- | :--- |
|  | JSR PUSHSL | Save SELECT pointer on stack. |
| JSR GET.SL | Get operand byte. |  |
|  | PHA | Save it on the stack. |
|  | JSR INC.SL | Increment SELECT pointer so it points <br> to the opcode following the relative |
|  | branch instruction. (Relative branches <br> are relative to the next opcode.) |  |
|  |  | RLA |
| PLA | Is it plus operand byte to accumulator. |  |


|  | BPL FORWRD | If plus, it means a forward branch. Since operand byte is minus, we'll be branching backward. |
| :---: | :---: | :---: |
|  | DEC SELECT +1 | Branching backward is like branching forward from a location 256 bytes lower in memory. |
| FORWRD | CLC | Add operand byte to the address |
|  | ADC SELECT | of the opcode following the |
|  | BCC RELEND | branch instruction. |
|  | INC SELECT+1 |  |
| RELEND | STA SELECT | Now SELECT points to the address specified by the operand of the relative branch instruction. Let's print it. |
|  | JSR PR.ADR |  |
|  | JSR POP.SL | Restore SELECT pointer. |
|  | LDX \#1 | $X$ holds number of bytes in operand. |
|  | LDA \#4 | A holds number of characters in operand. |
|  | RTS | Return to caller, with SELECT pointer once again pointing to the operand byte of the relative branch instruction. |

## Zero-Page Mode: ZEROPG

To print the operand of an instruction that uses the zero-page addressing mode, we could simply print a 1-byte operand. But I find listings more readable when all zero-page addresses are shown with the leading zeros (eg: " 00 FE " rather than " FE " to represent address $\$ 00 \mathrm{FE}$ ). Therefore, let's print all zero-page operands with a leading zero. That simply requires us to print two ASCII zeros and then to print the 1-byte operand. This will cause the bytes " 852 A " to be disassembled as:

STA 002A 852 A

ZEROPG LDA \#O
JSR PR.BYT
JSR ONEBYT
LDX \#1
LDA \#4
RTS

Print two ASCII zeroes to all currently selected devices. Print the 1-byte operand. X holds number of bytes in operand. A holds number of characters in operand. Return to caller.

## Zero-Page Indexed Modes: ZERO.X, ZERO.Y

To print the operand of an instruction that uses the zero-page $X$ or zero-page $Y$ addressing mode, we need only print the zero-page address, a comma, and then an " X " or a "Y." Thus, "B5 6C" will disassemble as:

> LDA 006C,X B5 6C
and "B6 53" will disassemble as:

LDX 0053,Y B6 53

| ZERO.X | JSR ZEROPG | Print the zero-page address. |
| :--- | :--- | :--- |
|  | JSR XINDEX | Print a comma and the letter "X." |
| LDX \#1 | X holds number of bytes in operand. |  |
|  | LDA \#6 | A holds number of characters in |
| ZERO.Y | operand. |  |

## A Pseudo-Addressing Mode for Embedded Text

Now we have subroutines to disassemble machine code in any of the 6502's thirteen legal addressing modes. But what about text embedded in a machinelanguage program? We know that our programs already include text strings, where each text string begins with a TEX character (\$7F) and ends with an ETX (\$FF). The disassembler, however, doesn't know anything about embedded text. If we try to disassemble a machine-language program that includes embedded text, the disassembler will assume that the TEX character, and the text string itself, are 6502 opcodes and operands; because it doesn't know about text, it will misinterpret the text string.

Wouldn't it be nice if the disassembler could recognize the TEX character for what it is, and then print out the text string as text, rather than as opcodes and operands? When it has finished printing a text string, the disassembler could then
resume treating the bytes following the ETX as conventional 6502 opcodes and operands.

Such behavior is not hard to implement. We need only define a pseudoaddressing mode, called TEXT mode, and say that the TEX character is the only opcode that has the TEXT addressing mode. Then we'll write a special addressing mode subroutine, called TXMODE, to print operands that are in the TEXT mode. TXMODE will print an operand in the TEXT mode by printing the text that follows the TEX character and ends with the first ETX character.

Here's some source code to implement such behavior:

| TXMODE | PLA | Pop return address |
| :--- | :--- | :--- |
|  | PLA | to OPERND. |
|  | PLA | Pop return address |
| TXLOOP | PLA | JSR NEXTSL |$\quad$| AdvLINE. |
| :--- | :--- |

Now that we have the desired addressing mode subroutines, we can make up the table of addressing mode subroutines:

SUBS
.WORD ABSLUT
.WORD ABS.X
.WORD ABS.Y
.WORD ACC
.WORD IMPLID
.WORD IMMEDT
.WORD INDRCT

> .WORD IND.X
> . WORD IND.Y
> .WORD RELATV
> .WORD ZEROPG
> WORD ZERO.X
> .WORD ZER.Y

Each addressing mode subroutine will return with SELECT pointing at the last byte in the instruction, with A holding the number of characters in the operand field, and with $X$ holding the number of bytes in the operand ( 0,1, or 2 ). Each addressing mode subroutine will return to OPERND, which will finish the line by calling FINISH.

Finishing the Line: FINISH
FINISH must space over to the proper column for field 3, which will hold the address of the opcode. Then it must print the address of the opcode and dump 1,2 or 3 bytes, as necessary. FINISH will end by advancing the printhead to a new line and by advancing SELECT so that it points to the first byte following the disassembled line (unless it has disassembled through EA, the ending address, in which case it will return with SELECT $=$ EA). FINISH returns PLUS if more bytes must be disassembled before EA is reached; it returns MINUS if it disassembled through EA.

| FINISH | STA OPCHRS STX OPBYTS | Save the length of the operand, in characters and in bytes. |
| :---: | :---: | :---: |
|  | DEX | If necessary, decrement the |
|  | BMI SEL.OK | SELECT pointer so it |
| LOOP. 1 | JSR DEC.SL | points to the opcode. |
|  | DEX |  |
|  | BPL LOOP. 1 |  |
| SEL.OK | SEC | Space over to the |
|  | LDA ADRCOL | column for the address field: |
|  | SBC \#4 | Operand field started in column 4... |
|  | SBC OPCHRS | .... and includes OPCHRS characters. |
|  | TAX | So now we need $X$ spaces. |
|  | JSR SPACES | Send enough spaces to reach address column. |
|  | JSR PR.ADR | Print address of opcode. |
| LOOP. 2 | JSR SPACE | Space once. |
|  | JSR DUMPSL | Dump selected byte. |
|  | JSR INC.SL | Select next byte. |


|  | DEC OPBYTS | Completed last byte in instruction? |
| :--- | :--- | :--- |
|  | BPL LOOP. 2 | If not, do next byte. <br> Back up SELECT to last byte in |
|  | JSR DEC.SL | operand. <br> FINEND |
|  | JSR CR.LF | Advance to a new line. |
| OPBYTS | RTS | Return to caller. |
| OPCHRS | .BYTE | Number of bytes in operand. |
| ADRCOL | .BYTE 16 | Number of characters in operand. |
| .BYTE | Starting column for address field. |  |

Now we can disassemble a line. So let's write the disassemblers, one for the printer and one for the screen. These routines will have much the same structure as TVDUMP and PRDUMP, which direct hexdumps to the printer or to the screen.

## Disassemble to Screen: TV.DIS

| TV.DIS | LDA DISLNS | Initialize line counter with |
| :---: | :---: | :---: |
|  | STA LINUM | number of lines to be disassembled. |
|  | LDA \#\$FF | Set end address to \$FFFF, |
|  | STA EA | so NEXTSL will always increment |
|  | STA EA+1 | the SELECT pointer. |
|  | JSR TVT.ON | Select TVT as an output device. (Other selected devices will echo the disassembly.) |
| TVLOOP | JSR DSLINE | Disassemble one line. |
|  | DEC LINUM | Completed last line yet? |
|  | BNE TVLOOP | If not, disassemble next line. |
|  | RTS | If so, return. |
| DISLNS | .BYTE 5 | DISLNS holds number of lines to be disassembled by TV.DIS. To disassemble one line set DISLNS $=1$ |
| LINUM | .BYTE 0 | This variable keeps track of the number of lines yet to be disassembled. |

## Printing Disassembler: PR.DIS

The printing disassembler (PR.DIS) will announce itself by displaying "PRINTING DISASSEMBLER" on the screen, but not on the printer. It will then let the user set the starting and ending addresses, in the same manner as PRDUMP. When the user has specified the block of memory to be disassembled, the PR.DIS will print a disassembly of the specified block of memory, echoing its output to the screen.

| PR.DIS | JSR PR.OFF | Deselect printer. |
| :--- | :--- | :--- |
|  | JSR TVT.ON | Select TVT. |
|  | JSR PRINT: | Display title: |
|  | .BYTE TEX |  |
|  | .BYTE CR,LF |  |
|  | .BYTE | PRINTING DISASSEMBLER' |
|  | .BYTE CR,LF,ETX |  |
|  | JSR.SETADS | Let user set starting address |
|  |  | and end address. |
|  | JSR GOTOSA | Set SELECT = Start address. |
| PRLOOP | JSR PR.ON | Select the printer. |
|  | JSR DSLINE | Disassemble one line. |
|  | BPL PRLOOP | If it wasn't the last line, disassemble the |
|  |  | next one. |
|  | RTS | Return to caller. |

With PR.DIS and TV.DIS, you can disassemble any block of memory, directing the disassembly to the screen or to the printer. See Chapter 12 for guidance on mapping these two disassemblers to function keys in the Visible Monitor.

## Chapter 10:

## A General MOVE Utility

Many computer programs spend a lot of time moving things from one place to another. Such programs should be able to call a move utility for most of this work. A move utility should:

- Be general enough to move anything of any size from any place in memory to anywhere else.
- Not be upset when the origin block overlaps the destination.
- Have entry points with input configurations convenient to different callers.
- Preserve its inputs.
- Be fast.

This routine will be called often. A calling program doesn't want to spend all its time here. The cost of that speed is size, because we'll use straight-line, dedicated code to handle each of several special cases, but even so this move code will weigh in at less than 200 bytes. That's less than three percent of the memory available on a system with 8 K bytes of programmable memory.

## Input Configurations

Different callers may find different input configurations convenient, so let's provide more than one entry point, each requiring different parameters to be set. The following two subroutine entry points are likely to meet the needs of most callers:

MOV.EA Move a block, defined by its starting address (SA), its ending
address (EA), and its destination address (DEST). Move a block, defined by its starting address, the number of bytes in the block (NUM), and the destination of the block.

MOV.EA will simply be a "front end" for MOVNUM. It will set NUM = ending address - starting address of the source block.

## Handling Overlap

There will be no problem with overlap if we always move from the leading edge of the source block - that is, copy up beginning with the highest byte to be moved, and copy down beginning with the lowest byte to be moved. This way, if a byte in the source block is overwritten it will already have been copied to its destination.

## Going Up?

To avoid overlap, MOVNUM must determine whether it's copying up or down. Therefore, before moving anything it must see if the destination address is greater or lesser than the starting address. Then it can branch to MOVE-UP or MOVE-DOWN as appropriate.

Figure 10.1: Top level of block move. Flowchart of MOVE.EA and MOVNUM routines.


Using the flowchart of figure 10.1 as a guide, let's write source code for the top level of MOV.EA and MOVNUM:

| MOV.EA | $\begin{aligned} \text { GETPTR } & =0 \\ \text { PUTPTR } & =\text { GETPTR }+2 \end{aligned}$ | This is the input-page pointer. This is the output-page pointer. |
| :---: | :---: | :---: |
|  | SEC | Set NUM = EA - SA |
|  | LDX EA +1 |  |
|  | LDA EA |  |
|  | SBC SA |  |
|  | STA NUM |  |
|  | BCS MOVE. 1 |  |
|  | DEX |  |
|  | SEC |  |
| MOVE. 1 | TXA |  |
|  | SBC SA+1 |  |
|  | STA NUM +1 |  |
|  | BCS MOVNUM | Now NUM $=\mathrm{EA}-\mathrm{SA}$. |
| ER.RTN | LDA \#ERROR | If EA less than SA, |
|  | RTS | return with error code. |
| MOVNUM <br> SAVE | LDY \#3 | Save the 4 zero-page |
|  | LDA GETPTR, Y | bytes we'll use. |
|  | PHA |  |
|  | DEY |  |
|  | BPL SAVE |  |
|  | SEC | Is DEST less than START? |
|  | LDA SA +1 |  |
|  | CMP DEST+1 |  |
|  | BCC MOVEUP | If so, we'll move down. |
|  | BNE MOVEDN | If not, we'll move up. |
|  | LDA SA | SA , destination are in the same page. |
|  | CMP DEST | If SA more than destination, we'll |
|  | BCC MOVEUP | move down. If SA less than destination, |
|  | BNE MOVEDN | we'll move up. If they are equal, we'll return bearing okay code. |
| OK.RTN RESTOR | LDY \#0 | Restore 4 zero-page bytes that were |
|  | PLA | used by the move code. |
|  | STA GETPTR, $Y$ |  |
|  | INY |  |
|  | CPY \#4 | Restored last byte yet? |
|  | BNE RESTOR | If not, restore next one. If so, |
|  | RTS | return, with move complete and zero page preserved. |
| NUM | .WORD 0 | This 16-bit variable holds the number of bytes to be moved. |

## Optimizing for Speed

Moving a page at a time is the fastest way to move data, and for large blocks we can move most of the bytes this way. Therefore, when moving data we'll move one page at a time until there is less than a page to move; then we'll move a byte at a time until the entire source block is moved. MOVE-UP and MOVE-DOWN must test to see if they have more or less than a page to move, and then branch to dedicated code that either moves a page or moves less than a page.

Figure 10.2: Move a block up. Flowchart of the MOVEUP routine.


## MOVE-UP

Using figure 10.2 as a guide, we can write source code for MOVE-UP:
$\left.\begin{array}{ll}\text { MOVEUP } & \begin{array}{l}\text { LDA NUM +1 } \\ \text { BEQ LESSUP }\end{array} \\ & \begin{array}{l}\text { More than one page to move? } \\ \text { If not, move less than a page up. } \\ \text { To move more than a page, set the page } \\ \text { pointers GETPTR and PUTPTR to the }\end{array} \\ \text { highest pages in the source and destina- } \\ \text { tion blocks. To do this, treat } \mathrm{X} \text { as the } \\ \text { high byte and Y as the low byte of a } \\ \text { pointer, which we'll call (X,Y). First set } \\ \text { (X,Y) = NUM - \$FF, the relative ad- } \\ \text { dress of the highest page in the block. }\end{array}\right\}$

| PAGEUP | LDX NUM +1 LDY \#SFF | Load $X$ with number of pages to move Move a page up. |
| :---: | :---: | :---: |
| UPLOOP | LDA (GETPTR), Y | Get a byte from origin block. |
|  | STA (PUTPTR), Y | Put it in destination block. |
|  | DEY | Adjust index for next byte down. |
|  | BNE UPLOOP | Loop if not the last byte. |
|  | LDA (GETPTR), $Y$ | Move last byte. |
|  | STA (PUTPTR), Y |  |
|  | DEC GETPTR +1 | Decrement page pointers. |
|  | DEC PUTPTR+1 |  |
|  | DEX | Still more than a page to move? |
|  | BNE PAGEUP | If so, move up another page. |
| LESSUP | JSR LOPAGE | Set GETPTR, PUTPTR to bottom of origin and destination blocks. |
|  | LDY NUM | Set index to number of bytes to be moved. |
| SOMEUP | LDA (GETPTR), Y STA (PUTPTR), Y | Move a byte. |
|  | DEY | About to move last byte? |
|  | CPY \#\$FF |  |
|  | BNE SOMEUP | If not, move another. |
|  | JMP OK.RTN | If so, return bearing "OK" code. |
| LOPAGE | LDA SA | Set page pointers to the bottom |
|  | STA GETPTR | of the origin and destination |
|  | LDA SA+1 | blocks. |
|  | STA GETPTR +1 |  |
|  | LDA DEST |  |
|  | STA PUTPTR |  |
|  | LDA DEST+1 |  |
|  | STA PUTPTR+1 |  |
|  | RTS | Return to caller. |

## Move-Down: MOVEDN

Figure 10.3 shows an algorithm for moving a block of data down through memory.


Using figure 10.3 as a guide, we can write source code for the move-down routine:

| MOVEDN | JSR LOPAGE | Set page pointers to bottom of origin <br> and destination blocks. |
| :--- | :--- | :--- |
|  | LDY \#0 | Y must equal zero whether we move <br> more or less than a page. |
|  | LDX NUM+1 | More than one page to move? <br> If not, move less than a page down. |
|  | BEQ LESSDN | Move a page down. <br> PAGEDN |
| LDA (GETPTR),Y | Get a byte from origin block <br> and put it in destination block. |  |
|  | STA (PUTPTR),Y | Moved last byte in page? |


|  | BNE PAGEDN |  |
| :---: | :---: | :---: |
|  | INC GETPTR + 1 | Increment page pointers. |
|  | INC PUTPTR+1 |  |
|  | DEX | Still more than a page to move? |
|  | BNE PAGEDN | If so, move another page down. |
|  | LDY \#0 | Move less than a page down starting at the bottom. |
| LESSDN | LDA (GETPTR), Y | Get a byte from origin... |
|  | STA (PUTPTR), Y | and put it in destination block. |
|  | INY | Adjust index for next byte. |
|  | SEC |  |
|  | CPY NUM | Moved last byte yet? |
|  | BCC LESSDN | If not, move another. |
|  | JMP OK.RTN | If so, return to caller, bearing " $\mathrm{OK}^{\prime}$ code. |

INC GETPTR + 1
INC PUTPTR + 1 DEX
BNE PAGEDN
LDY \#O
LESSDN
LDA (GETPTR), Y
STA (PUTPTR), Y
INY
SEC
CPY NUM
BCC LESSDN
JMP OK.RTN

Increment page pointers.
Still more than a page to move?
If so, move another page down.
Move less than a page down starting at the bottom.
Get a byte from origin... and put it in destination block. Adjust index for next byte.

Moved last byte yet?
If not, move another. If so, return to caller, bearing " $\mathrm{OK}^{\prime}$ code.

## Speed

For large blocks of data, most bytes will be moved by the page-moving code: PAGE-UP and PAGE-DOWN. Since the processor spends most of its time in these loops, let's see how long they will take to move a byte. (Appendix A5, Instruction Execution Times, provides information on the number of cycles required for each 6502 operation.) Ordinarily I would not go into great detail concerning the speed of execution of a small block of code, but these two loops form the heart of the move utility, because they move most of the bytes in any large block. By making those two loops very efficient, we can make the move utility very fast. In fact, these loops will let us move blocks bigger than one page, at a rate approaching 16 cycles/byte moved. (By way of a benchmark, that's more than twice as fast as the time required to move large blocks with MOVIT, a smaller move program published in The First Book of KIM. * MOVIT, made tiny [ 95 bytes] to use as little as possible of the KIM's limited programmable memory, requires at least 33 cycles/bytes moved.)

MOVE.EA and MOVNUM are move utilities because they have input configurations and performance suitable for many calling programs. But they are not very convenient to the human user who simply wants to move something. With the Visible Monitor and the move utility, you can move something from one place to

[^0]another, but you have to know what addresses to set and you have to know the address of the move utility itself.

That's too much for me to remember. I want a tool, which will know the addresses and won't require me to remember them.

When I'm developing programs with the Visible Monitor and I want to move some data or code from one place to another, I'd like to be able to call up a move tool with a single keystroke - say "M." It's easier for me to remember " M ' for Move" than it is to remember the address of the move utility and the addresses of its inputs.

Let's say I'm using the Visible Monitor and I press " M ." This invokes the move tool. The first thing it should do is let me know that it's active. What if I hit the " M " key by mistake? The computer should let me know that I've invoked a new program.

It should put up a title: "MOVE TOOL." Then it should let me specify the start, end, and destination addresses of a given block in memory. When these addresses are set, the move tool can call MOV.EA, which will actually perform the move, based on the addresses set by the user.

The top level of the move tool is therefore quite simple. Figure 10.4 shows the flowchart for the following routine:


Figure 10.4: A move tool. Flowchart of MOVER routine.

## MOVER

| MOVER | JSR TVT.ON | Select screen as an output device. <br> JSR PRINT: |
| :--- | :--- | :--- |
|  | Put a title on the screen. |  |

Of course, MOVER can work only if we have a routine that lets the user set the destination address. Let's write such a routine, and we'll be all set to move whatever we like, to wherever we want it.

## Set Destination Address: SET.DA

SET.DA JSR TVT.ON

JSR PRINT:
.BYTE TEX
.BYTE CR,LF,LF
.BYTE
.BYTE
.BYTE ETX JSR VISMON

DAHERE LDA SELECT STA DEST LDA SELECT +1
STA DEST +1
RTS
DEST .WORD 0

Select TVT as an output device. All other selected output devices will echo the screen output. Put prompt on the screen:
"SET DESTINATION ADDRESS " "AND PRESS Q."

Call the Visible Monitor, so user can specify a given address. Set destination address equal to address set by the user.

Return to caller.
Pointer to destination of block to be moved.

See Chapter 12, Extending the Visible Monitor, to learn how to hook the move tool into the Visible Monitor by mapping it to a given key. Then to move anything in memory to anywhere else, you need only strike that key and the move tool will do the rest.

## Chapter II:

## A Simple Text Editor


#### Abstract

With the Visible Monitor you can enter ASCII text into memory by placing the arrow under field 2 and striking character keys. But you must strike two keys for every character in the message: first the character key, to enter the character into the displayed address, and then the space bar, to select the next address. Furthermore, if you want to enter an ASCII space or carriage return into memory, you'll have to place an arrow under field 1 and enter the hexadecimal representation of the desired character: $\$ 20$ for a space; $\$ 0 \mathrm{D}$ for a carriage return. Then, of course, you'll have to hit the space bar to select the next address, and the "greater than" key to move the arrow back underneath field 2 , so that you can enter the next character into memory.

If you only need to enter up to a dozen ASCII characters at a time, then the Visible Monitor should meet your needs. When you need to enter longer messages into memory, you'll find yourself wanting a more suitable tool - a simple text editor.

Text editors come in many different shapes, sizes and formats. A line-oriented editor, suitable for creating and editing program source files, requires that you enter and edit text a line at a time. Usually each line must be numbered when it is entered; then, in order to edit a line, you must first specify it by its line number.

On the other hand, a character-oriented editor allows you to overstrike, insert, or delete characters anywhere in a given string of characters. Character-oriented editors are frequently found in word processors for office applications, but don't get your hopes up; this chapter will not present software nearly as sophisticated as that available in even the humblest of word processors. However, it will present a very simple character-oriented editor that will enable you to enter and edit text strings, such as prompts, anywhere in memory.


## Structure

The text editor will have the three-part structure shown in figure 11.1. From this we can write source code for the top level of the text editor:


Figure 11.1: Structure of simple text editor.

| EDITOR | JSR SETBUF | Initialize pointers and variables required <br> by the editor. <br> Show the user a portion of the text <br> buffer. |
| :--- | :--- | :--- |
| EDLOOP | JSR SHOWIT | Let the user edit the buffer or move <br> about within it. |
|  | JSR EDITIT | CLC |
|  | BCC EDLOOP | Loop back to show the current text. |

Look familiar? It should. This is essentially the same structure used in the Visible Monitor. It's a simple structure, well-suited to the needs of many interactive display programs.

## SETBUF

The text editor will operate on text in a portion of memory called the text buffer. Because the editor must be able to change the contents of the text buffer, the buffer must occupy programmable memory and may not be used for any other purpose. This exemplifies a problem familiar to programmers: how to allocate memory in the most effective manner. Memory used to store a program cannot be used at the same time to store text; nor can memory allotted to the text buffer be used for stor-
ing programs or variables.
How do you get five pounds of tomatoes into a four-pound-capacity sack without crushing the tomatoes or tearing the sack? You don't. If you want to store a lot of text in your computer's programmable memory, you might not have room for much of a text editor. On the other hand, an elaborate text editor, requiring a good deal of programmable memory for its own code, may not leave much room in your system for storing text.

Therefore, this text editor leaves the allocation of memory for the text buffer to the discretion of the user. A subroutine called SETBUF sets pointers to the starting and ending addresses of the text buffer. The rest of the editor then operates on the text buffer defined by those pointers.

SETBUF sets the starting and ending addresses of the edit buffer. If you always want to enter and edit text in the same buffer, then substitute your own subroutine to set the starting and ending addresses to the values you desire. Otherwise, use the following version of SETBUF, which lets the user define a new text buffer each time it is called.

For testing purposes, you might even want to set the text buffer completely inside screen memory. This allows you to see exactly what's happening inside the text buffer.

## SETBUF

| SETBUF | JSR TVT.ON | Select TVT. |
| :--- | :--- | :--- |
|  | JSR PRINT: | Display "SET UP EDIT BUFFER." |
|  | .BYTE TEX,CR,LF,LF |  |
|  | .BYTE 'SET UP EDIT BUFFER' |  |
| GETADS | .BYTE CR,LF,LF,ETX |  |
|  |  | Let user set starting address and end ad- <br>  <br>  <br>  <br>  <br>  <br>  <br>  <br>  <br> RSR GOTS |
|  |  | dress of edit buffer. |
|  | Now SELECT = starting address of edit |  |

This version of SETBUF allows the user to set the text buffer anywhere in memory, provided that the ending address is not lower in memory than the starting address. It returns with the SELECT pointer pointing at the starting address of the buffer.

## SHOWIT

Now that SETBUF has set the pointers associated with the text buffer, let's figure out how to display part of that buffer.

Figure 11.2 shows the simple 3 -line display to be used by the text editor. " X " marks the home position of the edit display. Everything in the edit display is relative to the home position. Thus, to move the edit display about on your screen (ie: from the top of the screen to the bottom of the screen), you need only change the home position, which is set by SHOWIT.

LINE 1:
LINE 2:

LINE 3:


Figure 11.2: Three-line display of simple text editor.

Line 1 is entirely blank. Its only purpose is to separate the text displayed in line 2 from whatever you may have above it on your screen.

Line 2 displays a string of characters from the edit buffer. The central character in line 2 is the current character. The current character is indicated by an upwardpointing arrow as in line 3 . The address of the current character is given by the four hexadecimal characters represented by " HHHH " in line 3.

The letter " M " in line 3 shows you where a graphic character will indicate the current mode of the editor.

## Modes

This editor will have two modes: overstrike mode and insert mode. In overstrike mode you overstrike, or replace, the current character with the character from the keyboard. In insert mode, you insert the keyboard character into the text buffer just before the current character. How one sets these modes, a function for the subroutine EDITIT, will be discussed later. But SHOWIT must know the current mode in order to display the proper graphic in line 3 of the editor display.

Since we're going to have two modes, let's keep track of the current mode of the editor with a 1-byte variable called EDMODE. We'll assign the following values to EDMODE:

EDMODE $=0$ when the editor is in overstrike mode. EDMODE $=1$ when the editor is in insert mode.

Any other value of EDMODE is undefined and therefore illegal. If SHOWIT should find that EDMODE has an illegal value, then it should set EDMODE to some legal default value - say, zero. That would make overstrike the default mode for the editor.

We'll also need two graphics characters, INSCHR and OVRCHR, to indicate insert and overstrike modes, respectively. In this chapter, the character to indicate a given edit mode will simply be the first initial of the mode name: " 0 " for overstrike mode, "I" for insert mode.

## SHOWIT

SHOWIT JSR TVPUSH
JSR TVHOME
LDX TVCOLS
LDY \#3
JSR CLR.XY
JSR TVHOME
JSR TVDOWN
JSR TVPUSH
JSR LINE. 2
JSR TV.POP
JSR TVDOWN
JSR LINE. 3
JSR TV.POP RTS

Save the zero-page bytes we'll use.
Set home position of the edit display.
Clear 3 rows for the edit display.

Restore TV.PTR to home position of edit display.
Set TV.PTR to beginning of
line 2 and save it.
Display text in line 2.
Set TV.PTR to beginning of line 3 .
Display line 3.
Restore zero-page bytes used.
Return to caller, with edit display on screen, rest of screen unchanged, and zero page preserved.

Of course, SHOWIT can work only if it can call a couple of routines (LINE. 2 and LINE.3) to display lines 2 and 3 of the editor display, respectively. Let's write those routines.

## Display Text Line

To display the text line, we simply need to copy a number of characters from the text buffer to the second line of the editor display. Since the screen is TVCOLS wide, we should display TVCOLS number of characters in such a way that the central character in the display is the currently selected character. We can do that if we decrement SELECT by TVCOLS/2 times, and then display TVCOLS number of characters:

| LINE. 2 | JSR PUSHSL | Save SELECT pointer. |
| :--- | :--- | :--- |
|  | LDA TVCOLS | Set X equal |
| to half the width |  |  |
|  | LSR A | of the screen. |
|  | TAX |  |
|  | DEX |  |
|  | DEX |  |
|  | JSR DEC.SL | Decrement SELECT X times. |
|  | DEX |  |
|  | BPL LOOP. 1 |  |
|  | LDA TVCOLS | Initialize COUNTR. (We're |
|  | STA COUNTR | going to display TVCOLS characters.) |
|  | JSR GET.SL | Get a character from buffer. |
|  | JSR TV.PUT | Put it on screen. |
|  | JSR IVSKIP | Go to next screen position. |
|  | DEC COUNTR | Advance to next byte in buffer. |
|  | BPL LOOP. 2 | Done last character in row? |
|  | JSR POP.SL | If not, do next character. |
|  | RTS | Restore SELECT from stack. |
|  |  | Return to caller. |

## LINE. 2

## Display Status Line

Line 3 of the editor display provides status information: identifying the current mode of the editor, pointing at the current character in line 2 of the edit display, and providing the address of the current character.

## LINE. 3

| LINE. 3 | LDA TVCOLS |  |
| :---: | :---: | :---: |
|  | LSR A | A $=$ TVCOLS $/ 2$ |
|  | SBC \#2 | $\mathrm{A}=(\mathrm{TVCOLS} / 2)-2$ |
|  | JSR TVPLUS | Now TV.PTR is pointing 2 characters to the left of center of line 3 of the edit display. |
|  | LDA EDMODE | What is current mode? |
|  | CMP \#1 | Is it insert mode? |
|  | BNE OVMODE | If not, it must be overstrike mode. |
|  | LDA \#INSCHR | If so, load A with the insert graphic. |
|  | CLC |  |
|  | BCC TVMODE |  |
| OVMODE | LDA \#OVRCHR | Load A with the overstrike graphic. |
| TVMODE | JSR TV.PUT | Put mode graphic on screen. |
|  | LDA \#2 |  |
|  | JSR TVPLUS | Now TVPTR is pointing at the center of line 3 of the edit display. |
|  | LDA ARROW | Display an up-arrow here, |
|  | JSR TV.PUT | pointing up at the current character. |
|  | LDA \#2 |  |
|  | JSR TVPLUS | Now TV.PTR is pointing at the position reserved for the address of the current character. |
|  | LDA SELECT +1 | Display address of current |
|  | JSR VUBYTE | character. |
|  | LDA SELECT |  |
|  | JSR VUBYTE |  |
|  | RTS | Return to caller. |

We've chosen to define the editor's current character as the character pointed to by SELECT. We've already developed some subroutines that operate on the SELECT pointer and on the currently selected byte, so we won't have to write many new editor utilities; instead, we can use many of the SELECT utilities presented in earlier chapters.

## Edit Update

Now we can display the three lines of the edit display. What else must the editor do? Oh, yes: it must let us edit. Here's a reasonably useful, if small, set of editor functions:

- Allow the user to move forward through the message.
- Allow the user to move backward through the message.
- Allow the user to overstrike the current character.
- Allow the user to delete the current character.
- Allow the user to delete the entire message.
- Allow the user to insert a new character at the current character position.
- Allow the user to change modes from insert to overstrike and back again.
- Print the message.
- Allow the user to terminate editing, thus causing the editor to return to its caller.

What keys will perform these functions? I'll leave that up to you by treating the editor function keys as variables and keeping them in a table called EDKEYS (see Appendix C11). To assign a given function to a given key, store the character code generated by that key in the appropriate place in the table:

## EDITIT

ENDEDT PLA
PLA
PLA
RTS
NOTEND STA TEMPCH
PLA
JSR DO.KEY
LDA TEMPCH
JSR GETKEY CMP QUITKY
BNE NOTEND

| ENDEDT | PLA |
| :--- | :--- |
|  | PLA |
|  | PLA |
| NOTEND | RTS |
| STA TEMPCH |  |

## EDITIT

| DO.KEY | CMP MODEKY | Is it the "change mode" key? |
| :---: | :---: | :---: |
|  | BNE IFNEXT | If not, perform the next test. |
|  | DEC EDMODE | If so, change the editor's mode... |
|  | BPL DO.END | * |
|  | LDA \#1 |  |
|  | STA EDMODE |  |
| DO.END | RTS | and return. |
| IFNEXT | CMP NEXTKY | Is it the "next" key? |
|  | BNE IFPREV | If not, perform the next test. |
|  | JSR NEXTCH | If so, advance the current position by one character... |
|  | RTS | and return. |
| IFPREV | CMP PREVKY | Is it the "previous" key? |
|  | BNE IF.RUB | If not, perform the next test. |
|  | JSR PREVCH | If so, back up the current position by one character... |
|  | RTS | and return. |
| IF.RUB | CMP RUBKEY | Is it the "delete" key? |
|  | BNE IF.PRT | If not, perform the next test. |
|  | JSR DELETE | If so, delete the current character... |
|  | RTS | and return. |
| IF.PRT | CMP PRTKEY | Is it the "print" key? |
|  | BNE IFFLSH | If not, perform the next test. |
|  | JSR PRTBUF | If so, print the buffer... |
|  | RTS | and return. |
| IFFLSH | CMP FLSHKY | Is it the "flush" key? |
|  | BNE CHARKY | If not, perform the next test. |
|  | JSR FLUSH | If so, flush all text in the edit buffer... |
|  | RTS | and return. |
|  |  | OK. It's not an editor function key, so it must be a regular character key. Therefore, if we're in overstrike mode we'll overstrike the current character with the new character, and if we're in insert mode we'll insert the new character at the current character position. |
| CHARKY | LDX EDMODE | Are we in overstrike mode? |
|  | BEQ STRIKE | If so, overstrike the character. |
|  | JSR INSERT | If not, insert the character... |
|  | RTS | and return. |
| STRIKE | JSR PUT.SL | Put the character into the currently |


|  | JSR NEXTSL RTS | Advance to the next character position, and return to caller. |
| :---: | :---: | :---: |
| INSERT | PHA | Save the character to be inserted, while we make space for it in the edit buffer... |
|  | JSR PUSHSL | Push the address of the current character onto the stack. |
|  | LDA SA +1 | Push starting address of the buffer |
|  | PHA | onto stack. |
|  | LDA SA |  |
|  | PHA |  |
|  | LDA EA+1 | Push ending address of the buffer |
|  | PHA | onto stack. |
|  | LDA EA |  |
|  | PHA |  |
|  | JSR SAHERE | Set SA $=$ SELECT, so current character will be the start of the block we'll move. |
|  | JSR NEXTSL | Advance to next character position in the text buffer. |
|  | BMI ENDINS | If we're at the end of the buffer, we'll overstrike instead of inserting. |
|  | JSR DAHERE | Set DEST $=$ SELECT, so destination of block move will be 1 byte above block's start address (ie, we'll move a block up by 1 byte). |
|  | LDA EA | Decrement end address |
|  | BNE NEXT | so we won't move text |
|  | DEC EA +1 | beyond the end of |
| NEXT | DEC EA | the text buffer. |
|  |  | Now the starting address is the current character, the destination address is the next character, and the ending address is one character shy of the last character in the buffer. We're ready now to move a block. |
| OPENUP | JSR MOV.EA | Open up 1 byte of space at the current character's location, by moving to DEST the block specified by SA and EA. |
| ENDINS | PLA | Restore EA so it points to the last byte |
|  | STA EA | in the edit buffer. |
|  | PLA |  |
|  | STA EA +1 |  |
|  | PLA | Restore SA so it points to the first byte |
|  | STA SA | in the edit buffer. |

## PLA

STA SA+1
JSR POP.SL
PLA
Restore SELECT so it points to the current character.
Reload the accumulator with the character to be inserted. Since we've created a 1-byte space for this character, we need only overstrike it.
JSR STRIKE RTS

Return to caller.

EDITIT looks like it will do what we want it to do - provided that it may call the following (as yet unwritten) subroutines:

- NEXTCH - Select next character.
- PREVCH - Select previous character.
- FLUSH - Flush the buffer.
- PRTBUF - Print the buffer.

Let's write them.

## Select Next Character

We want to be able to advance through the text buffer, but we don't want to be able to go beyond the end of the buffer or beyond the end of the message. The end of the message will be indicated by one or more ETX (end-of-text) characters. ETX characters will fill from the last character in the message to the end of the buffer. So if the current character is an ETX, we shouldn't be allowed to advance through memory. Or, if the current character is the last byte in the edit buffer, we shouldn't be allowed to advance through memory. But if we aren't at the end of our text for one reason or another, select the next character by calling the NEXTSL subroutine:

## NEXTCH

NEXTCH JSR GET.SL CMP \#ETX BEQ AN.ETX

Get currently selected character. Is it an ETX?
If so, return to caller, bearing a negative return code.

JSR NEXTSL
RTS

AN.ETX LDA \#\$FF
RTS

If not, select next byte in the buffer, and return positive if we incremented SELECT; negative if SELECT already equaled EA .
Since we are on an ETX, we won't increment
SELECT; we'll just return with a negative return code.

## Select Previous Character

The PREVCH (select-previous-character routine) should work in a manner similar to that used by NEXTCH. NEXTCH increments the SELECT pointer and returns plus, unless SELECT is greater than or equal to EA, in which case NEXTCH preserves SELECT and returns minus. Conversely, PREVCH should decrement SELECT and return plus, unless SELECT is less than or equal to SA, in which case it should preserve SELECT and return minus:

## PREVCH

| PREVCH | SEC |
| :--- | :--- |
|  | LDA SA+1 |
|  | CMP SELECT +1 |
|  | BCC SL.OK |
|  | BNE NOT.OK |

LDA SA CMP SELECT BEQ NO.DEC BNE NOT.OK

SL.OK JSR DEC.SL

LDA \#0
RTS
NOT.OK LDA SA
STA SELECT
LDA SA +1

Prepare to compare.
Is SELECT in a higher page than SA?
If so, SELECT may be decremented.
If SELECT is in a lower page than SA, then it's not okay. We'll have to fix it. SELECT is in the same page as SA.
Is SELECT greater than SA?
If SELECT $=$ SA, don't decrement it. If SELECT is less than SA, it's not okay, so we'll have to fix it.
SELECT is OK, because it's greater than
SA. Thus, we may decrement it and it will remain in the edit buffer.
Set a positive return code... and return.
Since SELECT is less than SA, it is not even in the edit buffer. So give SELECT a legal value, by setting it $=S A$.

```
    STA SELECT+1
    RTS.
NO.DEC LDA #$FF
    RTS
```

    LDA \#0 Set a positive return code...
    and return.
    SELECT $=$ SA, so change nothing. Set
a negative return code and return.

Flush Buffer
To flush the buffer, we'll just fill the buffer with ETX characters:

## FLUSH

| FLUSH | JSR GOTOSA | Set SELECT to the first character position in the buffer. |
| :---: | :---: | :---: |
| FLOOP | LDA \#ETX | Load accumulator with an ETX character... |
|  | JSR PUT.SL | and put it into the buffer. |
|  | JSR NEXTSL | Advance to next byte. |
|  | BPL FLOOP | If we haven't reached the last byte in the buffer, let's repeat the operation for this byte. |
|  | JSR GOTOSA | If we have reached the last byte in the buffer, let's set SELECT to the beginning of the buffer... |
|  | JSR RTS | and return. |

## Print Buffer

To print the buffer, we must print the characters in the edit buffer up to, but not including, the first ETX. Even if there is no ETX in the buffer, we must not print characters from beyond the end of the buffer:

## PRTBUF

PRTBUF JSR GOTOSA PRLOOP JSR GET.SL CMP \#ETX BEQ ENDPRT

Set SELECT to the start of the buffer. Get the currently selected character. Is it an ETX character? If so, stop printing and return.

JSR PR.CHR
JSR NEXTCH
BPL PRLOOP

ENDPRT RTS

If not, print it on all currently selected devices.
Advance SELECT by 1 byte within the buffer.
If we haven't reached the end of the buffer, let's get the next character from the buffer, and handle it.
Since we reached the end of the buffer, let's return.
When this routine returns, the current character is at the end of the message.

## Delete Current Character

To delete the current character, we'll take all the characters that follow it in the text buffer and move them to the left by 1 byte. Here's some code to implement such behavior:

| DELETE | JSR PUSHSL |
| :--- | :--- |
|  | LDA SA +1 |
|  | PHA |
|  | LDA SA |
|  | PHA |
|  | JSR DAHERE |

JSR NEXTSL
JSR SAHERE

JSR MOV.EA
PLA
STA SA
PLA
STA SA+1

Save address of current character.
Save buffer's start address.

Set DEST = SELECT, because we'll move a block of text down to here, to close up the buffer at the current character.
Advance by 1 byte through text buffer, if possible.
Set SA $=$ SELECT, because the block we'll move starts 1 byte above the current character. (Note: the end address of the block we'll move is the end address of the text buffer.)
Move block specified by SA, EA, and DEST.
Restore initial SA (which is the start address of the text buffer, not of the block we just moved).

| JSR POP.SL | Restore SELECT $=$ address of the cur- <br> rent character. |
| :--- | :--- |
| RTS | Return to caller. |

That's the last of the utilities we need. We now have enough code to comprise a simple text editor. Appendices C10 and C11 are listings of this text editor, showing key assignments that work on an Ohio Scientific C-IP. If you have a different system or prefer your editor functions mapped to different keys, simply change the values of the variables in the key table. If you don't want to have a given function, then for that function store a keycode of zero. You'll find this editor very handy for entering tables of ASCII characters into memory, and for entering, editing, and printing short text strings such as titles for your hexdumps and disassembler listings.

## Chapter I2:

## Extending the Visible Monitor

At this point you have the Visible Monitor, the print utilities, two hexdump tools, a table-driven disassembler, a move tool, and a simple text editor. Wouldn't it be nice if they were all combined into one interactive software package? Then you could call any tool or function with a single keystroke. Since the Visible Monitor already uses several keys (0 thru 9; A thru F; G; Space; Return; and Rubout or Clear-Screen), we'll have to map these new functions into unused keys.

Here's a list of keys and the functions they will have in the extended monitor:

H Call a HEXDUMP tool (TVDUMP if the printer is not selected; PRDUMP if the printer is selected).
M Call MOVER, the move tool.
P Toggle the printer flag.
T Call the text editor.
U Toggle the user output flag.
? Call the disassembler (TV.DIS if the printer is not selected; PR.DIS if the printer is selected).

With this assignment of keys to functions, we can select or deselect the printer at any time just by pressing " P ," and likewise the user-driven output device just by pressing "U." We can print or display a hexdump just by pressing "H" and print or display a disassembly just by pressing " f " (which is almost mnemonic if we think of the disassembler as an answer to our question, "What's in the machine?"). We can move anything from anywhere to anywhere else by pressing " M " for move, and we can enter and edit text just by pressing " T " for text editor.

Here's some code to provide these features. Since we want to extend the monitor, this subroutine is called EXTEND:

## EXTEND

EXTEND CMP \#'P
BNE IF.U
LDA PRINTR
EOR \#SFF
STA PRINTR RTS
IF.U CMP \#U
BNE IF.H
LDA USR.FN
EOR \#SFF
STA USR.FN
RTS
IF.H CMP \#'H
BNE IF.M
LDA PRINTR
BNE NEXT. 1
JSR TVDUMP RTS
NEXT. 1 JSR PRDUMP
RTS
IF.M CMP \#M
BNE IF.DIS
JSR MOVER
RTS
IF.DIS CMP \#'?
BNE IF.T
LDA PRINTR
BNE NEXT. 2
JSR TV.DIS
RTS
NEXT. 2 JSR PR.DIS
RTS
IF.T CMP \#'T
BNE EXIT

When EXTEND is called by the Visible Monitor's UPDATE routine, a character from the keyboard is in the accumulator.
Is it the " P " key?
If not, perform the next test.
If so, toggle the printer flag...
and return to caller.
Is it the " U " key?
If not, perform the next test.
If so,
toggle the user-output
flag...
and return.
Is it the " H " key?
If not, perform the next test.
Is the printer selected?
If so, print a hexdump.
If not, dump to screen... and return.
Print a hexdump... and return.
Is it the " M " key?
If not, perform the next test.
If so, call the move tool.
...and return.
Is it the " $\mathrm{Z}^{\prime}$ " key?
If not, perform the next test.
Is the printer selected?
If so, print a disassembly.
If not, dump to screen...
and return.
Print a disassembly... and return.
Is it the " T " key?
If not, return.

|  | JSR EDITOR | If so, call the text editor... |
| :--- | :--- | :--- |
|  | RTS | and return. |
| EXIT | RTS | Extend this subroutine by adding more <br> test-and-branch code here. |

The only remaining step is to modify the Visible Monitor's UPDATE routine so that it calls EXTEND, rather than DUMMY, before it returns. Currently, the Visible Monitor's UPDATE routine calls DUMMY just before it returns, with the bytes $\$ 20$, $\$ 10$, and $\$ 10$ at addresses $\$ 13 D 1, \$ 13 D 2$, and $\$ 13 D 3$, respectively. To make the Visible Monitor's UPDATE routine call EXTEND (instead of DUMMY), you must change $\$ 13 \mathrm{D} 2$ from $\$ 10$ to $\$ \mathrm{BO}$.

You can change this byte with the Visible Monitor itself, provided that you are very careful not to touch any key except the keys that are legal to the unextended Visible Monitor. Once you have changed \$13D2, you may strike any key, but while you are changing \$13D2, striking a key that is not legal within the unextended Visible Monitor will cause the Visible Monitor to crash. Be careful. Once you have changed \$13D2, try out your new extensions of the Visible Monitor by pressing the now legal keys: "H," "M," "P," "U," "?," and "T."

# Chapter 13: Entering the Software into Your System 

Chapters 5 thru 12 present software that will do useful work for you, but only if you can get it into your computer's memory. How you do that will depend on the system you have.

If you have an Apple II, you have an extended machine-language monitor built into your system. If the monitor doesn't come up on RESET, you can invoke it from BASIC with the following BASIC command:

## POKE 0,0:CALL 0 [RETURN]

(The string "[RETURN]" means press the carriage return key.)
This writes a 6502 BRK instruction into location $\$ 0000$, and then executes a call to a machine-language subroutine at location $\$ 0000$. The 6502, upon encountering the BRK instruction, will pass control to the Apple II ROM monitor. You'll know you're in the Apple II monitor because you'll see an asterisk (*) on the screen. Your Apple II documentation should tell you how to use this monitor to enter data into memory, dump memory, etc.

The Ohio Scientific C-IP has a much simpler monitor than the Apple II built into its ROM (read-only memory). Press BREAK on the Ohio Scientific C-IP and then press "M." You'll get the ROM monitor display and can use the ROM monitor to enter hexadecimal object code into memory. Unfortunately, although the Ohio Scientific ROM monitor lets you enter a machine-language program into memory by hand, or even from a cassette file in the proper format, it provides no facility for
recording a machine-language program onto a cassette. So unless you plan to key the Visible Monitor into memory and then leave your computer on forever, you're out of luck. However, you can SAVE a BASIC program on cassette, and then LOAD it from cassette. And that's the key: we'll use the OSI C-1P's ROM BASIC interpreter to help get machine-language programs into memory.

And what if you have an Atari or a PET Computer? Each of these systems features a BASIC interpreter in ROM (read-only memory), but lacks a machine-language monitor. How can you enter hexadecimal object code into memory using only a BASIC interpreter? Perhaps more importantly, even if we manage to enter that object code into memory, how can we save that object code onto a cassette? If all we have is a BASIC interpreter, the simplest solution is to make our object code look like a BASIC program.

That's not so hard. A BASIC program may contain DATA statements, so a simple BASIC program can contain a number of DATA statements, where the DATA statements actually represent, in decimal, the values of successive bytes in the object code. Then the BASIC program can READ those DATA statements and POKE the values it finds into the appropriate section of memory.

## Using BASIC to Load Machine Language

The software in this book can be entered into your computer by RUNning just such a series of BASIC programs. Each of these programs consists of an OBJECT CODE LOADER followed by some number of DATA statements. The first two DATA statements specify the range of DATA statements that follow. Each of the following DATA statements contains ten values: the first value is the start address at which object code from the line is to be loaded; the next eight values represent bytes to be loaded into memory, beginning at the specified address; and the tenth value is the checksum. The checksum is simply the total of the first nine values in the DATA statement. Of these ten values, the first and the tenth will always be greater than 4000 , and the others will always be less than 256.

Appendices E1 through E11 contain this book's object code in the form of such DATA statements. You must type each of these DATA statements into your computer, but the BASIC OBJECT CODE LOADER is designed to let you know if you've made a mistake. It won't catch any error you might make while typing, but it will catch the most likely errors. How? The answer is in the checksum. If you make a mistake while typing in one of these DATA lines, the checksum will almost certainly fail to match the sum of the address and the 8 bytes in the line. Then, when the OBJECT CODE LOADER detects a checksum error, it will identify the offending data statement by printing its line number as well as the address specified by the offending line.

The object code loader will use the following variables:

A The address specified by a data line. Object code from that data line is to be loaded into memory beginning at that address.
BYTE An array of DIMension 8, containing the values of 8 consecutive bytes of object code as specified by a data line.
CHECK The checksum specified by a data line.
FIRST The number of the first DATA statement containing object code.
LAST The number of the last DATA statement containing object code.
LINE A line counter, tracking the number of data lines of object code already loaded into memory.
SUM The calculated sum of the 8 bytes of object code and the address specified by a given data line. If SUM equals the checksum specified by that data line, then the data is probably correct.
TEMP A temporary variable.
Here is the object code loader:

100 REM
110 REM
120 DIM BYTE (8)
130 READ FIRST
140 REM
150 READ LAST
160 REM
170 FOR LINE $=$ FIRST TO LAST
180 GOSUB 300
190 NEXT LINE
200 PRINT "LOADED LINES", FIRST "THROUGH", LAST,"SUCCESSFULY."
210 END
220 REM
230 REM
240 REM
300 READ A
310 SUM $=\mathrm{A}$
320 FOR J=1 TO 8
321 REM
330 READ TEMP: BYTE $(\mathrm{J})=$ TEMP
340 SUM $=$ SUM + BYTE (J)
341 REM
350 NEXT J
360 REM
370 READ CHECK
380 IF SUM < > CHECK THEN 500

## OBJECT CODE LOADER by Ken Skier

:REM Initialize BYTE array.
:REM Get the line number of the first DATA statement containing object code. :REM Get the line number of the last DATA statement containing object code. :REM Read the specified DATA lines. :REM Load next data line into memory. :REM If not done, read next DATA line.
:REM If done, say so.
Subroutine at 300 handles one
DATA statement.
:REM Get address for object code.
:REM Initialize calculated sum of data. :REM Get 8 bytes of object code from data.
:REM Put them in the byte array, and :REM add them to the calculated sum of data.
:REM Now we have the 8 bytes, and we have calculated the sum of the data. :REM Get checksum from data line. :REM If checksum error, handle it.

| $390 \mathrm{FOR} \mathrm{J}=1$ TO 8 | :REM Since there is no checksum error, |
| :---: | :---: |
| 400 POKE A + J -1, BYTE(J) | :REM poke the data into the specified |
| 410 NEXT J | :REM portion of memory, |
| 420 RETURN | :REM and return to caller. |
| 430 REM |  |
| 440 REM | Checksum error-handling code follows. |
| 500 PRINT "CHECKSUM ER | ATA LINE",LINE |
| 510 PRINT "START ADDRE | IN BAD DATA LINE IS", A |
| 520 END |  |
| 530 REM | The next two DATA statements specify |
| 540 REM | the range of DATA statements that |
| 550 REM | contain object code. |
| 570 REM |  |
| 600 DATA ? ? ? | :REM This should be the number of the |
| 610 REM | first DATA statement containing object |
| 611 REM | code. |
| 612 REM |  |
| 620 DATA ? ? ? | :REM This should be the number of the |
| 630 REM | last DATA statement containing object |
| 631 REM | code. |

Once you've entered the BASIC OBJECT CODE LOADER into your computer's memory, SAVE it on a cassette. Remember that by itself the BASIC OBJECT CODE LOADER can do nothing; it needs DATA statements in the proper form to be a complete, useful program. When you're ready to create such a program, LOAD the BASIC OBJECT CODE LOADER from cassette back into memory. Now you're ready to append to it DATA statements from one of the E Appendices - for example, from Appendix E1. Do not append DATA statements from more than one appendix to the same BASIC program. Append as many DATA lines as you can, without using memory above $\$ 0 F F F$ (decimal 4095). You can insure that you don't run over this limit by setting 4095 as the top of memory available to your system's BASIC interpreter. How do you set the top of memory available to the BASIC interpreter? That varies from system to system, so consult the B Appendix for your system.

Before you can append to the OBJECT CODE LOADER all the DATA statements from Appendix E1, your BASIC interpreter may give you an OUT OF MEMORY error (MEMORY FULL). When that happens, delete the last DATA line you appended to the OBJECT CODE LOADER. Let's say you've appended DATA
lines 1000 thru 1022 when you get an OUT OF MEMORY error. Delete DATA line 1022. Now enter the line numbers of the first and last of the object code DATA statements into DATA lines 600 and 620 , like this:

| 600 | DATA | 1000 |
| :--- | :--- | :--- |
| 620 | DATA | 1021 |

DATA lines 600 and 620 , the very first DATA lines in your program, tell the BASIC OBJECT CODE LOADER how many DATA lines of object code follow. Now the OBJECT CODE LOADER can "know" how many DATA lines to read, without reading too few or too many. In this case, DATA lines 600 and 620 tell the OBJECT CODE LOADER that the object code may be found in DATA lines 1000 thru 1021.

Note that DATA lines 600 and 620 each contain one value, whereas the remaining DATA lines each contain ten values.

Now you are ready to RUN the OBJECT CODE LOADER. Unless you're a better typist than I am, you probably made some mistakes while typing in the DATA lines from Appendix E1. Don't worry; the incorrect data will not be blindly loaded into memory. If the BASIC OBJECT CODE LOADER detects a checksum error, it will tell you so, like this:

$$
\begin{array}{ll}
\text { CHECKSUM ERROR IN DATA STATEMENT } & 1012 \\
\text { START ADDRESS GIVEN IN BAD DATA LINE IS } & 4442
\end{array}
$$

This means that data statement 1012 has a checksum error: ie, bad data. To help you double check, the second line of the error message specifies the start address given by the bad data line: this is the first number in the offending data line. These two items of information should make it easy for you to find the bad data line-just look for the DATA statement whose line number is 1012 and whose first value is 4442 . That's the DATA statement you entered incorrectly. Now you need only eyeball the ten numbers in that line, comparing them to the corresponding DATA statement in Appendix E1, and you should quickly find the number or numbers you entered incorrectly. Fix that DATA statement, and RUN the LOADER again.

When you have entered all of the DATA statements correctly, RUNning the LOADER will load the object code they specify into memory. The OBJECT CODE LOADER will then print:

## LOADED LINES aaaa THROUGH bbbb SUCCESSFULLY

where 'aaaa' is the number of the first DATA line of object code, and 'bbbb' is the number of the last DATA line of object code in the program. This message tells you that the BASIC OBJECT CODE LOADER has read and POKE'd the indicated range of DATA statements into memory.

When you see this message, you have verified the program, so SAVE it on a cassette. Then make up a new BASIC program, containing the OBJECT CODE LOADER and the next group of DATA statements from an E Appendix. (Remember not to append DATA lines from more than one E Appendix to the same BASIC program.) Store in lines 600 and 620 the line numbers of the first and last DATA statements you copied from the E Appendix. Verify and SAVE this program as well, and then continue in this manner until you have entered, verified, and SAVE'd BASIC programs containing all of the DATA statements in Appendices E1 thru E10, as well as the DATA statements in the E Appendix containing system data for your computer (one of the Appendices E11 thru E14). RUNning all of those BASIC programs will then enter all of the software presented in this book into your computer's memory.

At this point, you should be ready to transfer control from your computer's BASIC interpreter to the VISIBLE MONITOR.

## Activating the Visible Monitor

Once you have entered the object code for the Screen Utilities, the Visible Monitor, and the System Data Block into your system, you can activate the Visible Monitor by causing the 6502 in your computer to execute a JSR (jump to subroutine) to $\$ 1207$.

Using the Ohio Scientific C-IP ROM monitor, you can activate the Visible Monitor simply by typing:

## 1207G

Using the Apple II ROM monitor, you can call the Visible Monitor with the command:

## G1207 [RETURN]

Using the Atari 400 or 800 with its BASIC cartridge plugged in, you can invoke the Visible Monitor with the BASIC command:

$$
X=\operatorname{USR}(4615) \text { [RETURN] }
$$

In Atari BASIC, you can call a machine-language subroutine by passing the address of that subroutine as a parameter to the USR function. Since $\$ 1207$ is 4615 in decimal, the command $\mathrm{X}=\mathrm{USR}(4615)$ causes Atari BASIC to call the subroutine at $\$ 1207$. (The value returned by that subroutine will then be stored in the BASIC variable $X$ - not in the 6502's $X$ register. But that doesn't concern us because the Visible Monitor isn't designed to return a value to its caller.)

Using the PET 2001, you can invoke the Visible Monitor from BASIC in the immediate mode with the following BASIC command:

SYS (4615)

When you press (RETURN), you'll see the Visible Monitor display, because SYS (4615) causes BASIC to call the subroutine at address 4615 decimal, which is $\$ 1207$-the entry point for the Visible Monitor.

If and when you press " $Q$ " to quit the Visible Monitor, the Visible Monitor will return to its caller - PET BASIC. (The Visible Monitor doesn't leave much room for a PET BASIC program, since your BASIC program and its arrays, variables, etc cannot require memory beyond $\$ 0$ FFF, but the Visible Monitor should work very well with a small PET BASIC program. In any case, it's reassuring to have a new program such as the Visible Monitor return to a familiar one such as the PET BASIC interpreter.)

Once you have activated the Visible Monitor, you should see its display on the screen. If you don't see such a display, then the Visible Monitor has not been entered properly into your system's memory; perhaps you failed to enter the display code properly.

If you do see the Visible Monitor display on the screen, press the space bar. The display should change - specifically, the displayed address should increment, and fields 1 and 2, immediately to the right of the displayed address, may also change.

If nothing changes when you press the space bar, then the display code probably works fine, but you failed to enter the UPDATE code properly.

If the space bar does change the display, then test out the other functions of the Visible Monitor: press RETURN to decrement the selected address; press hexadecimal keys to select a different address; then select an address somewhere in screen memory and place new data into that address. If you picked a place in display memory that is not cleared by the Visible Monitor (ie: a place not in the top five rows of the screen), then you should be able to place arbitrary characters on the screen just by using the Visible Monitor to store arbitrary values in the selected address.

If your Visible Monitor fails to perform properly, you may have entered it into memory incorrectly. Compare the DATA statements you appended to the OBJECT

CODE LOADER with the DATA statements in the E Appendices. Remember: if even 1 byte is entered incorrectly, then in all likelihood the Visible Monitor will fail to function.

To extend the Visible Monitor as described in Chapter 12, store a $\$$ BO in address $\$ 13 \mathrm{D} 2$. To disable the features described in Chapter 12, store a $\$ 10$ in address \$13D2. Now you're really getting your hands on the machine, reaching into memory and operating on the bytes, and with that kind of control, you can do almost anything.

NOTE:
The author intends to provide the software in this book for sale on cassettes compatible with the Apple II, Atari, Ohio Scientific, and PET computers. If you prefer to load your software from cassette, rather than enter it in by hand, contact the author through BYTE Books.

## Appendices

## Appendix AI:

## Hexadecimal Conversion Table

| HEX | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | F | 00 | 000 |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 0 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 0 | 0 |
| 1 | 16 | 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 | 28 | 29 | 30 | 31 | 256 | 4096 |
| 2 | 32 | 33 | 34 | 35 | 36 | 37 | 38 | 39 | 40 | 41 | 42 | 43 | 44 | 45 | 46 | 47 | 512 | 8192 |
| 3 | 48 | 49 | 50 | 51 | 52 | 53 | 54 | 55 | 56 | 57 | 58 | 59 | 60 | 61 | 62 | 63 | 768 | 12288 |
| 4 | 64 | 65 | 66 | 67 | 68 | 69 | 70 | 71 | 72 | 73 | 74 | 75 | 76 | 77 | 78 | 79 | 1024 | 16384 |
| 5 | 80 | 81 | 82 | 83 | 84 | 85 | 86 | 87 | 88 | 89 | 90 | 91 | 92 | 93 | 94 | 95 | 1280 | 20480 |
| 6 | 96 | 97 | 98 | 99 | 100 | 101 | 102 | 103 | 104 | 105 | 106 | 107 | 108 | 109 | 110 | 111 | 1536 | 24576 |
| 7 | 112 | 113 | 114 | 115 | 116 | 117 | 118 | 119 | 120 | 121 | 122 | 123 | 124 | 125 | 126 | 127 | 1792 | 28672 |
| 8 | 128 | 129 | 130 | 131 | 132 | 133 | 134 | 135 | 136 | 137 | 138 | 139 | 140 | 141 | 142 | 143 | 2048 | 32768 |
| 9 | 144 | 145 | 146 | 147 | 148 | 149 | 150 | 151 | 152 | 153 | 154 | 155 | 156 | 157 | 158 | 158 | 2304 | 36864 |
| A | 160 | 161 | 162 | 163 | 164 | 165 | 166 | 167 | 168 | 169 | 170 | 171 | 172 | 173 | 174 | 175 | 2560 | 40960 |
| B | 176 | 177 | 178 | 179 | 180 | 181 | 182 | 183 | 184 | 185 | 186 | 187 | 188 | 189 | 190 | 191 | 2816 | 45056 |
| C | 192 | 193 | 194 | 195 | 196 | 197 | 198 | 199 | 200 | 201 | 202 | 203 | 204 | 205 | 206 | 207 | 3072 | 49152 |
| D | 208 | 209 | 210 | 211 | 212 | 213 | 214 | 215 | 216 | 217 | 218 | 219 | 220 | 221 | 222 | 223 | 3328 | 53248 |
| E | 224 | 225 | 226 | 227 | 228 | 229 | 230 | 231 | 232 | 233 | 234 | 235 | 236 | 237 | 238 | 239 | 3584 | 57344 |
| F | 240 | 241 | 242 | 243 | 244 | 245 | 246 | 247 | 248 | 249 | 250 | 251 | 252 | 253 | 254 | 255 | 3840 | 61440 |

## Appendix A2: ASCII Character Codes

| Code | Char | Code | Char | Code | Char | Code | Char |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | NUL | 20 | SP | 40 | @ | 60 | ‘ |
| 01 | SOH | 21 | . | 41 | A | 61 | a |
| 02 | STX | 22 | " | 42 | B | 62 | b |
| 03 | ETX | 23 | \# | 43 | C | 63 | c |
| 04 | EOT | 24 | \$ | 44 | D | 64 | d |
| 05 | ENQ | 25 | \% | 45 | E | 65 | e |
| 06 | ACK | 26 | \& | 46 | F | 66 | f |
| 07 | BEL | 27 | ' | 47 | G | 67 | g |
| 08 | BS | 28 | ( | 48 | H | 68 | h |
| 09 | HT | 29 | ) | 49 | I | 69 | i |
| OA | LF | 2A | * | 4A | J | 6 A | j |
| OB | VT | 2B | + | 4B | K | 6B | k |
| ${ }^{0} \mathrm{C}$ | FF | 2 C | , | 4C | L | 6C | 1 |
| OD | CR | 2D | - | 4D | M | 6D | m |
| OE | SO | 2E | - | 4E | N | 6 E | $n$ |
| OF | SI | 2F | 1 | 4F | O | 6 F | O |
| 10 | DLE | 30 | 0 | 50 | P | 70 | p |
| 11 | DC1 | 31 | 1 | 51 | Q | 71 | q |
| 12 | DC2 | 32 | 2 | 52 | R | 72 | r |
| 13 | DC3 | 33 | 3 | 53 | S | 73 | S |
| 14 | DC4 | 34 | 4 | 54 | T | 74 | t |
| 15 | NAK | 35 | 5 | 55 | U | 75 | u |
| 16 | SYN | 36 | 6 | 56 | V | 76 | v |
| 17 | ETB | 37 | 7 | 57 | W | 77 | w |
| 18 | CAN | 38 | 8 | 58 | X | 78 | x |
| 19 | EM | 39 | 9 | 59 | Y | 79 | y |
| 1 A | SUB | 3A | : | 5 A | Z | 7A | z |
| 1B | ESC | 3B | ; | 5B | , | 7 B | 1 |
| 1C | FS | 3 C | $<$ | 5 C | 1 | 7 C | 1 |
| 1D | GS | 3D | = | 5D | ] | 7 D | \} |
| 1 E | RS | 3E | $>$ | 5 E | $\wedge$ | 7E |  |
| 1F | US | 3F | ? | 5 F | - | 7F | DEL |

## Appendix A3:

## 6502 Instruction Set - Mnemonic List

| ADC | Add Memory to Accumulator with Carry |
| :--- | :--- |
| AND | "AND" Memory with Accumulator |
| ASL | Shift Left One Bit (Memory or Accumulator) |
| BCC | Branch on Carry Clear |
| BCS | Branch on Carry Set |
| BEQ | Branch on Result Zero |
| BIT | Test Bits in Memory with Accumulator |
| BMI | Branch on Result Minus |
| BML | Branch on Result not Zero |
| BPL | Branch on Result Plus |
| BRK | Force Break |
| BVC | Branch on Overflow Clear |
| BVS | Branch on Overflow Set |
|  |  |
| CLC | Clear Carry Flag |
| CLD | Clear Decimal Mode |
| CLI | Clear Interrupt Disable Bit |
| CLV | Clear Overflow Flag |
| CMP | Compare Memory and Accumulator |
| CPX | Compare Memory and Register X |
| CPY | Compare Memory and Register Y |
|  |  |
| DEC | Decrement Memory |
| DEX | Decrement Register X |
| DEY | Decrement Register Y |
|  |  |
| EOR | "Exclusive Or" Memory with Accumulator |
| INC | Increment Memory |
| INX | Increment Register X |
| INY | Increment Register Y |

JMP Jump to New Location
JSR Jump to New Location Saving Return Address
LDA Load Accumulator with Memory
LDX Load Register X with Memory
LDY $\quad$ Load Register Y with Memory
LSR Shift Right One Bit (Memory or Accumulator)
NOP No Operation
ORA "OR" Memory with Accumulator
PHA Push Accumulator on Stack
PHP Push Processor Status on Stack
PLA Pull Accumulator from Stack
PLP Pull Processor Status from Stack
ROL Rotate One Bit Left (Memory or Accumulator)
ROR Rotate One Bit Right (Memory or Accumulator)
RTI Return from Interrupt
RTS Return from Subroutine
SBC Subtract Memory from Accumulator with Borrow
SEC Set Carry Flag
SED Set Decimal Mode
SEI Set Interrupt Disable Status
STA Store Accumulator in Memory
STX Store Register X in Memory
STY Store Register Y in Memory
TAX Transfer Accumulator to Register $X$
TAY Transfer Accumulator to Register $Y$
TSX Transfer Stack Pointer to Register X
TXA Transfer Register $X$ to Accumulator
TXS Transfer Register X to Stack Pointer
TYA Transfer Register Y to Accumulator

## Appendix A4:

## 6502 Instruction Set — Opcode List

00 - BRK
01 - ORA - (Indirect, X)
02 - Future Expansion
03 - Future Expansion
04 - Future Expansion
05 - ORA - Zero Page
06 - ASL - Zero Page
07 - Future Expansion
08 - PHP
09 - ORA - Immediate
OA - ASL - Accumulator
OB - Future Expansion
OC - Future Expansion
OD - ORA - Absolute
OE - ASL - Absolute
OF - Future Expansion

10 - BPL
11 - ORA - (Indirect), Y
12 - Future Expansion
13 - Future Expansion
14 - Future Expansion
15 - ORA - Zero Page, X
16 - ASL - Zero Page,X
17 - Future Expansion

18 - CLC
19 - ORA - Absolute, Y
1A - Future Expansion
1B - Future Expansion
1C - Future Expansion
1D - ORA - Absolute, X
1E - Future Expansion
1F - Future Expansion

20 - JSR
21 - AND - (Indirect, X)
22 - Future Expansion
23 - Future Expansion
24 - Bit - Zero Page
25 - AND - Zero Page
26 - ROL - Zero Page
27 - Future Expansion
28 - PLP
29 - AND - Immediate
$2 \mathrm{~A}-\mathrm{ROL}$ - Accumulator
2B - Future Expansion
2C - BIT - Absolute
2D - AND - Absolute
2 E - ROL - Absolute
2F - Future Expansion
$30-\mathrm{BMI}$
31 - AND - (Indirect), Y
32 - Future Expansion
33 - Future Expansion
34 - Future Expansion
35 - AND - Zero Page, X
36 - ROL - Zero Page, X
37 - Future Expansion
38 - SEC
39 - AND - Absolute, Y
3A - Future Expansion
3B - Future Expansion
3C - Future Expansion
3D - AND - Absolute, X
3F - Future Expansion

40 - RTI
41 - EOR - (Indirect, X)
42 - Future Expansion
43 - Future Expansion
44 - Future Expansion
45 - EOR - Zero Page
46 - LSR - Zero Page
47 - Future Expansion
48 - PHA
49 - EOR - Immediate
4A - LSR - Accumulator
$4 B$ - Future Expansion
4C - JMP - Absolute
4D - EOR - Absolute
4E - LSR - Absolute
4 F - Future Expansion

50 - BVC
51 - EOR - (Indirect), Y
52 - Future Expansion
53 - Future Expansion
54 - Future Expansion
55 - EOR - Zero Page, X
56 - Zero Page, X
57 - Future Expansion

58 - CLI
59 - EOR - Absolute, Y
5A - Future Expansion
5B - Future Expansion
5C - Future Expansion
5D - EOR - Absolute, X
5E - LSR - Absolute, X
5F - Future Expansion

60 - RTS
61 - ADC - (Indirect, $X$ )
62 - Future Expansion
63 - Future Expansion
64 - Future Expansion
65 - ADC - Zero Page
66 - ROR - Zero Page
57 - Future Expansion
68 - PLA
69 - ADC - Immediate
6A - ROR - Accumulator
6B - Future Expansion
6C - JMP - Indirect
6D - ADC - Absolute
6E - ROR - Absolute
6F - Future Expansion

70 - BVS
71 - ADC - (Indirect), Y
72 - Future Expansion
73 - Future Expansion
74 - Future Expansion
75 - ADC - Zero Page, $X$
76 - ROR - Zero Page, X
77 - Future Expansion
78 - SEI
79 - ADC Absolute, $Y$
7A - Future Expansion
7B - Future Expansion
7C - Future Expansion
7D - ADC - Absolute, X
7E - ROR - Absolute, X
7F - Future Expansion

| 80 - Future Expansion | A8 - TAY |
| :---: | :---: |
| 81 - STA - (Indirect, X) | A9 - LDA - Immediate |
| 82 - Future Expansion | AA - TAX |
| 83 - Future Expansion | AB - Future Expansion |
| 84 - STY - Zero Page | AC - LDY - Absolute |
| 85 - STA - Zero Page | AD - LDA - Absolute |
| 86 - STX - Zero Page | AE - LDX - Absolute |
| 87 - Future Expansion | AF - Future Expansion |
| 88 - DEY |  |
| 89 - Future Expansion |  |
| 8A - TXA | B0 - BCS |
| 8B - Future Expansion | B1 - LDA - (Indirect), Y |
| 8C - STY - Absolute | B2 - Future Expansion |
| 8D - STA - Absolute | B3 - Future Expansion |
| 8E - STX - Absolute | B4 - LDY - Zero Page, X |
| 8F-Future Expansion | B5 - LDA - Zero Page, X |
|  | B6 - LDX - Zero Page, Y |
|  | B7 - Future Expansion |
| $90-\mathrm{BCC}$ | B8-CLV |
| 91 - STA - (Indirect), Y | B9 - LDA - Absolute, Y |
| 92 - Future Expansion | BA - TSX |
| 93 - Future Expansion | BB - Future Expansion |
| 94 - STY - Zero Page, X | BC - LDY - Absolute, X |
| 95 - STA - Zero Page, X | BD - LDA - Absolute, X |
| 96 - STX - Zero Page, Y | BE - LDX - Absolute, Y |
| 97 - Future Expansion | BF - Future Expansion |
| 98 - TYA |  |
| 99 - STA - Absolute, Y |  |
| 9A - TXS | CO - CPY - Immediate |
| 9 B - Future Expansion | C1 - CMP - (Indirect, X ) |
| 9 C - Future Expansion | C2- Future Expansion |
| 9D - STA - Absolute, X | C3-Future Expansion |
| 9 E - Future Expansion | C4-CPY - Zero Page |
| 9 F - Future Expansion | C5 - CMP - Zero Page |
|  | C6- DEC - Zero Page |
|  | C7 - Future Expansion |
| A0 - LDY - Immediate | C8 - INY |
| A1 - LDA - (Indirect, X ) | C9 - CMP - Immediate |
| A2 - LDX - Immediate | CA - DEX |
| A3 - Future Expansion | CB - Future Expansion |
| A4 - LDY - Zero Page | CC - CPY - Absolute |
| A5 - LDA - Zero Page | CD - CMP - Absolute |
| A6 - LDX - Zero Page | CE - DEC - Absolute |
| A7 - Future Expansion | CF - Future Expansion |

DO - BNE
D1 - CMP - (Indirect), Y
D2 - Future Expansion
D3 - Future Expansion
D4 - Future Expansion
D5 - CMP - Zero Page, $X$
D6 - DEC - Zero Page, X
D7 - Future Expansion
D8 - CLD
D9 - CMP - Absolute, Y
DA - Future Expansion
DB - Future Expansion
DC - Future Expansion
DD - CMP - Absolute, X
DE - DEC - Absolute, X
DF - Future Expansion

EO - CPX - Immediate
E1 - SEC - (Indirect, X)
E2 - Future Expansion
E3 - Future Expansion
E4 - CPX - Zero Page
E5 - SBC - Zero Page
E6 - Zero Page
E7 - Future Expansion

E8 - INX
E9 - SBC - Immediate
EA - NOP
EB - Future Expansion
EC - CPX - Absolute
ED - SBC - Absolute
EE - INC - Absolute
EF - Future Expansion

FO - BEQ
F1 - SBC - (Indirect), Y
F2 - Future Expansion
F3 - Future Expansion
F4 - Future Expansion
F5 - SBC - Zero Page, X
F6 - INC - Zero Page, X
F7 - Future Expansion
F8 - SED
F9 - SBC - Absolute, Y
FA - Future Expansion
FB - Future Expansion
FC - Future Expansion
FD - SBC - Absolute, X
FE - INC - Absolute, X
FF - Future Expansion

## Appendix A5:

Instruction Execution Times (indackergcles)
Accumulator
Immediate
Zero Page
Zero Page, $X$
Zero Page, $Y$
Absolute
Absolute, $X$
Absolute, $Y$
Implied
Relative
(Indirect), $X$
(Indirect), $Y$
Absolute Indirect


$$
\begin{aligned}
& \begin{array}{l}
\text { Accumulator } \\
\text { Immediate } \\
\text { Zero Page } \\
\text { Zero Page, X } \\
\text { Zero Page, Y }
\end{array} \\
& \begin{array}{l}
\text { Absolute } \\
\text { Absolute, X } \\
\text { Absolute, Y } \\
\text { Implied } \\
\text { Relative }
\end{array} \\
& \begin{array}{l}
\text { (Indirect), } \mathrm{X} \\
\text { (Indirect), } \mathrm{Y}
\end{array} \\
& \text { Add one cycle if indexing across page boundary } \\
& \text { ** Add one cycle if branch is taken, Add one additional if branching operation } \\
& \text { crosses page boundary }
\end{aligned}
$$

## Appendix A6:

## 6502 Opcodes by Mnemonic and Addressing Mode



| ADC | 6D | 7 D | 79 | . | 69 |  | 61 | 71 | . | 65 | 75 | , |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| AND | 2D | 3D | 39 |  | 29 | . | 21 | 31 | , | 25 | 35 | , |
| ASL | OE | 1E | , | OA | , | . | . | . | . | 06 | 16 | . |
| BCC | . | . | . | . | - | . | . | - | 90 | . | . | . |
| BCS | - | , | . | . | . | . | . | . | B0 | , | . | . |
| BEQ | . | . | . | , | . | . | . | . | F0 | . | . | , |
| BIT | 2 C | . | . | . | . | . | . | . | . | 24 | . | - |
| BMI | . | - | - | , | , | , | . | , | 30 | , | , | , |
| BNE | . | . | - | . | . | . | . | . | D0 | . | . | . |
| BPL | , | . | . | . | . | . | , | . | 10 | . | . | . |
| BRK | . | , | - | . | . | 00 | . | . | . | , | . | . |
| BVC | . | . | . | . | . | , | , | . | 50 | , | , | - |
| BVS | - | . | . | - | - | . | - | . | 70 | . | - | - |
| CLC | . | , | : | . | . | 18 | . | . | . | . | , | . |
| CLD | - | . | . | , | . | D8 | , | - | - | . | . |  |
| CLI | , | , | , | . | . | 58 | . | . | . | . | . |  |



ABSOLUTE

思 品 要

X＇GOVd OYBZ
خ
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N
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N

Mnemonics ．．．．． 60 60
E9
E1 F1
E5 F5
SEC ．．．．． 38
SED ．．．．．F8
SEI ．．．．． 78
STA 8D 9D 99 ．．．． 81 ． 81 ． 85
STX 8E ．．．．．．．．． 86
STY 8C ．．．．．．．．． 8494
TAX ．．．．．AA
TAY ．．．．．A8
TSX
BA
TXA
8A
TXS
9A
TYA
98


## Appendix BI: The Ohio Scientific Challenger I-P

The Ohio Scientific Challenger I-P is the simplest of the systems considered in this book. Its screen is mapped in the manner described in Chapter 5: the lowest screen address is in the upper left corner, and the screen addresses increase uniformly as you move to the right and down the screen. Any ASCII character stored in screen memory will be displayed properly on the video screen; it is not necessary to replace the ASCII character with a system-specific display code. Therefore, the system data block may be initialized as shown in Appendices C13 and E12.

Incidentally, the OSI C-IP's screen TVT subroutine at $\$ \mathrm{BF} 2 \mathrm{D}$ stores the relative location of the cursor in $\$ 0200$. Modify $\$ 0200$ and you change the next location at which a character will be printed to the screen.

If you have an Ohio Scientific BASIC-in-ROM system other than the Challenger I-P, it may have different character input/output routines. If so, examine the following locations:

| BASIN | \$FFEB | General character-input routine for OSI <br> BASIC-in-ROM. <br> BASOUT |
| :--- | :--- | :--- |
|  | \$FFEE | General character-output routine for |
|  |  | OSI BASIC-in-ROM. |

For example, in the OSI C-IP you can get a character from the keyboard by calling $\$$ FEED, or you may call OSI's general character-input routine at $\$$ FFEB. This routine gets a character from the keyboard unless the SAVE flag is set, in which case it gets a character from the cassette input port. Similarly, in the OSI C-IP you can print a character to the screen by calling $\$ \mathrm{BF} 2 \mathrm{D}$, or send a character to the cassette output port by calling $\$$ FCB1. Or, you can simply call OSI's general characteroutput routine at $\$$ FFEE, which outputs the accumulator to the screen and, if the SAVE flag is set, echoes to the serial port as well.

Thus, even if you don't know the addresses of your OSI system's specific I/O routines, you can set ROMKEY $=\$$ FFEB and ROMTVT $=\$$ FFEE. When you RESET
your system, the Ohio Scientific Operating System will automatically "hook" those routines to your keyboard for input and to your screen for output.

## Setting the Top of Memory

If you wish to load object code using the BASIC OBJECT CODE LOADER (see Chapter 13) you must first set the top of memory available to your BASIC interpreter to \$OFFF. Do this as part of cold-starting BASIC. To cold-start BASIC, turn on your OSI computer, press the (BREAK) key, and then press ' C '. The screen will prompt, "Memory Size?" Type "4095" and then press (RETURN). Now BASIC will use the lowest 4 K of RAM, leaving memory from $\$ 1000$ and up available to machine-language programs.

With the top of memory set to \$OFFF, you may enter and RUN the BASIC programs that load object code into your computer's memory.

## Calling Machine-Language Code from BASIC

To call a machine-language subroutine from BASIC, first set the pointer at $\$ 000 \mathrm{~B}, 000 \mathrm{C}$ so it points to the subroutine, and then call that subroutine with BASIC's USR function, either in the immediate mode or from within a BASIC program. For example, let's say you wish to call the Visible Monitor from BASIC. The Visible Monitor's entry point is at $\$ 1207$, so we must make $\$ 000 \mathrm{~B}, 000 \mathrm{C}$ point to $\$ 1207$. This means storing 07 in $\$ 000 B$, and storing $\$ 12$ (decimal 18) in $\$ 000 \mathrm{C}$. The following line will do that for us:

POKE 11,7:POKE 12,18

Now we may invoke the Visible Monitor with the line:

$$
X=\operatorname{USR}(X)
$$

or with any other line that uses the USR function.
Note that the USR function does not set a BASIC variable equal to the contents of some register in the 6502; in fact, the line $X=\operatorname{USR}(X)$ will not change the value of the BASIC variable $X$ at all. Thus, the USR function lets you activate any desired machine-language subroutine, but it doesn't let you capture a value returned by such
a subroutine. If you want a machine-language subroutine to return some value which you can then use in a BASIC program, you'll have to make the machinelanguage subroutine store its value or values somewhere in memory, and then have the BASIC program PEEK that memory location after it has called the machinelanguage subroutine via the USR function.

# Appendix B2: The PET 200I 

## Display Memory

The PET screen is mapped conventionally, with the HOME address at $\$ 8000$ ( 32,768 decimal). It has 25 rows, each consisting of 40 characters. The address of each screen location is $40(\$ 28)$ greater than the address of the screen location directly above it. Thus, the screen parameters for the PET 2001 are:

| HOME | .WORD \$8000, |  |
| :--- | :--- | :--- |
| ROWINC | .BYTE \$28 |  |
| TVCOLS | .BYTE 39 | (We count columns from zero.) |
| TVROWS | .BYTE 24 | (We count rows from zero.) |

## PET Character Set

However, although the PET screen buffer is mapped conventionally, you cannot simply store an ASCII character in screen memory if you wish to see that ASCII character on the screen. The PET character generator introduces a few wrinkles and you must compensate carefully if you are to display ASCII characters properly on the screen.

For example, if you store $\$ 31$ (the code for an ASCII " 1 ") in the PET's display memory, then you will see a " 1 " displayed on the screen. So far, so good. The same is true for all ASCII digits and for some ASCII punctuation marks. But if you store $\$ 45$ (ASCII code for an upper case " $E$ ") in screen memory, then you won't see an " $E$ " on the screen: you'll see either a lowercase " e " or else a horizontal line segment much longer than a hyphen. What's happening?

The PET 2001 features a memory location, \$E84C (59468) which has a special effect on the video-display circuitry. The value stored in that address selects for the video display one character set or another.

To see how the choice of character set affects the display, enter the following BASIC program into your PET:

```
100 REM DISPLAY PET CHARACTER SET
110 REM IN 16 BY 16 MATRIX
120 REM
130 HOME =32768
140 CHAR =0
150 FOR ROW=0 TO 15
160 FOR COL=0 TO 15
170 POKE (HOME + COL) + (40*ROW),CHAR
180 CHAR = CHAR +1
190 NEXT COL
200 NEXT ROW
210 END
```

Before running this program, clear the screen by holding down the PET's SHIFT key at the same time that you depress the CLR/HOME key. When the screen is clear, use the CRSR SOUTH key to move the cursor down seventeen rows. Then type RUN and press RETURN. You'll see one PET character set appear in a 16 by 16 matrix in the upper left portion of your PET's screen.

What you'll see on your screen will look like table B2.1 (without the labeled axes).

Table B2.1: The PET character set.
RIGHT NYBBLE OF CHARACTER
$-0-1-2-3-4-5-6-7-8-9-A-B-C-D-E-F$
LEFT NYBBLE OF CHARACTER

| 0- | (1) | A | B | D | E | F | G | H |  |  | J |  | L | M |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1- | P | Q | R | T | U | V | W | X |  |  |  |  | , |  |  |  |
| $2-$ |  | ! | " | \$ | \% | \& | , | ( |  |  | * | + |  |  |  |  |
| $3-$ | 0 | 1 | 2 | 4 | 5 | 6 | 7 | 8 |  |  |  | ; | < |  |  |  |
| 4 - | - | a | b | d | e | f | g | h |  |  | j | k | 1 |  | , |  |
| $5-$ | p | q | r | t | u | v | w | x |  |  |  |  |  |  |  |  |
| 6- |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $7-$ |  | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 8 - | @ | A | B | D | E | F | G | H |  |  | I | K | L |  | , | O |
| $9-$ | P | Q | R | T | U | V | W, | X | Y |  |  | 1 |  |  |  | - |
| A- |  | ! | " | \$ | \% | \& | , | ( |  |  |  | $+$ |  |  |  |  |
| B- | 0 | 1 | 2 | 4 | 5 | 6 | 7 | 8 |  |  |  | ; | $<$ |  |  |  |
| C- | - | a |  | d | e | f | g | h |  |  |  | k |  |  |  |  |
| D- | P | q | r | t | $u$ | $v$ | w | x |  |  |  |  |  |  |  |  |
| E- |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

In this chart, special graphic characters are indicated by an underline. Look at your PET screen to see those special graphics in all their glorious detail.

Note that the characters for $\$ 80$ thru $\$ F F$ are the same as for $\$ 00$ thru $\$ 7 \mathrm{~F}$, but in reverse intensity. The low 128 characters ( $\$ 00$ thru $\$ 7 \mathrm{~F}$ ) are "normal" - that is, white characters on a dark background; whereas the high 128 characters ( $\$ 80$ thru $\$ \mathrm{FF}$ ) are in reverse video - dark characters in a white background. An " A " in normal intensity may be displayed by storing an $\$ 01$ somewhere in the screen memory; a reverse intensity " A " may be displayed by storing an $\$ 81$ somewhere in screen memory. From this pattern we can derive a handy corollary: to reverse the intensity of any character on the screen, simply reverse its bit 7. You don't even have to know what the character represents; just toggle bit 7 and you change its intensity.

The chart in figure B2.1 (and on your PET screen) shows one complete character set because the BASIC program stores every 8 -bit value, from $\$ 00$ thru $\$ \mathrm{FF}$, into the screen buffer. But I mentioned two character sets. What must you do to see the second character set?

If the cursor is within three rows of the bottom of the screen, move it up so that it is at least three rows above the bottom of the screen. This will insure that you don't scroll part of the character set up off the screen when you execute the following BASIC command in the immediate mode:

POKE 59468,12

Did that change the display? If not, then execute the following BASIC command in the immediate mode (again being sure that the cursor is at least three rows from the bottom of the screen):

POKE 59468,14

Depending on the value stored in 59468 (\$E84C), one or another character set will be displayed. The values of the bytes stored in screen memory will not change when you change the contents of $\$ E 84 \mathrm{C}$, but in some cases the displayed characters will change. In the ranges 00 thru $\$ 3 \mathrm{~F}$ and $\$ 80$ thru $\$ B F$, the two character sets are identical. But in the ranges $\$ 40$ thru $\$ 7 \mathrm{~F}$ and $\$ \mathrm{CO}$ thru $\$ \mathrm{FF}$, they differ.

Both character sets include numbers, uppercase letters, and certain punctuation marks; but only one character set includes lowercase letters and the remaining punctuation marks. The second character set lacks lowercase letters and these punctuation marks, offering instead a set of special graphics characters, including playingcard suits. POKE 59468,14 to select the former character set (thereby making possible the display of all printable ASCII characters); POKE 59468,12 to select the latter character set (thereby making possible the display of the gaming graphics).

## FIXCHR

Note that neither character set corresponds directly to ASCII. If you have an ASCII character in the accumulator and you want to display the appropriate graphic character on the screen, you must first call FIXCHR (as TV.PUT does, in Chapter 5). When an ASCII character is passed in the accumulator, FIXCHR must return in the accumulator the proper PET display code for that character. FIXCHR's caller may then store this display code in memory, thereby placing on the screen an appropriate image of the original ASCII character.

How will FIXCHR work? By examining the PET character set and comparing it to Appendix A2, ASCII codes, we can see a solution in the form of the following algorithm:

- If a character is in the range $\$ 40$ thru $\$ 5 \mathrm{~F}$, subtract $\$ 40$ and return.
- If a character is in the range $\$ 20$ thru $\$ 3 F$, return.
- If a character is in the range $\$ 60$ thru $\$ 7 \mathrm{~A}$, store a decimal 14 in 59468 to select the character set that has lower case letters; and return.
- All other input characters are either ASCII control codes, for which there are no agreed-upon graphics, or else PET special graphics characters, so just return.

Examine the tables yourself to see if this algorithm will work.

## FIXCHR

| FIXCHR | AND \#\$7F | Clear bit 7, so the character will be in the legal ASCII range. |
| :---: | :---: | :---: |
|  | SEC | Prepare to compare. |
|  | CMP \#\$40 | If it's less than \$40, return. |
|  | BCC FIXEND |  |
|  |  | Okay. The character is greater than \$40. |
|  | CMP \#\$60 | Is it greater than \$5F? |
|  | BCS LOWERC | If so, handle it as lowercase. |
|  |  | Okay. The character is in the range \$40-\$5F |
|  | SBC \#\$40 | Subtract \$40 for proper display code. |
|  | RTS |  |
| LOWERC | LDX \#14 | Since we have a lowercase letter, let's select the character set that |
|  | STX 59468 | has lowercase letters. |
| FIXEND | RTS | Return, bearing PET display code for character originally in accumulator. |

Call FIXCHR with an ASCII character in the accumulator. FIXCHR will return with the corresponding PET display code in the accumulator. When it returns, its caller may store the accumulator anywhere in screen memory, thus displaying an image of the original ASCII character.

## PET Keyboard Input Routine

To get an ASCII character from the PET keyboard, call the following subroutine:

## PETKEY

JSR \$FFE4
CMP \#0
BEQ PETKEY

AND \#\$7F
RTS

Call PET ROM key scan routine.
Zero means no key.
If no key, scan again.
A new key is in the accumulator. If the shift key was down, bit 7 is set. So clear bit 7 , just to be sure we've got a legal ASCII character.
Return with ASCII character in the accumulator.

This subroutine yields the uppercase ASCII code for any letter key that you depress, and the proper ASCII code for any digit key or punctuation key.

## PET TVT Routine

To print an ASCII character to the screen, call \$FFD2, a PET ROM routine I will refer to as PETTVT.

Any printable ASCII character passed to \$FFD2 (or, apparently, to \$E3EA or $\$ F 230$ ) will be printed properly to the screen at the PET's current TVT screen location. You may change the PET's current TVT screen location (which is not the same as the current location used by the screen utilities in Chapter 5) by calling PETTVT with the accumulator holding any of the control codes from Table B2.1.

Table B2.1: Control codes that affect the next character to be printed by PETTVT.
Character Name Code Function

| CURSOR NORTH | \$91 | Move current location up by |
| :---: | :---: | :---: |
| CURSOR EAST | \$1D | Move current location one column to the right. |
| CURSOR SOUTH | \$11 | Move current location down by one row. |
| CURSOR WEST | \$9D | Move current location left by one column. |
| INSERT | \$94 | Move current character, and all characters to its right, one column to the right. |
| DELETE | \$14 | Move current character, and all characters to its right, one column to the left. |
| HOME | \$13 | Set current location to upper left of screen. |
| CLEAR | \$93 | Set current location to the upper left corner and clear the screen. |
| REVERSE | \$12 | Select reverse video for following characters. |
| REVERSE-OFF | \$92 | Select normal video mode for following characters. |

These control codes may be passed directly to PETTVT, or they may be included within a string of characters to be printed by "PRINT:" or "PR.MSG." For example, if you wish to clear the screen before printing a message, just put the CLEAR character (\$93) at the beginning of your message string, immediately following the STX. The message-printing subroutine will get the CLEAR character and pass it to PR.CHR, which, in turn, will pass it through the ROMTVT vector on to the PETTVT routine. The PETTVT routine will then clear the screen and set the current location to the upper left corner of the screen.

The next character in the string will then be printed in the upper left corner of a clear screen. If, instead of printing your message at the top row of a clear screen, you'd prefer to print it in the fifth row of a clear screen, just follow the CLEAR character with four CURSOR-SOUTH characters ( $\$ 11, \$ 11, \$ 11, \$ 11$ ), and follow the four cursor-south characters with the text of your message. Following the text of your message, of course, you must include an ETX (\$FF).

You might never use the PETTVT control codes, but it's good to know they're available, should you ever want your PET's display screen to perform as something more than a glass teletype.

## System Data Block

To run on a PET 2001, the software in this book requires the system data block shown in Appendices C14 and E13.

## Setting the Top of Memory

Before you can use the BASIC OBJECT CODE LOADER (presented in Chapter 12) to load object code into your PET's memory, you must insure that your PET's BASIC interpreter leaves undisturbed all memory above \$0FFF ( 4095 decimal). The PET BASIC interpreter will do as we wish if we set its top-of-memory pointer appropriately. The top-of-memory pointer specifies the highest address that may be used for the storage of BASIC program lines, variables, and strings. Memory above that address is off-limits to BASIC.

As you may know, there is more than one version of the PET 2001 by Commodore. Some PET's have software in "old" ROMS (REV 2 ROMS), and others have software in "new" ROMS (REV 3 ROMS). As far as the software in this book is concerned, old ROM PETS and new ROM PETS are the same, since the ROM routines we care about are accessible from the same addresses in both old and new ROM PETS. Therefore, until now I haven't even mentioned that the PET 2001 comes in two flavors. But now you must discover whether you have an old ROM or a new ROM PET, because otherwise you won't be able to set the top of memory.

Old ROM and new ROM PETS each contain a machine-language subroutine to clear the screen, but in new ROM PETS that subroutine is at $\$$ E229 (57897 decimal), and in old ROM PETS that subroutine is as \$E236 (57910 decimal). To see what ROMS are in your PET, use the PET's screen editor to place some characters on the screen, and then type:

SYS (57897)
and press (RETURN). Does the screen clear? If so, you've got a new ROM PET. If not, turn off your PET, turn it on, place some characters on the screen, and then type:

SYS (57910)
and press (RETURN). Does the screen clear? If so, you've got an old ROM PET. If not, then your PET contains neither Rev 2 ROMS nor Rev 3 ROMS, and you'll have to consult your system's documentation carefully to discover the address of the top-of-memory pointer.

On old ROM PETS, the top-of-memory pointer is at 134 and 135 ( $\$ 86,87$ ). On new ROM PETS, the top-of-memory pointer is at 52 and 53 ( $\$ 34,35$ ). Regardless of the location of the top-of-memory pointer, we want to set the low byte of that pointer equal to $\$ \mathrm{FF}$ ( 255 decimal), and the high byte of that pointer equal to $\$ 0 \mathrm{~F}$ ( 15 decimal), so that the pointer itself points to $\$ 0$ FFF. That will leave memory from
$\$ 1000$ and up available to machine-language programs.
Thus, we set the top of memory on an old ROM PET with:

POKE 134,255:POKE 135,15

Similarly, we set the top of memory on a new ROM PET with:

POKE 34,255:POKE 35,15

Once you have set the top of memory available to your PET's BASIC interpreter, you may enter the BASIC OBJECT CODE LOADER and the DATA statements from Appendices E1 thru E11, and from Appendix E13. Remember to set the top of memory not only when typing in these DATA statements, but when RUNning the OBJECT CODE LOADER, as well.

# Appendix B3: <br> The Apple II 

## Apple Display

The display memory of the Apple II is mapped in a manner that is much more complex than the Ohio Scientific or PET computers. On each of these other systems, only one portion of memory is mapped to the screen. The screen cannot display the contents of any other bank of memory (unless, of course, you copy the contents of another bank of memory into the display memory). But the Apple II may display the contents of any of four banks of memory: Low-Resolution Graphics and Text Page 1, Low-Resolution Graphics and Text Page 2, High-Resolution Graphics Page 1, and High-Resolution Graphics Page 2. Table B3.1 summarizes the locations of these pages in memory.

Table B3.1: Banks of display memory in the Apple II.

Hexadecimal
Low-Resolution Graphics and Text Page 1:
Low-Resolution Graphics and Text Page 2: Hi-Resolution Graphics Page 1: Hi-Resolution Graphics Page 2:
\$0400-\$07FF

Decimal

Note that each of these display pages takes up much more than one hexadecimal page ( 256 bytes). A display page is simply an area of any size memory, whose contents may be displayed on the screen. Each low-res display page occupies four hexadecimal pages, and each hi-res display page occupies 32 hexadecimal pages. Why are the hi-res display pages bigger than the low-res display pages? Hi-res means highresolution, and higher resolution requires more information.

How do you make the video screen show the contents of a given display page? You need only store a zero in a particular address. Certain addresses in the Apple II signal the video-display circuitry whenever data are written to them. The videodisplay circuitry responds to these signals by displaying the contents of a given bank of memory. These special addresses, or display selectors, are given in Table B3.2.

Table B3.2: Addresses that affect the APPLE II Display.

| Hexadecimal | Decimal | Label | Purpose of Address |
| :---: | :---: | :---: | :---: |
| \$C050 | -16304 | TXTCLR | Store a 0 here to set graphics mode. |
| \$C051 | -16303 | TXTSET | Store a 0 here to set text mode. |
| \$C052 | -16302 | MIXCLR | Store a 0 here to set bottom four lines to graphics. |
| \$C053 | -16301 | MIXSET | Store a 0 here to select text/ graphics mix (bottom four lines text). |
| \$C055 | -16299 | HISCR | Store a 0 here to select Page 2. |
| \$C056 | -16298 | LORES | Store a 0 here to select lowresolution graphics and text page. |
| \$C057 | -16297 | HIRES | Store a 0 here to select highresolution graphics. |

Space limitations prohibit a discussion in this book of the power of highresolution graphics. The Apple II documentation, however, provides an excellent step-by-step guide to the design, display, saving, and loading of high-resolution images. I must stress, however, that the software in this book expects the host system to have low-resolution graphics, so you'd better tell your Apple II to have lowresolution graphics. The software in this book uses the Apple's low-resolution graphics with text page 1 as the screen memory. To select this display page, simply press the RESET button on your Apple. If, on the other hand, you wish to select this display page under software control, you can do it by calling the subroutine LORES1:

| LORES1 | PHP | Save processor flags. |
| :--- | :--- | :--- |
|  | PHA | Save accumulator. |
|  | LDA \#0 | Store a o in |
|  | STA LOWSCR | LOWSCR to select Page 1, |
|  | STA LORES | and in LORES to select low-resolution |
|  |  | graphics. |
|  | PLA | Restore accumulator. |
|  | PLP | Restore processor flags. |
|  | RTS | Return to caller. |

This subroutine will select low-resolution graphics and text page 1. It preserves all flags and registers, and is completely relocatable.

Even when you've configured your Apple II to low-resolution graphics, your job isn't done. The low-res display of the Apple II is mapped in an unusual manner. For any other system you can assume that the address of a given location on the screen is simply the address of the location above it, plus some row increment. On the Apple II this is not always true. See Table B3.3, Apple II low-res display memory map.

Table B3.3: Apple II low-resolution display.
Page 1

| Row <br> Number | Address of <br> Leftmost Column | Address of <br> Rightmost Column |
| :--- | :--- | :--- |
| $\$ 00$ | $\$ 400$ | $\$ 427$ |
| $\$ 01$ | $\$ 480$ | $\$ 4 \mathrm{~A} 7$ |
| $\$ 02$ | $\$ 500$ | $\$ 527$ |
| $\$ 03$ | $\$ 580$ | $\$ 5 \mathrm{~A} 7$ |
| $\$ 04$ | $\$ 600$ | $\$ 627$ |
| $\$ 05$ | $\$ 680$ | $\$ 6 \mathrm{~A} 7$ |
| $\$ 06$ | $\$ 700$ | $\$ 727$ |
| $\$ 07$ | $\$ 780$ | $\$ 7 \mathrm{~A} 7$ |
|  |  |  |
| $\$ 08$ | $\$ 428$ | $\$ 44 \mathrm{~F}$ |
| $\$ 09$ | $\$ 4 \mathrm{AB}$ | $\$ 4 \mathrm{CF}$ |
| $\$ 0 \mathrm{~A}$ | $\$ 528$ | $\$ 54 \mathrm{~F}$ |
| $\$ 0 B$ | $\$ 5 \mathrm{~A} 8$ | $\$ 5 \mathrm{CF}$ |
| $\$ 0 \mathrm{C}$ | $\$ 628$ | $\$ 64 \mathrm{~F}$ |
| $\$ 0 \mathrm{D}$ | $\$ 6 \mathrm{~A} 8$ | $\$ 6 \mathrm{CF}$ |
| $\$ 0 \mathrm{E}$ | $\$ 728$ | $\$ 74 \mathrm{~F}$ |
| $\$ 0 \mathrm{~F}$ | $\$ 7 \mathrm{~A} 8$ | $\$ 7 \mathrm{CF}$ |
|  |  |  |
| $\$ 10$ | $\$ 450$ | $\$ 477$ |
| $\$ 11$ | $\$ 4 \mathrm{D}$ |  |
| $\$ 12$ | $\$ 550$ | $\$ 4 \mathrm{~F} 7$ |
| $\$ 13$ | $\$ 5 \mathrm{DO}$ | $\$ 577$ |
| $\$ 14$ | $\$ 650$ | $\$ 5 \mathrm{~F} 7$ |
| $\$ 15$ | $\$ 6 \mathrm{DO}$ | $\$ 677$ |
| $\$ 16$ | $\$ 750$ | $\$ 6 \mathrm{~F} 7$ |
| $\$ 17$ | $\$ 7 \mathrm{D} 0$ | $\$ 777$ |
|  |  | $\$ 7 \mathrm{~F} 7$ |

Page 2

| Row Number | Address of Leftmost Column | Address of Rightmost Column |
| :---: | :---: | :---: |
| \$00 | \$800 | \$827 |
| \$01 | \$880 | \$8A7 |
| \$02 | \$900 | \$927 |
| \$03 | \$980 | \$9A7 |
| \$04 | \$A00 | \$A27 |
| \$05 | \$A80 | \$AA7 |
| \$06 | \$B00 | \$B27 |
| \$07 | \$B80 | \$BA7 |
| \$08 | \$828 | \$84F |
| \$09 | \$8A8 | \$8CF |
| \$0A | \$928 | \$94F |
| \$0B | \$9A8 | \$9CF |
| \$0C | \$A28 | \$A4F |
| \$0D | \$AA8 | \$ACF |
| \$0E | \$B28 | \$B4F |
| \$0F | \$BA8 | \$BCF |
| \$10 | \$850 | \$877 |
| \$11 | \$8D0 | \$8F7 |
| \$12 | \$950 | \$977 |
| \$13 | \$9D0 | \$9F7 |
| \$14 | \$A50 | \$A77 |
| \$15 | \$AD0 | \$AF7 |
| \$16 | \$B50 | \$B77 |
| \$17 | \$BD0 | \$BF7 |

Note that the display addresses do not increase uniformly as we move down, row-by-row, through low-res display page 1 or 2 . The addresses increase uniformly from row 0 thru row 7 , but from row 7 to row 8 the display addresses do not increase; they decrease! Then they increase uniformly through line $\$ 0 \mathrm{~F}$ ( 15 decimal), but from line $\$ 0 \mathrm{~F}$ to line $\$ 10$ ( 15 to 16 decimal), the display address plummets again. Then from row $\$ 10$ to row $\$ 17$ ( 16 thru 23) the display addresses again increase uniformly.

If you'd like to take a visual tour of the Apple II's low-res display memory, run the BASIC program in listing B3.1. This program will simply poke a blank into each address in low-res display page 1, starting at the lowest address and moving to the highest address. You'll see that the screen does not fill with blanks in a contiguous manner, but follows a pattern of three interleaved parts.

Listing B3.1: APPLE II low-resolution display, memory-mapper program.
100 REM APPLE II LOW-RESOLUTION DISPLAY, MEMORY-MAPPER
105 REM
108 REM BY KEN SKIER
110 REM
120 FIRST $=1024$ : REM START OF LOW-RESOLUTION PAGE 1.
130 LAST = 2043: REM END OF LOW-RESOLUTION PAGE 1.
140 CHAR $=32$ : REM CHARACTER TO BE POKED INTO SCREEN
150 REM
160 REM
170 FOR X = FIRST TO LAST
175 REM FOR EACH ADDRESS IN LOW-RESOLUTION PAGE 1.
180 POKE X,CHAR
185 REM POKE A WHITE BLANK. THEN,
190 GOSUB 1000: REM WAIT A MOMENT...
200 NEXT X: REM BEFORE POKING NEXT ADDRESS.
210 END
220 REM
230 REM
1000 FOR WAIT $=0$ TO 100
1005 REM THIS IS A WAIT SUBROUTINE.
1010 NEXT WAIT: REM IT SLOWS DOWN PROGRAM SO YOU
1020 RETURN: REM CAN FOLLOW THE ACTION.
Must we now write a whole new set of display procedures to accommodate the unusual mapping of the Apple II low-res display pages? We could. But the screen utilities presented in Chapter 5 will work for the Apple II if we think of the Apple low-res screen as three separate screens: the top eight rows are one screen, the middle eight rows are another screen, and the bottom eight rows are a third screen. Each of these "screens" has a set of screen parameters.

The sceen utilities in this book will work fine if you limit their scope to a given third of the screen. Use TVTOXY only to set a relative screen position within the third of the screen that you have selected. Use the screen utilities only for the top third of the screen. The middle and bottom thirds of the screen may still be used by the PRINT utilities.

To limit the screen utilities to the top third of low-res display page 1, initialize the screen parameters as follows:

SCREEN .WORD \$0400
TVCOLS .BYTE \$27
TVROWS .BYTE $\$ 07$
ROWINC .BYTE $\$ 80$

If you want to keep text from scrolling into the upper third of the screen, store $\$ 08$ in address $\$ 0022$. (In BASIC you may do this with the command POKE 34,8.)

There's one more quirk to the Apple display. If you store an ASCII character in display memory, then you will display a blinking or inverse version of the character. Setting bit 7 in an ASCII character code will cause that character to be displayed in normal mode (a white character on a black background), rather than as a black character on a white background or as a blinking character.

You may experiment with this feature of the Apple II by using the Apple II monitor to store $\$ 41$ (an ASCII " A ") in a location in low-res display page 1. You'll see a blinking "A." Now store \$C1 in a location in low-res display page 1. You'll see a normal "A." Why? Because \$C1 is $\$ 41$ with bit 7 set. To understand what's happening here, look at the Apple II's character set given in Table B3.4.

Table B3.4: The Apple II character set.
RIGHT NYBBLE OF CHARACTER
$-0-1-2-3-4-5-6-7-8-9-\mathrm{A}-\mathrm{B}-\mathrm{C}-\mathrm{D}-\mathrm{E}-\mathrm{F}$
LEFT NYBBLE
OF CHARACTER

| 0- | @ | A | B | C | D | E | F | G | H | I | J | K | L | M | N | O |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1- | P | Q | R | S | T | U | V | W | X | Y | Z |  | $\backslash$ | ] |  | -- |
| 2- |  | Q | " | \# | \$ | \% | , | $($ | ) | * | + |  | - | . | 1 |  |
| 3- | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | : | ; | $<$ | = | > | $?$ |
| 4- | @ | A | B | C | D | E | F | G | H | I | J | K | L | M | N | O |
| 5- | P | Q | R | S | T | U | V | W | X | Y | Z | [ | $\backslash$ | ] |  | -- |
| 6- |  | ! | " | \# | \$ | \% | , | $($ | ) | * | + | , | - | . | 1 |  |
| 7- | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | : | ; | $<$ | $=$ | > | $?$ |
| 8- | @ | A | B | C | D | E | F | G | H | I | J | K | L | M | N | O |
| $9-$ | P | Q | R | S | T | U | V | W | X | Y | Z |  | $\backslash$ | ] |  | -- |
| A- |  | $!$ | " | \# | \$ | \% | , | ( | ) | * | + |  | - |  | 1 |  |
| B- | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | : | ; | $<$ | $=$ | $>$ | ? |
| C- | @ | A | B | C | D | E | F | G | H | I | J | K | L | M | N | O |
| D- | P | Q | R | S | T | U | V | W | X | Y | Z |  | $\backslash$ | ] |  | - |
| E- |  | $!$ | " | \# | \$ | \% | , | ( | ) | * | + |  | - |  | 1 |  |
| F- | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | : | ; | $<$ | = | > | ? |

The Apple II really has only 64 characters in its character set, but it has four ways of displaying each character. Thus, the table shows a set of characters at $\$ 00$ thru $\$ 3 \mathrm{~F}$; the same characters, in the same sequence, appear again at $\$ 40$ thru $\$ 7 \mathrm{~F}$, at $\$ 80$ thru $\$ \mathrm{BF}$, and at $\$ \mathrm{CO}$ thru $\$ \mathrm{FF}$. These represent what I call the first, the second, the third, and the fourth quadrants of the character set.

Character codes in this first quadrant ( $\$ 00$ thru $\$ 3 \mathrm{~F}$ ) will be displayed in reverse video: as black characters on a white background. Character codes in the second quadrant ( $\$ 40$ thru $\$ 7 \mathrm{~F}$ ) will be displayed in a blinking mode. Character codes in the third and fourth quadrants ( $\$ 80$ thru $\$ B F$ and $\$ C 0$ thru $\$$ FF) will be displayed in normal mode: as white characters on black background.

Before we store any ASCII character in screen memory, we must first call FIXCHR, to convert, if necessary, the ASCII character to the host system's corresponding display code. In the Apple II, FIXCHR is very simple:

## FIXCHR <br> ORA \#\$80

## RTS

Set bit 7, so character will be displayed in normal mode.
Return appropriate display code to caller.

## I/O Vectors

The Apple II has a subroutine in read-only memory to get a character from the keyboard, and another subroutine to print a character on the screen. However, the key-in routine at $\$$ FD35 does not return an ASCII code when you press the key for an ASCII character; instead, it returns the appropriate ASCII code with bit 7 set. Similarly, the screen-printing routine at $\$$ FBFD will print an ASCII character to the screen, but the character will be in reverse video or blinking. In order to print an ASCII character to the screen, you must first set bit 7 and then call \$FBFD. Conversely, to get an ASCII character from the keyboard, you must first call \$FD35 and then clear bit 7. Therefore, the following patches are offered:

## Subroutine to Print an ASCII Character to Apple II Screen

APLTVT ORA \#\$80
JSR \$FBFD RTS

Set bit 7 in the ASCII code. Call the ROM screen printer. Return to caller, now that ASCII character originally in accumulator has been printed to screen in normal mode.

## Subroutine to Get an ASCII Character from Apple II Keyboard

Get ASCII character from keyboard with bit 7 set. (Note: you may call \$FD35 instead of calling \$FDOC.)

ORA \#\$80
RTS

Clear bit 7 , leaving the accumulator holding a conventional ASCII code. Return to caller, bearing ASCII character code for depressed key.

## Apple II System Data Block

The I/O vectors ROMTVT and ROMKEY should be initialized to point to APLTVT and APLKEY, respectively. This has been done in the Apple II system data block. You must enter the Apple II system data block into your system's memory if any of the software in this book is to run on your Apple II. See Appendices C15 and E14.

# Appendix B4: The Atari 800 

## Screen

The Atari 800 microcomputer has the most flexible - and, perhaps the most confusing - video-display hardware of any system discussed in this book. Unlike the other systems, almost any portion of the Atari computer's memory may be mapped to the screen. Furthermore, there are many different screen-display modes. When the Atari computer is powered-up, the screen is in text mode zero. That's comparable to the Apple II's low-resolution graphics and text display, which is comparable to the only video-display mode available on the Ohio Scientific or PET computers.

The Atari computer makes other screen modes available to the programmer, but the software in this book assumes a low-resolution text display, so you'd better leave your Atari in screen mode zero if you expect to see any of the displays driven by the software in this book. In other words, if you change the screen mode, the Visible Monitor may well become invisible.

I mentioned that the screen buffer may be almost anywhere in memory. If that's true (and it is), how can you determine the HOME address upon which all the displays in this book are based? It's easy. A pointer at $\$ 58, \$ 59$ ( 88,89 decimal) points to the lowest address in screen memory: the address we refer to as HOME. Before running any of the software in this book, you must set HOME properly for your system. Simply set HOME equal to the value of that pointer. HIPAGE, the value of the highest page in screen memory, is equal to (the high byte of HOME) plus three.

Once we've set HOME and HIPAGE properly, we're home free. The other screen parameters are fixed:

| ROWINC | .BYTE 40 |
| :--- | :--- |
| TVCOLS | .BYTE 39 |
| TVROWS | .BYTE 23 |
| SPACE | .BYTE \$20 |
| ARROW | .BYTE $\$ 7 B$ |

Note that the top of screen memory is always at the top programmable memory, so if you add more programmable memory to your Atari 800, you'll move the screen memory up higher in the address space.

## Proper Display of ASCII Characters

Like the PET, and to a lesser extent the APPLE II, the Atari screen requires that we perform a conversion before we can properly display an ASCII character on the screen. To determine the nature of this conversion, let us first look at the ATARI character set in Table B4.1.

Table B4.1: The Atari character set ATASCI.


A quick examination shows that ASCII characters $\$ 20$ thru $\$ 5 \mathrm{~F}$ are ATASCI (Atari's character set) characters $\$ 00$ thru $\$ 3 F$. Thus, if an ASCII character is in the range of $\$ 20$ thru $\$ 5 \mathrm{~F}$, we can convert it to the appropriate ATASCI character simply by subtracting $\$ 20$.

Further inspection reveals that ASCII characters $\$ 61$ thru \$7A correspond to ATASCI characters $\$ 61$ through $\$ 7$ A. Thus, if an ASCII character is in the range of $\$ 61$ thru $\$ 7 \mathrm{~A}$, it needs no conversion to ATASCI; it already is the corresponding ATASCI character.

Finally, if an ASCII character is not in the range $\$ 20$ thru $\$ 5 \mathrm{~F}$ or $\$ 61$ thru $\$ 7 \mathrm{~A}$, it's not a printable character and has no agreed-upon graphic representation. For those cases we'll just leave them alone.

Figure B4.1 flow-charts this algorithm.


Figure B4.1: Flowchart of routine to convert an ASCII character for display on Atari screen.

Using the flowchart in figure B4.1 as a guide, we can write source code for FIXCHR , which takes an ASCII character as input and returns an Atari display code so that the character may be properly displayed on the video screen.

## FIXCHR

FIXCHR AND \#\$7.
SEC
CMP \#\$20

Clear bit 7 so character is a legitimate ASCII character.
Prepare to compare. Character less than $\$ 20$ ?

|  | BCC BADCHR | If so, it's not a printable ASCII <br> character, so return a blank. |
| :--- | :--- | :--- |
|  | CMP \#\$60 | Character less than $\$ 60$ ? |
| BCC SUB $\$ 20$ | If so, subtract $\$ 20$ and return. |  |
|  | CMP \#\$7B | Character less than $\$ 7 B ?$ |
| BCC EXIT | If so, return with the character. <br> If not less than $\$ 7 B$, <br> the character is not a printable ASCII <br> character, so return a blank. |  |
| EXIT | RTS | Subtract $\$ 20$ and <br> SUB $\$ 20$ |
|  | SBC $\# \$ 20$ | return. |

## Keyboard Input

If no key has been pressed, then address $\$ 02 \mathrm{FC}$ ( 764 decimal) contains $\$$ FF. But whenever you depress a key on the Atari keyboard - even if a program is not scanning the keys - an electronic circuit will sense that a key has closed and will store the hardware code for that key in address $\$ 02 \mathrm{FC}$. However, the code in $\$ 02 \mathrm{FC}$ will be a hardware code, not obviously related to ASCII or ATASCI.

Table B4.2: Atari Hardware Key-Codes.

| Hex | Decimal | Key | Hex | Decimal | Key |
| :---: | :---: | :---: | :---: | :---: | :---: |
| \$00 | 0 | L | \$20 | 32 | , |
| 1 | 1 | J | 1 | 33 | SPACE |
| 2 | 2 | ; | 2 | 34 | - |
| 3 | 3 |  | 3 | 35 | N |
| 4 | 4 |  | 4 | 36 |  |
| 5 | 5 | K | 5 | 37 | M |
| 6 | 6 | + | 6 | 38 | / |
| 7 | 7 | * | 7 | 39 | ATARI |
| 8 | 8 | 0 | 8 | 40 | R |
| 9 | 9 |  | 9 | 41 |  |
| A | 10 | P | A | 42 | E |
| B | 11 | U | B | 43 | Y |
| C | 12 | RETURN | C | 44 | TAB |
| D | 13 | I | D | 45 | T |
| E | 14 | - | E | 46 | W |
| F | 15 | $=$ | F | 47 | Q |


| $\$ 10$ | 16 | V | $\$ 30$ | 48 | 9 |
| ---: | :---: | :---: | ---: | :---: | :---: |
| 1 | 17 |  | 1 | 49 |  |
| 2 | 18 | C | 2 | 50 | $\varnothing$ |
| 3 | 19 |  | 3 | 51 | 7 |
| 4 | 20 |  | 4 | 52 | BACK S |
| 5 | 21 | B | 5 | 53 | 8 |
| 6 | 22 | X | 6 | 54 | $<$ |
| 7 | 23 | Z | 7 | 55 | $>$ |
| 8 | 24 | 4 | 8 | 56 | F |
| 9 | 25 |  | 9 | 57 | H |
| A | 26 | 3 | A | 58 | D |
| B | 27 | 6 | B | 59 |  |
| C | 28 | ESC | C | 60 | LOWR |
| D | 29 | 5 | D | 61 | G |
| E | 30 | 2 | E | 62 | S |
| F | 31 | 1 | F | 63 | A |

The Hex and Decimal Columns give the low 6 bits of the hardware key-code stored in address \$02FC ( 764 decimal) when the given keys are pressed. Either SHIFT key sets bit 6. CTRL key sets bit 7 .

In order to convert that hardware code to ASCII, we need to understand its nature. The six low-order bits of the hardware key-code uniquely identify the key. (See Table B4.2.) Bits 6 and 7 identify its shift state. Bit 6 is set if the key is typewriter-shifted; bit 7 is set if the key is control-shifted. The key is typewritershifted if either SHIFT key is down; the CAPS/LOWR key has no effect on the typewriter-shift state as reflected in the hardware key-code. The keyboard is control-shifted if the CTRL key is down.

If you don't care about the keyboard's shift state, but merely want to determine which physical key has been pressed, then you can clear the two high-order bits in the hardware key-code and you'll be left with a number from 0 to 63 decimal ( 00 to $\$ 3 F$ ) uniquely identifying the key most recently depressed. If you care about the keyboard's typewriter-shift state but are indifferent to its control-shift state, then you can clear bit 7 in the hardware key-code and you'll be left with a number from 0 to 127 decimal ( 00 to $\$ 7 \mathrm{~F}$ ), which means the keyboard can generate twice as many characters as it has physical keys. To enable control-shifting, simply preserve the hardware key-code, and you double once again the number of characters that the keyboard (and hence the user) may generate.

Since the simple text editor presented in Chapter 11 assigns certain functions to control-shifted keys, and since you never know when you might need some additional character codes from your keyboard, Appendix C16 presents a key-handling subroutine for the Atari. This subroutine is capable of generating different
characters in each of the four different shift-states (unshifted, typewriter-shifted, control-shifted, typewriter- and control-shifted).

It's a simple matter to use the eight-bit hardware keycode as an index into a keyboard definition table. For any given hardware key-code, we may assign any character we like. The keyboard definition table presented in Appendix C16 assigns standard ASCII characters to all letter, number, and punctuation keys, in both the unshifted and typewriter-shifted states. Other keys are assigned values consistent with their expected use by the software in this book (eg: Control-P generates a $\$ 10$, thus making it a PRINT key in the eyes of the simple text editor). All keys and shift states that have no special meaning to this software have been assigned character codes of zero; feel free to change these character codes to any values you desire.

Assuming that we have in memory a keyboard definition table called ATRKYS, we can get an ASCII character from the Atari keyboard with the following subroutine, ATRKEY:

| ATRKEY | LDA \$02FC |
| :--- | :--- |
|  | CMP \#\$FF |
|  | BEQ ATRKEY |

TAY
LDA ATRKYS,Y
RTS

Has a key been depressed?
$\$$ FF means no key. If not, look again. A key has gone down and the accumulator holds its hardware key-code.
Prepare to use that code as an index.
Look up character for that key and shift state.
Return with ASCII character corresponding to that key and shift state.

## Print a Character to the Screen

The Atari 400 and 800 computers each provide a powerful I/O (input/output) routine which allows the programmer to get characters from virtually any source, and to send characters to virtually any device - the screen, the printer, the cassette recorder, and the disk. But, as in the case of Atari's varied screen modes, power breeds complexity. I have found it easier to substitute my own simple routine to print a character on the TV screen, bypassing the Atari I/O routines entirely.

Incidentally, this routine will work with any 6502 -based computer that has a low-resolution memory-mapped display. If you need a simple TVT simulator for your home-brew 6502-based system with a video display, TVTSIM might meet your needs. In any event, it prints characters to the screen, and avoids the necessity of plumbing the depths of the many modes and data structures associated with Atari's central I/O routine.

With your system data block initialized as shown in Appendices C16 and E15 (which includes the TVT simulator as the subroutine to print characters to the screen), you are almost ready to run the software in this book on your own system.

## Setting the Top Of Memory

Address $\$ 2 \mathrm{E} 6$ ( 742 decimal) holds the number of pages of RAM available to the BASIC interpreter. Store a \$0D (13 decimal) in that location and BASIC will use memory up to $\$ 0 D F F$, but will not use $\$ 0 \mathrm{E} 00$ and up.

NOTE: On the Atari, the software in this book uses memory from \$0E80 to \$1FFF, which is the address space required by the ATARI DOS (Disk Operating System) and the ATARI RS-232 serial interface, so you may not use DOS or RS-232 if you expect to use the software in this book. However, there should be no conflict between software in this book and the cassette-based Atari 800.

Thus, we may set the top of memory with the following BASIC command:
POKE 742,13
When you have used the OBJECT CODE LOADER to READ and POKE object code from all the appropriate $E$ appendices into your Atari computer, run the following BASIC program. It will initialize screen parameters and the top of memory, and then pass control to the Visible Monitor.

100 REM
110 REM
120 REM
130 REM
140 REM
$150 \mathrm{LO}=\operatorname{PEEK}(88)$ :
$160 \mathrm{HI}=\operatorname{PEEK}(89)$ : OR DISASSEMBLER"
170 POKE 4096,LO: REM Set Low byte of HOME.
180 POKE 4097, HI: REM Set High byte of HOME. 200 REM

220 POKE 742,13:
230 REM
240 REM
$250 X=U S R(4615)$ :
260 END

165 IF HI < 32 THEN PRINT "ON AN 8 K ATARI YOU MAY NOT USE EDITOR

190 POKE 4101,HI + 3: REM Set HIPAGE = Highest page in screen memory.
210 REM Now set the top of memory available to BASIC.
Visible Monitor Start-Up Program for the Atari.
First, set the screen parameters.
A pointer at 88,89 points to lowest screen address.
REM Set LO to the low byte of HOME.
REM Set HI to the high byte of HOME.

Tell BASIC to use only memory up to \$0DFF.
Now call the Visible Monitor.
REM Call the Visible Monitor as a subroutine.

# Appendix CI: 

## Screen Utilities



| 590 |  | ; |  | POSITION IN UPPER LEFT CORNER. |
| :---: | :---: | :---: | :---: | :---: |
| E00 |  | ; |  |  |
| 610 | $1002=$ | ROWINC=PARAMS +2 |  |  |
| 620 |  | ; |  | ROWINC IS A EYTE GIUING |
| 630 |  | ; |  | RDDRESS DIFFERENCE FROM ONE |
| 540 |  | ; |  | ROW TO THE NEXT. |
| 551 |  | ; |  |  |
| 563 | $1003=$ | TUCOLS=FARAMS +3 |  |  |
| 670 |  | ; |  | TUCOLS IS A BYTE GIUING |
| Ect |  | ; |  | NUMEER OF COLUMNS ON SCREEN. |
| 693 |  | ; |  | (COUNTING FROM ZERO.) |
| 700 |  | ; |  |  |
| 710 | 1004= | TUROWS=PARAMS +4 |  |  |
| 720 |  | ; |  | TUROWS IS A BYTE GIUING |
| 730 |  | ; |  | NUMEER OF ROWS ON SCREEN, |
| 749 |  | ; |  | (COUNTING FROM ZERO.) |
| 750 |  | ; |  |  |
| 750 | 1005= | HIFAGE=PARAMS +5 |  |  |
| 770 |  | ; |  | Hiffge is the high byte of |
| 780 |  | ; |  | THE HIGHEST RDDRESS ON SCREEN. |
| 730 |  | ; |  |  |
| 806 |  | ; |  |  |
| 810 | $1006=$ |  | ELANK=PRRAMS + E | YOUR SYSTEM 5 CHARACTER |
| 820 |  | ; |  | CODE FOR A BLANK. |
| 830 |  | ; |  |  |
| 840 | $1007=$ |  | ARROW=PARAMS+7 | YOUR SYSTEM'S CHARRCTER |
| 850 |  | ; |  | FOR AN UP-GRRROW. |
| 6ED |  | ; |  |  |
| 878 | 1611= |  | FIXCHR=PARAMS + | \$11 |
| 830 |  | ; |  | FIXCHR IS A SUBROUTINE THAT |
| ESO |  | ; |  | RETURNS YOUR SYSTEM'S |
| 509 |  | ; |  | DISPLAY CODE FOR ASCII. |
| 910 |  | ; |  | CODE. |
| 929 |  | ; |  |  |
| 930 |  | ; |  |  |
| 540 |  | ; |  |  |
| 951 |  | ; |  |  |
| 960 |  | ; |  |  |
| 970 | 1100 |  | * $=$ \$1100 |  |
| 880 |  | ; |  |  |
| 950 |  | ; |  |  |
| 1800 |  | ; |  |  |
| 1010 |  | ; |  |  |
| 1020 |  | ; |  |  |
| 1030 |  | ; |  |  |
| 1040 |  | ; |  |  |
| 1050 |  | ; |  |  |
| 1060 |  | ; |  |  |
| 1070 |  | ; |  |  |
| 1080 |  | ; | CLEAR SCR | EEN |
| 1090 |  | ; |  |  |
| 1100 |  | ; |  |  |
| 1110 |  | ; |  |  |
| 1120 |  | ; |  |  |
| 1130 |  | ; |  |  |
| 1140 |  | ; |  | . |
| 1150 |  | ; |  |  |
| 1150 |  | ; |  |  |


| 1176 |  | ; | CLEAR SCREEN. | PRESERUING THE ZERO. PRGE. |
| :---: | :---: | :---: | :---: | :---: |
| 1180 |  | ; |  |  |
| 1190 |  | ; |  |  |
| 12010 |  | ; |  | . |
| 1210 |  | ; |  |  |
| 1220 | 1100200411 | CLR.TU | JSR TUPUSH | SAUE ZERO PAGE bytes that |
| 1230 |  | ; |  | WILL BE CHANGED. |
| 1240 | 1103202811 |  | JSR TUHOME | SET SCREEN LOCATIOM TO UPPER |
| 1250 |  | ; |  | LEFT CORNER OF THE SCREEM. |
| 1250 | 1106 feg3io |  | LDX TUCOLS | LOAD $X, Y$ REGISTERS WITH |
| 1270 | 1109 ACD410 |  | LDY TUROWS | $X, Y$ DIMENSIONS OF SCREEN. |
| 1280 | 1100201311 |  | JSR CLR. XY | CLEAR $\times$ COLUMNS; Y ROMS |
| 1290 |  | - |  | FROM CURRENT SCREEN LOCATION. |
| 1300 | 110F 200311 |  | JSR TU.POP | RESTORE ZERO PAGE BYTES THAT |
| 1310 |  | ; |  | WERE CHANGED. |
| 1320 | 111260 |  | RTS | RETURM TO CRLLER, WITH ZERO |
| 1330 |  | ; |  | PAGE PRESERUED. |
| 1343 |  | ; |  |  |
| 1350 |  | ; |  |  |
| 1360 |  | ; |  |  |
| 1370 |  | ; |  |  |
| 1360 |  | ; |  |  |
| 1350 |  | ; |  |  |
| 1400 |  | ; |  |  |
| 1410 |  | ; |  |  |
| 1420 |  | ; |  |  |
| 1430 |  | ; |  |  |
| 1440 |  | ; |  |  |
| 1450 |  | ****** |  |  |
| 1450 |  | ; |  |  |
| 1470 |  | ; | CLEAR PORT | TION OF SCREEM |
| 1460 |  | ; |  |  |
| 1450 |  |  |  |  |
| 1590 |  | ; |  |  |
| 1516 |  | ; |  |  |
| 1520 |  | ; |  |  |
| 1530 |  | ; |  |  |
| 1540 |  | ; |  | Clear $\times$ COLUMNS, Y ROWS |
| 1550 |  | ; |  | FROM CURRENT SCREEN LOCATION. |
| 1563 |  | ; |  | MOUES TU.PTR DOWN BY Y ROWS.' |
| 1579 |  | ; |  |  |
| 159 |  | ; |  |  |
| 1590 |  | ; |  |  |
| 1500 | 1113 gezall | CLR.XY | STX COLS | SET THE NUMEER OF COLUMNS |
| 1¢10 |  | ; |  | TO BE CLEARED. |
| 1520 | 111698 |  | TYA |  |
| 1639 | 1117 AA |  | TAK | NOW $\times$ HOLDS NUMEER OF ROWS |
| 1640 |  | ; |  | TO BE CLEfiRED. |
| 1650 |  | ; |  |  |
| 165ด | 1118 ADEE10 | CLRROW | LDA BLANK | WE' LL CLEAR THEM BY |
| 1670 |  | ; |  | WRITING BLAMKS TO THE |
| 1ESQ |  | ; |  | SCREEN. |
| 1630 | 111E AC2A11 |  | LDY COLS | LOED Y WITH NUMEER OF |
| 1706 |  | ; |  | COLUMTS TO PE CLEARED. |
| 1710 | 111E 9100 | CLRFOS | STA (TU.PTR), Y | CLEAR A POSITION BY |
| 1729 |  | ; |  | WRITING A BLANK INTO IT. |
| 1730 |  | ; |  |  |
| 1740 | 112088 |  | DEY | ADJUST INDEX FOR NEXT |



| 2330 |  | ; |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 2340 | 1132 ADO 418 | CENTER | LIA TUROWS | LOAD A WITH TOTAL ROWS. |
| 2350 | 1135 4A |  | LSR A | DIUIDE IT BY TWO. |
| 2360 | 1135 ค8 |  | TAY | $Y$ NOW HOLDS THE NUMBER OF |
| 2370 |  | ; |  | THE SCREEN' 5 CEMTRAL ROW. |
| 2380 |  | ; |  |  |
| 2350 | 1137 ADD310 |  | LDA TUCOLS | LOAD A WITH TOTAL COLUMNS. |
| 2400 | 113月 4A |  | LSR A | DIUIDE IT BY TWO. |
| 2410 | 1138 AP |  | TAX | $X$ NOW HOLDS THE NUMBER OF |
| 2420 |  | ; |  | THE SCREEN' 5 CENTRAL COLUMN. |
| 2430 |  | ; |  |  |
| 2440 |  | ; |  |  |
| 2450 |  | ; |  | $X$ RND $Y$ REGISTERS NOW HOLD |
| 2460 |  | ; |  | $X, Y$ COORDINATES OF CENTER |
| 2470 |  | ; |  | OF SCREEN. |
| 2480 |  | ; |  |  |
| 2450 |  | ; |  | 50 NOW LET'S SET THE SCREEN |
| 2509 |  | ; |  | LOCATION TO THOSE X,Y |
| 2510 |  | ; |  | COORDINATES: |
| 2520 |  | ; |  |  |
| 2530 |  | ; |  |  |
| 2540 |  | ; |  |  |
| 2550 |  | ; |  |  |
| 2550 |  | ; |  |  |
| 2570 |  | ; |  |  |
| 2580 |  | ; |  |  |
| 2550 |  | ; |  |  |
| 2600 |  | ; |  |  |
| 2610 |  | ; **** | ******************* |  |
| 2620 |  | ; |  |  |
| 2630 |  | ; | TUTOXY |  |
| 2540 |  | : |  |  |
| 2650 |  | ; ***** |  |  |
| 2550 |  | ; |  |  |
| 2579 |  | ; |  |  |
| 2680 |  | ; |  |  |
| 2550 |  | ; |  |  |
| 2700 |  | ; |  |  |
| 2710 | 113038 | TUTOXY | SEC | SET CURRENT SCREEN LOCATION |
| 2720 |  | ; |  | TO COORDINATES GIUEN BY |
| 2730 |  | ; |  | THE $X$ AIND Y REGISTERS. |
| 2740 |  | ; |  |  |
| 2750 | 1130 Ecas 10 |  | CPX TUCOLS | IS $X$ OUT OF RFANGE? |
| 2769 | 1140 9003 |  | BCC $\times$. OK | If NOT, LERUE IT ALONE. |
| 2770 |  | ; |  | IF $X$ IS OUT OF RANGE, GIUE |
| 2780 | 1142 AEg310 |  | LDX TUCOLS | IT ITS HIGHEST LEGAL URLUE. |
| 2790 |  | ; |  | NOW $X$ IS LEGRL. |
| 2800 |  | ; |  |  |
| 2810 | 114538 | X.OK | SEC | IS $Y$ OUT OF RANGE? |
| 2820 | 1146 CC0410 |  | CPY TUROWS |  |
| 2830 | 1149 9003 |  | BCC Y.OK | If not, lefve it flone. |
| 2840 |  | ; |  |  |
| 2850 |  | ; |  | If $Y$ IS OUT Of RANGE, GIUE |
| 2860 | 1148 ACO 410 |  | LDY TUROWS | $Y$ ITS HIGHEST LEGRL UALUE. |
| 2870 |  | ; |  | NOW Y IS LEGRL. |
| 2800 |  | ; |  |  |
| 2890 |  | ; |  |  |
| 2500 | 114E RDODIG | Y.OK | LDA HOME | SET TU.PTR = LOWEST SCREEM |


| 2910 | 11518500 |  | STA TU.PTR | ADDRESS. |
| :---: | :---: | :---: | :---: | :---: |
| 2320 | 1153 ADO110 |  | LDA HOME +1 |  |
| 2930 | 11568501 |  | STA TU.PTR+1 |  |
| 2940 |  | ; |  |  |
| 2550 | 1158 D8 |  | PHP | Sfue caller' 5 decimal flag. |
| 2960 | 1153 DE |  | CLD | CLEAR DECIMAL FOR BINARY |
| 2976 |  | ; |  | ADDITION. |
| 2590 |  | ; |  |  |
| 2930 | 115A 8A |  | T×A | ADD $\times$ TO TU.PTR |
| 3060 | 1158 18 |  | CLC |  |
| 3010 | 11506500 |  | ADC TU.PTR |  |
| 3020 | $115 \pm$ 9003 |  | BCC COLSET |  |
| 3030 | $1160 \mathrm{EGO1}$ |  | INC TU.PTR+1 |  |
| 3040 | 116218 |  | CLC |  |
| 3050 |  | ; |  |  |
| 30.80 |  | ; |  |  |
| 3070 | 1153 C000 | COLSET | CPY ${ }^{\text {¢ }}$ | ADD Y*ROWINC TO TU.PTR: |
| 3080 | 1165 FDGE |  | EEQ TU.SET |  |
| 3090 | 116718 | ADDROW | CLC |  |
| 3100 | 1159 509210 |  | RDC POWINC |  |
| 3110 | 11655002 |  | ECC : +4 |  |
| 3120 | 1160 ESGI |  | INIC TU. PTR +1 |  |
| 3130 | 116F 86 |  | DEY |  |
| 3140 | 1170 D0F5 |  | BNE ADMROW |  |
| 3150 |  | ; |  |  |
| 3169 |  | ; |  |  |
| 3170 | 11728509 | TU.SET | STA TU.PTR |  |
| 3180 | 117426 |  | FLP | RESTORE CALLER' 5 decimal flag |
| 3190 | 117560 |  | RTS | RETURN TO CALLER |
| 3200 |  | ; |  |  |
| 3210 |  |  |  |  |
| 3220 |  | ; |  |  |
| 3230 |  | ; |  |  |
| 3240 |  | ; |  |  |
| 3250 |  | ; |  |  |
| 3250 |  | ; |  |  |
| 3270 |  | ; |  |  |
| 3280 |  | ; |  |  |
| 3290 |  | ; |  |  |
| 3300 |  | ; ***** |  | ****************************** |
| 3310 |  | ; |  |  |
| 3329 |  |  | TUDOWN, | TUSKIP, and TUPLUS |
| 3530 |  | ; |  |  |
| 3345 |  | ; ***** |  |  |
| 3550 |  | ; |  |  |
| 3360 |  | ; |  |  |
| 3370 |  | ; |  |  |
| 3380 |  | , |  |  |
| 3330 |  | ; |  |  |
| 3400 | 1175 ADG210 | TUDOWN. | LDA ROWINC | MOUE TU.PTR DOWN BY ONE ROW. |
| 3410 | 117918 |  | CLC |  |
| 3420 | 11789005 |  | BCC TUPLUS |  |
| 3430 |  | ; |  |  |
| 3440 | 117C 203B11 | UuCHAR | JSR TU.PUT | PIJT CHARACTER ON SCREEN |
| 3450 |  | ; |  | AND THEN |
| 34E0 |  | ; |  |  |
| 3470 | 117F ASOL | TUEKIP | LDA \#1 | SKIP ONE SCREEN LOCATION |
| 3480 |  | ; |  | BY INCREMENTING TU.PTR |




| 4550 | 118E GSg | ; | ADC *6 | IF SO, IT MUST BE A-F. |
| :---: | :---: | :---: | :---: | :---: |
| 4670 |  | ; |  | ADD 36 HEX TO CONVERT IT. |
| 4800 |  | ; |  | TO. CORRESPONDING ASCII CHAR. |
| 459 | 11006930 | DECIML | ADC 4 \# 30 | IF A IS $\square-9$, ADD 36 HEX |
| 4720 |  | ; |  | TO CONUERT IT TO |
| 4712 |  | ; |  | CORRESPONDING ASCII CHAR. |
| 4720 |  | ; |  |  |
| 4730 | 110228 |  | PLF | RESTORE ORIGINAL DECIMAL |
| 4740 |  | ; |  | FLAG, AND |
| 4750 | 1103 E0 |  | RTS | RETURN TO CALLER |
| 4760 |  | ; |  |  |
| 4770 |  | ; |  |  |
| 4780 |  | ; |  |  |
| 4790 |  | ; |  |  |
| 4800 |  | ; |  |  |
| 4810 |  | ; |  |  |
| 4820 |  | ; |  |  |
| 4830 |  | ; |  |  |
| 4840 |  | ; |  |  |
| 4850 |  | ; |  |  |
| 4860 |  | ; |  |  |
| 4870 |  | ; |  |  |
| 4880 |  | ; |  |  |
| 4890 |  | ; ****** |  |  |
| 49010 |  | ; |  |  |
| 4910 |  | ; | TUPLISH |  |
| 4929 |  | ; |  |  |
| 4930 |  | ; ***** |  |  |
| 4540 |  | ; |  |  |
| 4950 |  | ; |  |  |
| 4SEU |  | ; |  |  |
| 4973 |  | ; |  | SAUE CURRENT SCREEN LOCATION |
| 4580 |  | ; |  | ON STACK, FOR CALLER. |
| 4950 |  | ; |  |  |
| 5000 |  | ; |  |  |
| 5010 |  | ; |  |  |
| 5023 |  | ; |  |  |
| 5030 |  | ; |  |  |
| 5040 | $11 C 458$ | TUFUSH | PLA | PULL RETURN RDDRESS FROM |
| 5050 | 1105 AR |  | TAX | STACK AND SAUE IT IN $X$ AND |
| 5060 | 1106 ES |  | FLA | Y REGISTERS. |
| 5070 | 1107 fi |  | TAY |  |
| 5090 |  | ; |  |  |
| 5050 |  | ; |  |  |
| 5100 | 1108 A 501 |  | LDA TU.PTR+1 | GET TU.PTR AND |
| 5110 | 11CA 48 |  | FH9 |  |
| 5120 | 11CB f500 |  | LEA TU.PTR | PUSH IT ONTO THE STACK. |
| 5130 | 11CD 48 |  | PHA |  |
| 5140 |  | ; |  |  |
| 5150 |  | ; |  |  |
| 5150 | 11C5 56 |  | TYA | PLACE RETURN ADDRESS |
| 5170 | 11CF 48 |  | PHA |  |
| 5180 | 110988 |  | TXA | BACK ON STACK. |
| 5190 | 110148 |  | PHA |  |
| 5206 |  | ; |  |  |
| 5210 |  | ; |  |  |
| 5220 | 110260 |  | RTS | THEN RETURN TO CALLER. |


| 5230 | ; |  | CALLER WILL FIND TU.PTR ON |
| :---: | :---: | :---: | :---: |
| 5240 | ; |  | STACK, LOW BYTE ON TOP. |
| 5250 | ; |  |  |
| 5250 | ; |  |  |
| 5270 | ; |  |  |
| 5280 | ; |  |  |
| 5290 | ; |  |  |
| 5300 | ; |  |  |
| 5310 | ; |  |  |
| 5320 | ; |  |  |
| 5330 | ; |  |  |
| 5340 | ; |  |  |
| 5350 | ***** |  | **************************** |
| 5350 | ; |  |  |
| 5370 | ; | TU. POP |  |
| 5380 | ; |  |  |
| 5390 | **** |  |  |
| 5400 | ; |  |  |
| 5410 | ; |  |  |
| 5426 | ; |  |  |
| 5430 | ; |  | RESTORE SCREEN LOCATION |
| 5440 | ; |  | PREUIOUSLY SAUED ON STACK. |
| 5450 | ; |  |  |
| 5460 | ; |  |  |
| 5470 | ; |  |  |
| 54801103 E8 | TU.POP | PLA | FULL RETURN ADDRESS FROM |
| 54301104 AA |  | TfX | STACK, SAUING IT IN X... |
| 5500110568 |  | PLA |  |
| 5510 115S Re |  | TAY | ...fnd In y |
| 5520 | ; |  |  |
| 5530 | ; |  |  |
| 5540 1107 E8 |  | PLA | RESTORE... |
| 5551108 8500 |  | STA TU.PTR | ...TU.PTR |
| 5564 11DR 68 |  | PLA | . FR OM |
| 557011088581 |  | STA TU.PTR+1 | . . STACK. |
| 5580 | ; |  |  |
| 5590 | ; |  |  |
| 5E0] 11DD 96 |  | TYA | FLACE RETURN ADDRESS |
| 5610 11DE 48 |  | PHA | BACK ... |
| 562011 DF EA |  | TXA |  |
| $563011 E 048$ |  | FHA | ...ON STACK. |
| 5640 | ; |  |  |
| 5650 | ; |  |  |
| 5EEO 11E1 6И |  | RTS | RETURN TO CALLER. |

## Appendix C2:

## Visible Monitor (Top Level and Display Subroutines)



| 590 | ; |  |
| :---: | :---: | :---: |
| $6000075=$ |  | RUBOUT $=\$ 7 \mathrm{~F}$ |
| 610 | , |  |
| 620 000n= |  | $C R=\$ 0 D \quad$ ASCII FOR CARRIAGE RETUARH. |
| E30 | ; |  |
| 640 | ; |  |
| 650 | ; |  |
| 66® | ; |  |
| 670 | ; |  |
| 680 | ; |  |
| 698 | ; |  |
| 700 | ; |  |
| 710 | ; |  |
| 720 | ; |  |
| 730 | ; |  |
| 740 | ; |  |
| 758 | ; |  |
| 760 | ; |  |
| 770 | ; | REQUIRED SUBROUTINES |
| 788 | ; |  |
| 790 | ; |  |
| 808 | ; |  |
| 810 | ; |  |
| 820 | ; |  |
| $8301100=$ |  | TUSIJBS $=\$ 1100$ |
| $8401100=$ |  | CLR.TU = TUSUBS |
| 850 1113= |  | CLR.XY $=$ TUSUBS + \$13 |
| 860 1128= |  | TUHOME $=$ TUSUBS+\$2B |
| B70 113C= |  | TUTOXY $=$ TUSUES + \$3C |
| 880 1175= |  | TUDOWN = TUSUBS+\$76 |
| 890 1176 |  | ULCHAR = TUSUBS+47C |
| 900 117F= |  | TUSKIP $=$ TUSUBS+ + \% 7 F |
| 910 1181= |  | TUPLUS $=$ TUSUBS+\$E1 |
| 320 11月3 $=$ |  | UUBYTE $=$ TUSUBS+\$A3 |
| $93011 \mathrm{EG}=$ |  | RSCII = TUSUBS+IBEG |
| 940 1iC4= |  | TUPUSH $=$ TUSUBS + \$5C4 |
| 950 1103= |  | TU.POP $=$ TUSUSS + \$D3 |
| 960 | ; |  |
| 970 | ; |  |
| 980 | ; |  |
| 9901200 |  | * $=$ \#1200 |
| 1000 | ; |  |
| 1010 | ; |  |
| 1020 | ; |  |
| 1030 12E3= |  | UFDATE $=$ *+\$E3 |
| 1040 | ; |  |
| 1050 | ; |  |
| 1860 | ; |  |
| 1070 | ; |  |
| 1080 | ; |  |
| 1090 | ; |  |
| 1100 | ; |  |
| 1110 | ; |  |
| 1120 | ; |  |
| 1130 | ; |  |
| 1140 | ; |  |
| 1150 | ; |  |
| 1150 | ; |  |


| 1170 |  | ; ****************************************** |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 1180 |  | USER-MODIFIABLE DATA |  |  |
| 1190 |  |  |  |  |
| 1200 |  | ; |  |  |
| 1210 |  | ; ******************************************* |  |  |
| 1220 |  | ; |  |  |
| 1230 |  | ; |  |  |
| 1240 |  | ; |  |  |
| 1250 |  | ; |  |  |
| 1260 |  | ; | . BYTE D | Number of Current field. |
| 1270 | 120000 | FIELD |  |  |
| 1280 |  | ; |  | (MUST BE. 0-6.) |
| 1250 |  | ; |  | IMAGE OF ACCUMULATOR. |
| 1300 | 120100 | REG.A | . BYTE D |  |
| 1310 |  | ; |  |  |
| 1320 | 120200 | REG. X | . BYTE 0 | IMAGE OF X-REGISTER. |
| 1330 |  | ; |  |  |
| 1340 | 120300 | REG.Y | . BYTE 0 | IMRGE OF Y-REGISTER. |
| 1350 |  | ; |  |  |
| 1360 | 120400 | REG.P | . BYTE D | IMAGE OF PROCESSOR STATUS |
| 1370 |  | ; |  | REGISTER. |
| 1390 |  | ; |  |  |
| 1390 | $1201=$ | REGS $=$ REG.A |  |  |
| 1400 |  | ; |  | POINTER TO CURRENTLY- |
| 1410 | 12950000 | SELECT | . WORD 0 |  |
| 1420 |  | ; |  |  |
| 1430 |  | ; |  |  |
| 1440 |  | ; |  |  |
| 1450 |  | ; |  |  |
| 1460 |  | ; |  |  |
| 1470 |  | ; |  |  |
| 140 al |  | ; |  |  |
| 1490 |  | ; |  |  |
| 1590 |  | ; |  |  |
| 1510 |  |  |  | ************************** |
| 1520 |  | ; |  |  |
| 1530 |  | ; | THE UISIBLE MONITOR |  |
| 1540 |  | ; |  |  |
| 1550 |  |  |  |  |
| 1550 |  | ; |  |  |
| 1579 |  | ; |  |  |
| 1580 |  | ; |  |  |
| 1590 |  | ; |  |  |
| 1600 | 120768 | UIEMON | $\begin{aligned} & \text { PHP } \\ & \text { CLD } \end{aligned}$ | SAUE CALLER' 5 STATUS FLAGS. CLEAR DECIMAL MODE, SINCE |
| 1510 | 1208 D8 |  |  |  |
| 1620 |  | ; |  | ARITHMETIC OPERATIONS IN THIS BOOK ARE RLINAYS BIMARY. |
| 1630 |  | ; |  |  |
| 1640 |  | ; |  | PUT MONITOR DISPLAY ON |
| 1550 | 1209201212 |  | JSR DSPLAY |  |
| 1650 |  | ; |  | SCREEN. |
| 1670 |  | ; TSR UPDATE |  | GET USER REQUEST AND handle IT. |
| 1680 | 1200205312 |  | JSR UPDATE |  |
| 1690 |  | $\because$ |  |  |
| 1700 | 129F 18 |  | CLC |  |
| 1710 | 1210 90F5 |  | BCC UISMON+1 | LOOP BACK 10 DISPLAY... |
| 1720 |  | ; |  |  |
| 1730 |  | ; |  |  |
| 1740 |  | ; |  |  |



| 2330 | 1230201311 |  | JER CLR. $X Y$ |  | CLEAR X COLUMNS, Y ROWS. |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 2340 |  | ; |  |  |  |
| 2350 | 123350 |  | RTS |  | RETURN TO CALLER. |
| 2366 |  | ; |  |  | . |
| 2370 |  | ; |  |  |  |
| 2380 |  | ; |  |  |  |
| 2390 |  | ; |  |  | \% |
| 2400 |  | ; |  |  |  |
| 2410 |  | ; |  |  |  |
| 2420 |  | ; |  |  |  |
| 2430 |  | ; | $\cdots$ |  |  |
| 2446 |  | ; |  |  |  |
| 2450 |  | ; |  |  |  |
| 2450 |  | ; **** |  | * |  |
| 2470 |  | ; |  |  |  |
| 2480 |  | ; | DISPLAY | AB | EL Line |
| 2450 |  |  |  |  |  |
| 2500 |  | ; **** |  | ** |  |
| 2510 |  | ; |  |  |  |
| 2520 |  | ; |  |  |  |
| 2530 |  | ; |  |  |  |
| 2540 |  | ; |  |  |  |
| 2550 |  | ; |  |  |  |
| 2560 | 1234 A200 | LINE. 1 | LDX \#13 |  | $X$-COORDINATE OF LABEL "A". |
| 2570 | 1236 FDOL |  | LDY \#2 |  | Y-COORDIMATE OF LABEL "A'. |
| 2580 | $12362613 C 11$ |  | ISR TUTOXY |  | SET TU.PTR TO POINT TO |
| 2540 |  | ; |  |  | SCREEN LOCATION "OF-LABEL "A |
| 26813 |  | ; |  |  |  |
| 2616 | 1238 A 000 |  | LDY \#\# |  | PUT LABELS ON SCREEN: |
| 2620 | 1230865112 |  | STY LBLCOL |  | INITIALIZE LABEL COLUATN |
| 2630 |  | ; |  |  | COUMTER. |
| 2 E 40 |  | ; |  |  |  |
| 2 ESG | $1240 \mathrm{ES5212}$ | LBLOOP | LDA LABELS, Y |  | GET A CHARACTER RND |
| 26Ei] | $1243257 C 11$ |  | JSR UUCHAR |  | PUT IT ON THE SCREEN. |
| 2570 | 1246 EES 112 |  | INC LBLCOL |  | PREPARE FOR NEXT CHARACTER. |
| 2600 | 1249 ACS 112 |  | LDY LBLCOL |  | DONE LAST CHARACTER? |
| 26ST | 1246 CEDA |  | CPY \#10 |  |  |
| 2780 | 124 E - 8 F |  | BNE LELOOP |  | IF NOT, DO NEXT CHARACTER. |
| 2710 |  | ; |  |  |  |
| 2720 | 125050 |  | RTS |  | RETURN TO CALLER. |
| 2730 | 125100 | LBLCOL | . BYTE D |  | DATA CELL: HOLDS COLUPTN |
| 2740 |  | ; |  |  | OF CHARACTER TO BE COPIED. |
| 2750 |  | ; |  |  |  |
| 27E日 |  | ; |  |  |  |
| 2770 |  | ; |  |  |  |
| 2780 |  | ; |  |  |  |
| 2750 | 125241 | LABELS | . EYTE A X | Y | $\mathrm{P}^{\prime}$ |
| 2730 | 125320 |  |  |  |  |
| 27:30 | 125420 |  |  |  |  |
| 2750 | 125559 |  |  |  |  |
| 2790 | 125620 |  |  |  |  |
| 2790 | 125720 |  |  |  |  |
| 2790 | 125959 |  |  |  |  |
| 2790 | 125920 |  |  |  |  |
| 2799 | 125420 |  |  |  |  |
| 2790 | 125B 50 |  |  |  |  |
| 2800 |  | ; |  |  |  |
| 2810 |  | ; |  |  |  |



| 3400 | 1283 A200 |  | LDX |  | START WITH ACCUMUCATOR |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 3410 |  | ; |  |  | IMRGE. |
| 3420 | 1285 Bn0112 | UUREGS | LDA | REGS, $X$ | LOOK UP THE REGISTER IMRGE. |
| 3430 | 128820 2311 |  | JSR | UUBYTE | DISPLAY IT IN HEX FORMAT. |
| 3440 | 128B 207F11 |  | JSR | TUSKIP | SKIP ONE SPACE AFTER HEX: |
| 3450 |  | ; |  |  | FIELD. |
| 3450 |  | ; |  |  | ! |
| 3470 | 128E E8 |  | INX |  | GET READY FOR NEXT REGISTER. $\because$ |
| 3480 | 120F EDO4 |  | CPX | \# 4 | DONE FOUR REGISTERS YET? |
| 3490 | 1291 DGF2 |  | BNE | UUREGS | IF NOT, DO NEXT ONE... |
| 3500 |  | ; |  | -. |  |
| 3510 | 129360 |  | RTS |  | If all Registers displayed, |
| 3520 |  | ; |  |  | RETURN. |
| 3530 |  | ; |  |  |  |
| 3540 |  | ; |  |  |  |
| 3550 |  | ; |  |  |  |
| 3550 |  | ; |  |  |  |
| 3570 |  | ; |  |  |  |
| 3580 |  | ; |  |  |  |
| 3550 |  | ; |  |  |  |
| 3600 |  | ; |  |  |  |
| 3610 |  | ; |  |  |  |
| 3620 |  | ; |  |  |  |
| 3630 |  | ; |  |  |  |
| 3640 |  | ; **** | ************ |  |  |
| 3650 |  | ; |  |  |  |
| 3660 |  | ; |  | GET SELEC | TEd byte |
| 3675 |  |  |  |  |  |
| 3680 |  | ; 赤** | ***** |  |  |
| 3690 |  | ; |  |  |  |
| 37 CD |  | ; |  |  |  |
| 3710 |  | ; |  |  |  |
| 3720 |  | ; |  |  |  |
| 3730 |  | ; |  |  |  |
| 3740 |  | , |  |  |  |
| 3750 | 1294 5502 | GET.SL | LDA | GETPTR | GET EYTE POINTED TO EY |
| 3760 | 129648 |  | FHA |  | THE SELECT POINTER |
| 3770 | 1297 fero3 |  | LDX | GETPTR+1 | (PRESERUING THE ZERO PAGE). |
| 3750 |  | ; |  |  |  |
| 3790 | 1299 ADG512 |  | LDA | SELECT |  |
| 3800 | 129C 6502 |  | STA | GETPTR |  |
| 3810 | 123E ADOSI2 |  | LDA | SELECT+1 |  |
| 3829 | 12 ค1 6503 |  | STA | GETPTR+1 |  |
| 38.30 |  | ; |  |  |  |
| 3840 | 12月3 F (20] |  | LDY |  |  |
| 3850 | 12 A 5 B102 |  | LIA | (GETPTR), Y |  |
| 3850 | 12 A 7 fB |  | TAY |  |  |
| 3875 | 12fig 68 |  | PLA |  |  |
| 38.80 | 12A9 8502 |  | STA | GETPTR |  |
| 3890 | $12 A B 8603$ |  | STX | GETPTR+1 |  |
| 3901 | 12AD 98 |  | TYA |  |  |
| 3910 | l2ate eb |  | RTS |  | RETURN TO CALLER. |
| 3920 |  | ; |  |  |  |
| 3930 |  | , |  |  |  |
| 3540 |  |  |  |  |  |
| 3950 |  | ; |  |  |  |
| 39 E 0 |  |  |  |  |  |
| 3970 |  | ; |  |  |  |


| 3980 | ； |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 3990 | ； |  |  |  |
| 4000 | ； |  |  |  |
| 4010 | ； |  |  |  |
| 4020 |  |  |  |  |
| 4030 |  | ； |  |  |
| 4040 |  | ；DISFLAY ARROW LINE |  |  |
| 4050 |  | ； |  |  |
| 4060 |  |  |  |  |
| 4070 |  | ； |  |  |
| 4080 |  | ； |  |  |
| 4090 |  | ； |  |  |
| 4100 |  | ； |  |  |
| 4110 |  | ； |  |  |
| 4120 | 12AF A202 | LINE． 3 | LDX \＃2 | LOAD $X$ WITH STARTING COLUMN． |
| 4130 | 12B1 A004 |  | LDY 4 | LOAD Y WITH ROW NUMBER： |
| 4140 | 12B3 203C11 |  | JSR TUTOXY | SET TU．PTR TO EEGINNING |
| 4150 |  | ； |  | OF ARROW LIME． |
| 4160 |  | ； |  |  |
| 4170 | 12BE ACBD12 |  | LDY FIELD | LOOK UP CURRENT FIELD． |
| 4180 | $12 \mathrm{B9} 38$ | SEC |  |  |
| 4190 | 12BA C007 | CPY \＃7 |  |  |
| 4200 | 12BC 9005 | BCC FLD．OK． |  |  |
| 4210 | I2BE AOQU | LDY \＃D |  |  |
| 4226 | 1200 ECDO12 | STY FIELD |  |  |
| 4230 | $12 \mathrm{C} 3 \mathrm{BSCD12}$ | FLD．OK | LDA FIELDS，$~$ | LOOK UP COLUFN NUMEER FOR |
| 4240 |  | ； |  | CURRENT FIELD． |
| 4250 |  | ； |  |  |
| 42E日 | $12 C 6$ A8 | TAY |  | USE THAT COLUMN NUMBER AS |
| 4270 |  | ； |  | AN INDEX INTO THE ROW． |
| 4280 |  | ； |  |  |
| 4290 | 12C7 ADQ710 |  | LDA ARROW | PLACE AN UP－ARROW IN |
| 4300 | 12CA 9160 | STA（TU．PTR），Y |  | COLUMN OF THE ARROW LINE． |
| 4319 | 12CC E0 |  | RTS | RETURN TO CALLER． |
| 4320 |  | ； |  |  |
| 4330 |  | ； |  |  |
| 4340 | 12CD 03 | FIELDS | ．EYTE 3，6，8 | THIS DATA PREA SHOWS WHICH |
| 4340 | 12CE 26 |  |  |  |
| 4340 | 12CF 08 |  |  |  |
| 4350 | 12D0 日B |  | ．EYTE \＄DE，\＄0E | COLUMN SHOLLD GET RN UP－ |
| 4350 | 12 D 1 DE |  |  |  |
| 4360 | 120211 |  | ．BYTE \＄11，\＄14 | RRROW TO INDICATE ANY ONE |
| 43ED | 12ロ3 14 |  |  |  |
| 4370 |  | ； |  | OF FIELDS O－E．CHANGING |
| 4380 |  | ； |  | ONE OF THESE UALUES WILL |
| 4350 |  | ； |  | CAUSE THE LIF－ARROW TO APFEAR |
| 4400 |  | ； |  | IN A DIFFERENT COLUMN WHEN |
| 4410 |  | ； |  | INDICATING A GIUEN FIELD． |
| 4420 |  | ； |  |  |
| 4430 |  | ； |  |  |

## Appendix C3:

Visible Monitor (Update Subroutine)


```
    APPENDIX CS: RSSEMBLER LISTING OF
                                    THE UISIBLE MONITOR
                                    UPDRTE SUSROUTINE
    SEE CHAPTER G OF BEYOND GAMES: SYSTEMS
SOFTWARE FOR YOUR G5OZ PERSONAL COMPUTER
            BY KEN SKIER
        EQUATES
    TU.PTP = 0
    GETPTR = 2
    PARAMS = $100G ADDRESS OF SYSTEM DATA
                        BLOCK.
    ARROW = PARAMS+7
                            THIS DATA BYTE HOLDS YOUR
                        SYSTEM'S CHARACTER CODE
                        FOR RIN UP-ARROW.
    ROMKEY = FARAMS+B
                                    ROMKEY IS A FOINTER TO
                                    YOUR S''STEM'S SUBROLITINE
                                    TO GET AN FSCII CHARACTER
                                    FROM THE KEYBOARD.
    DUITMY = PARAMST $10
                            DUMMY RETURNS WITHOUT DOING
    ANYTHING.
```






| 2000 |  | ； |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 2910 | 137 D A203 | ADRFLD | LDX \＃3 | SINCE ARROW IS UNDER ADDRESS |
| 2320 | 137F 18 | fidloce | CLC | FIELD，ROLL HEX DIGIT INTO |
| 2330 | 1300 beds 12 |  | ASL SELECT | ADDRESS FIELD BY ROLLING IT |
| 2940 | 1383250612 |  | ROL SELECT +1 | IT IIYTO THE POINTER THAT |
| 2550 | 1385 CA |  | DEX | SELECTS THE DISPLAYED |
| 2560 | 1357 10F6 |  | BPL ADLOOP | RDDRESS． |
| 2976 | 138958 |  | TYA |  |
| 2580 | 13 13 0 00512 |  | ORA SELECT |  |
| 2950 | 1300806512 |  | STA SELECT |  |
| 3000 | 139060 |  | RTS | THEN RETURN． |
| 3010 |  | ； |  |  |
| 3020 |  | ； |  |  |
| 3930 | 1391 E001 | NOTADR | CPX \＃1 | IS ARROW UNDER FIELD 1？ |
| 3040 | 1393 D818 |  | BNE REGFLD | IF NOT，IT MUST BE UNDER |
| 3050 |  | ； |  | A REGISTER FIELD． |
| 3000 |  | ； |  |  |
| 3070 | 1395 290\％ | ROL．SL | AND \＃\＄EF | ROLL 4 LSB IN A INTO |
| 3600 | 135748 |  | FHR | CURRENTLY－SELECTED BYTE． |
| 3099 | 1358209412 |  | JSR GET．SL | GET THE CURRENTLY－SELECTED |
| 3100 | 1358 RA |  | ASL A | BYTE AND SHIFT LEFT 4 TIMES．．． |
| 3110 | 135C DA |  | RSL A |  |
| 3120 | 1398 8A |  | RSL A |  |
| 3154 | 13SE［a |  | ASL A |  |
| 3144 | $139529 F 0$ |  | AND \＃\＃FFO |  |
| 3150 | 13 an 8pacis |  | STA TEMP |  |
| 3162 | 13月4 58 |  | PLA |  |
| 3176 | 1595 QDAC13 |  | ORA TEMP |  |
| 3109 | 13 AE 262013 |  | JSR PUT．SL | PUT IT IN CURREMTLY－SELECTED |
| 3150 | 13AE 60 |  | RTS | ADDRESS AND RETURN． |
| 3200 |  | ； |  |  |
| 3210 | 13AC 00 | TEMP | ．BYTE © |  |
| 32，${ }^{\text {co }}$ |  | ； |  |  |
| 3230 |  | ； |  |  |
| 3240 |  | ； |  |  |
| 3250 | 13ad CA | REGFLD | DEX | THE ARROL MUST EE UNDER A |
| 326s | 13PE CA |  | DEX | REGISTER IMRGE：FIELD 3. |
| 3278 | İAF CA |  | DEX | 4，5，OR E． |
| 3280 | 13808 Acou |  | LDY \＃3 |  |
| 3290 |  | ； |  |  |
| 3300 | 13 E 218 | RGLOOP | CLC | ROLL HEX DIGIT INTO |
| 3315 | 1383150112 |  | ASL REGS，$X$ | APPRROPRIATE REGISTER IMAGE． |
| 3こ2：0 | 13上： 39 |  | DEY |  |
| 3330 | 13571076 |  | BFL RGLOOP |  |
| 3345 | $13 \mathrm{ES} 1 \mathrm{nO112}$ |  | ORA REGS， X |  |
| 3350 | 139：300112 |  | STA REGS．$\times$ |  |
| 3360 | 13EF EQ |  | RTS |  |
| 3370 |  | ； |  |  |
| 3300 |  | ； |  |  |
| 3390 | 130068 | IF．CLR | FLA | RESTORE KEYBOARD CHARACTER． |
| 3400 | $1 \geq \mathrm{Cl}$ Cgrf |  | CMP \＃RUBOUT | IS IT RUBOUT？© IF YOUR |
| 3410 |  | ； |  | STSTEM DOESN＇T HAVE A |
| 34263 |  | ； |  | RUEOUT KEY，SUBSTITUTE THE |
| 3430 |  | ； |  | CODE FOR THE KEY YOU LL USE |
| 3446 |  | ； |  | TO CLEAR THE SCREEN． |
| 3456i |  | ； |  |  |
| 3460 | ： $30 \% 12604$ |  | EIE NOTCLR | If IT ISN＇$T$ THE＇CLEAR |
| 3470 |  | ； |  | GCREEM KEY，PERFORM NEXTT |


| 3480 |  | ; |  | TEST. |
| :---: | :---: | :---: | :---: | :---: |
| 3490 |  | ; |  |  |
| 3500 | 1305200011 |  | JSR CLR.TU | IF IT IS, THEN CLEAR THE |
| 3510 | 13 CE 50 |  | RTS | SCREEN AND RETURN. |
| 3520 | ; |  |  |  |
| 3530 |  | ; |  |  |
| 3545 | $1369 \mathrm{CS51}$ | MOTCLR | CMP ${ }^{\text {\# }}$ O | IS IT ' ${ }^{\prime}$ ' FOR QUIT? |
| 3550 | $13 C E$ DEO4 |  | ENE OTHER | IF NOT, FERFORM NEXT TEST. |
| 3560 |  | ; |  |  |
| 3570 |  | ; |  | IT IS ' $Q$ ' FOR QUIT. THE USER WANTS TO RETURN TO THE |
| 3500 |  | ; |  |  |
| 3590 |  | ; |  | USER WANTS TO RETURN TO THE CALLER OF THE UISIBLE |
| 3680 |  | ; |  | MONITOR. SO LET'S DO THAT: POP URDATE'S RETURM ADDRESS. |
| 3610 | 13CD 68 |  | PLA |  |
| 3520 | 13CE 58 |  | PLA |  |
| 3530 |  | ; |  |  |  |
| 3540 | 13CF 28 |  | PLP | RESTORE INITIAL E5日2 FLAGS. UISMON' 5 RETURN ADDRESS IS |
| 3650 |  | ; |  |  |
| 3550 |  | ; | RTS | NOW ON THE STACK. |
| 3670 | 130060 |  |  | SO RETURN TO CALLER OF |
| 3680 |  | ; ${ }^{\text {RTS }}$ |  | UISIMON. IN THIS WAY, |
| 3650 |  | ; |  | UISMON CAN EE USED BY ANY CALLER TO GET AN ADDRESS |
| 3780 |  | ; |  |  |
| 3710 |  | ; |  |  |
| 3720 |  | ; |  | FRGM THE USER. |
| 3730 |  | ; |  |  |
| 3740 | 1301201010 | OTHER | JSR DUMMY | REPLACE THIS CALL TO DUMMY WITH A CRLL TO RNY |
| 3750 |  |  |  |  |
| 3780 |  | ; |  | SUBROUTINE THAT EXTENDS |
| 3770 |  |  |  | FUNCTIONALITY OF THE |
| 3780 |  | ; | RTS | UISIELE MOMITOR. |
| 3790 | 13 D 450 |  |  | THEN RETURN. |
| 3862 |  |  | ; |  |
| 3810 |  | ; |  |  |
| 3820 |  | ; |  |  |
| 3830 |  | ; |  |  |
| 3540 |  | ; |  |  |
| 3850 |  | ; |  |  |
| 386] |  | ; |  |  |
| 3870 |  | ; |  |  |
| 3880 |  | ; |  |  |
| 3890 |  | ; |  |  |
| 350 |  | ; |  |  |
| 3910 |  | ; |  |  |
| 3920 |  |  |  |  |
| 3930 |  | , |  |  |
| 3540 |  | ASCII |  | BINRRY |
| 3950 |  | ; |  |  |
| 3560 |  |  |  |  |
| 3970 |  | ; |  |  |
| 3980 |  | ; |  |  |
| 3980 |  | ; |  |  |
| 4009 |  |  |  | IF RCCUMLUATOR HOLDS ASCII |
| 4010 |  | ; |  | g-9 OR A-F, THIS ROUTINE RETURN'S EIMARY EQUTUALENT-- |
| 4020 |  | ; |  |  |
| 4030 |  | ; |  | RETURHS EIHARY EQUIUALENT-OTHERNISE, IT RETIRNS SFF. |
| 4040 |  | ; |  |  |
| 4050 |  | ; |  |  |


| 4060 | 130538 | BINARY | SEC |  |
| :---: | :---: | :---: | :---: | :---: |
| 4070 | 13 DE ES30 |  | SBC | \＃ |
| 4080 | 1308 900F |  | ECC | BAD |
| 4090 | 13DA CSun |  | CMP |  |
| 4100 | 13DC 900e |  | BCC | G00D |
| 4110 | 130E Egor |  | SBC | \＃ |
| 4120 | 15 E 0 C910 |  | CMP | 咅等10 |
| 4130 | $13 \mathrm{E} 2 \mathrm{B0} 05$ |  | BCS | SAD |
| 4140 | $13 E 438$ |  | SEC |  |
| 4150 | $13 E 5$ CSOA |  | CMP | \％\＄0月 |
| 4160 | 13 E ？E003 |  | BCS | GOOD |
| 4170 | 13 ES ABff | EAD | LDA | \＃5FF |
| 4180 | 13EE 60 |  | RTS |  |
| 4150 |  | ； |  |  |
| 4200 | 13EC A200 | GOOD | LDX | \＃ |
| 4210 | 13EE S0 |  | RTS |  |

## Appendix C4:

## Print Utilities

| 15 | ; | APPENDIX C4: RSSEMBLER LISTING OF |
| :---: | :---: | :---: |
| 29 | ; | PRINT UTILITIES |
| 30 | ; |  |
| 4 IN | ; |  |
| 50 | ; |  |
| 60 | ; |  |
| 70 | ; | SEE CHAPTER 7 OF BEYOMD GAMES: SYSTEMS |
| 84 | ; | SOFTHARE FOR YOUR E502 PERSONAL COMPUTER |
| 90 | ; | -. |
| 106 | ; |  |
| 110 | ; | . |
| 120 | ; |  |
| 130 | ; |  |
| 148 | ; |  |
| 150 | ; |  |
| 160 | ; |  |
| 170 | ; |  |
| 180 | ; |  |
| 150 | ; |  |
| 206 | ; |  |
| 210 | ; |  |
| 220 | ; |  |
| 230 | ; | COMSTANTS |
| 240 | ; |  |
| 258 | ; |  |
| 2 EB | ; |  |
| 270 | ; |  |
| 230 | ; |  |
| 230 | ; |  |
| 330 | ; |  |
| 310 [000 = |  | $C R=$ \#SD $\quad$ CARRIAGE RETURM. |
| 320 | ; |  |
| $33 \mathrm{BLDFF}=$ |  | ETX = WFF $\quad$ THIS CHPRACTER MUST |
| 346 | ; | TERMINATE ANY MESSAGE STRING. |
| 350 | ; |  |
| 3603 ロ0¢ $=$ |  | LF = \$GA LINE FEED. |
| 370 | ; |  |
| 339 01000 |  | OFF $=\square$ |
| 399 | ; |  |
| 400 O日FF $=$ |  | $O N=\$ F F$ |
| 410 | ; |  |
| 420 | ; |  |
| 430 | ; |  |
| 446 | ; |  |
| 450 | ; |  |
| 4E0 | ; |  |
| 470 | ; |  |
| 480 | ; |  |
| 490 | ; |  |
| 500 | ; |  |
| 510 | ; |  |
| 520 | ; |  |
| 530 | ; | EXTERNAL ADDRESSES |
| 540 | ; |  |
| 559 | ; |  |
| 560 | ; |  |
| 570 | ; |  |
| 580 | ; |  |


| 590 | ; |  |
| :---: | :---: | :---: |
| 600 | ; |  |
| 610 | ; |  |
| 620 | ; |  |
| E30 | ; |  |
| 640 | ; |  |
| E50 | ; |  |
| $6501000=$ |  | PARAMS $=$ \$1DUG ADDRESS OF SYSTEM DATA BLOCK. |
| 670 | ; |  |
| 680 | ; |  |
| 690 | ; |  |
| $700100 \mathrm{C}=$ |  | ROMPRT = PARPM5+JUC |
| 710 | ; | POINTER TO ROM ROUTINE THAT |
| 720 | ; | SENDS CHAR TO SERIAL OUJTFUT. |
| 730 | ; |  |
| 740 | ; |  |
| 750 | ; |  |
| $760100 \mathrm{~A}=$ |  | ROMTUT = PARAMS + STDA |
| 770 | ; | POINTER TO ROM ROUTINE THAT |
| 780 | ; | PRINTS A CHPR TO THE SCREEN. |
| 750 | ; |  |
| 800 | ; |  |
| 810 | ; |  |
| 820 180E= |  | USROUT $=$ PARAMS+ \#DE |
| 830 | ; | POINTER TO USER-WRITTEN |
| 840 | ; | CHARACTER OUTPUT ROUTINE. |
| 850 | ; |  |
| 860 | ; |  |
| 870 | ; |  |
| 880 | ; |  |
| 850 1100= |  | TUSUBS $=\$ 1100$ |
| 900 1186= |  | ASCII $=$ TUSUES + \$86 |
| 910 | ; |  |
| 920 | ; |  |
| 930 | ; |  |
| 940 | ; |  |
| 550 1200= |  | UMPAGE $=$ \$1200 UISIbLE MONITOR STARTING |
| 960. | ; | PAGE |
| 970 | ; |  |
| 980 1205= |  | SELECT $=$ UMPAGE +5 |
| $5901254=$ |  | GET.SL $=$ UMPAGE +354 |
| $1000130 \mathrm{D}=$ |  | INC.SL = UMPAGE + F 100 |
| 1010 | ; |  |
| 1620 | ; |  |
| 1030 | ; |  |
| 1840 | ; |  |
| 1050 | ; |  |
| 1068 | ; |  |
| 1070 | ; |  |
| 1080 | ; |  |
| 1090 | ; |  |
| 1100 | ; | UARIGBLES |
| 1110 | ; |  |
| 1120 | ; |  |
| 1130 | ; |  |
| 1140 | ; |  |
| 1150 | ; |  |
| 1150 | ; |  |



| 1750 |  | ； |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 1760 |  | ； |  |  |
| 1770 |  | ； |  |  |
| 1780 |  | ； |  |  |
| 1790 |  | ； |  |  |
| 1800 | 1414 AgFF | FR．ON | LDA \＃ON | SELECT PRINTER FOR OUTPUT |
| 1810 | 1416 BDUQ14 |  | STA FRIMTR | by Setting its device flag． |
| 1620 | 141960 |  | RTS |  |
| 1830 |  | ； |  |  |
| 1840 |  | ； |  |  |
| 1250 |  | ； |  |  |
| 1360 |  | ； |  |  |
| 1870 |  | ； |  |  |
| 1880 | 141R Agロロ | PR．OFF | LDA \＃OFF | DE－SELECT PRINTER FOR OUTPUT |
| 1890 | 141 C 80014 |  | STA PRINTR | by Clearing ITS deuice flag． |
| 1900 | 141F EV |  | RTS |  |
| 1910 |  | ； |  |  |
| 1920 |  | ； |  |  |
| 1930 |  | ； |  |  |
| 1540 |  |  |  |  |
| 1950 |  | ； |  |  |
| 19E0 | 1420 ASFF | USR．ON | LDA EON | SELECT USER－HRITTEN |
| 1970 | 1422800214 |  | STA USER | SUBROUTINE EY SETTING |
| 1980 | 1425 60 |  | RTS | USER＇ 5 device flag． |
| 1930 |  | ； |  |  |
| 2000 |  | ； |  |  |
| 2010 |  | ； |  |  |
| 2020 |  | ； |  |  |
| 2030 |  | ； |  |  |
| 2040 | 1426 R900 | USROFF | LDA \＃OFF | DE－SELECT USER－WRITTEN |
| 2050 | 1423 880214 |  | STA USER | OUTPUT SUBRCUTINE EY |
| 2860 | 142B 60 |  | RTS | CLEARING ITS DEUICE FLAG． |
| 2070 |  | ； |  |  |
| 2080 |  | ； |  |  |
| 2090 |  | ； |  |  |
| 2100 |  | ； |  |  |
| 2110 |  | ； |  |  |
| 2120 | 142 C 200814 | ALL．ON | JSR TUT．ON | SELECT ALL OUTPUT DEUICES |
| 2130 | 142 F 201414 |  | JSR FR．ON | BY SELECTING EACH OUTPUT |
| 2140 | 1432202014 |  | JSR USR．ON | DEUICE INDIUIDUALLY． |
| 2150 | 143560 |  | RTS |  |
| 2150 |  | ； |  |  |
| 2170 |  | ； |  |  |
| 2180 |  | ； |  |  |
| 2190 |  | ； |  |  |
| 2200 |  | ； |  |  |
| 2210 | 14.3620014 | ALLOFF | JSR TUTOFF | De－select all output devices |
| 2220 | 1439 201月14 |  | JSR PR．OFF | BY DE－SELECTING ERCH OME |
| 2230 | $143 C 202514$ |  | JSR USROFF | INDIVIDUALLY． |
| 2240 | 143 F E0 |  | RTS |  |
| 2250 |  | ； |  |  |
| 2250 |  | ； |  |  |
| 2270 |  | ； |  |  |
| 2280 |  | ； |  |  |
| 2290 |  | ； |  |  |
| 2300 |  | ； |  |  |
| 2310 |  | ； |  |  |
| 2320 |  | ； |  |  |







| 52301 | 151160 |  | RTS | RETURN (TO BYTE IMMEDIATELY |
| :---: | :---: | :---: | :---: | :---: |
| 5240 |  | ; |  | FOLLOWING THE ETX.3 |
| 5250 |  | ; |  |  |
| 5260 |  | ; |  |  |
| 5270 |  | ; |  |  |
| 5280 |  | ; |  |  |
| 5290 |  | ; |  |  |
| 5300 |  | ; |  |  |
| 5310 |  | ; |  |  |
| 5320 |  | ; |  |  |
| 5330 |  | ; |  |  |
| 5340 |  | ; |  |  |
| 5350 |  | ; ***** |  |  |
| 5360 |  | ; |  |  |
| 5370 |  |  | SAUE, RESTOR | S SELECT POINTER |
| 5380 |  | ; |  |  |
| 5330 |  | ; ****** |  |  |
| 5400 |  | ; |  |  |
| 5410 |  | ; |  |  |
| 5420 |  | ; |  |  |
| 5430 |  | ; |  |  |
| 5440 |  | ; |  |  |
| 5450 | 151268 | PUSHSL | PLA | PULL RETURN ADDRESS FROM |
| 5460 | 1513 800614 |  | STA RETURN | Stack amd save it in return. |
| 5470 | 151668 |  | FLA |  |
| 5480 | 1517 200714 |  | STA RETURN+1 |  |
| 5430 |  | ; |  |  |
| 5500 |  | ; |  |  |
| 5510 | 151A ADEE12 |  | LDA SELECT+1 | FUSH SELECT POINTER ONTO |
| 5520 | 151048 |  | PHA | THE STACK. |
| 5530 | 151E AD0512 |  | LDA SELECT |  |
| 5540 | 152148 |  | PHA |  |
| 5550 |  | ; |  |  |
| 5560 |  | ; |  |  |
| 5570 | 1522 AD0714 |  | LDA RETURN+1 | FUSH RETURN RDDRESS ERCK |
| 5580 | 152548 |  | FHA | ON THE STACK. |
| 5550 | 1526 ADE614 |  | LDA RETURN |  |
| 5600 | 152948 |  | PHA |  |
| 5610 |  | ; |  |  |
| 5620 |  | ; |  |  |
| 5630 | 152A E0 |  | RTS | RETURN TO CAELER. CALLER |
| 5640 |  | ; |  | WILL FIND SELECT ON STACK. |
| 5650 |  | ; |  |  |
| 5660 |  | ; |  |  |
| 5670 |  | ; |  |  |
| 5650 |  | ; |  |  |
| 5690 |  | ; |  |  |
| 5700 |  | ; |  |  |
| 5716 |  | ; |  |  |
| 5720 |  | ; |  |  |
| 5730 | 152B 68 | POP. SL | PLA | SAUE RETURY GLDRESS. |
| 5740 | 152C 800614 |  | STA RETURN |  |
| 5750 | 152F 68 |  | PLA |  |
| 5750 | 1530800714 |  | STA RETURN+1 |  |
| 5779 |  | ; |  |  |
| 5780 |  | ; |  |  |
| 5790 | 153368 |  | PLA | LOAD SELECT FROM STACK |
| 5800 | 1534800512 |  | STA SELECT |  |



## Appendix C5: <br> Two Hexdump Tools

| 10 | ; | AFPENDIX C5: | ASSEMBLER LISTING OF |
| :---: | :---: | :---: | :---: |
| 20 | ; | THO | HEXDUMP TOOLS |
| 30 | ; |  | . |
| 40 | ; |  |  |
| 50 | ; |  |  |
| 50 | ; | SEE CHAPTER | 8 OF BEYOMD GAMES: SYSTEMS |
| 70 | ; | SOFTWRRE FOR YOUR | 6502 PERSOMAL COMPUTER |
| E] | ; |  |  |
| 90 | ; |  |  |
| 100 | ; | - | BY KEN SKIER |
| 110 | ; |  |  |
| 120 | ; |  |  |
| 130 | ; |  |  |
| 140 | ; |  |  |
| 150 | ; |  |  |
| 160 | ; |  |  |
| 170 | ; |  |  |
| 180 | ; |  |  |
| 190 | ; |  |  |
| 200 | ; |  |  |
| 210 | ; |  |  |
| 220 | ; |  |  |
| 230 | ; |  |  |
| 240 | ; |  |  |
| 250 | ; |  |  |
| 260 | ; |  |  |
| 275 | ; |  |  |
| 280 | ; | CONSTANTS |  |
| 239 | ; |  |  |
| 300 | ; |  |  |
| 310 | ; |  |  |
| 320] | ; |  |  |
| 300 | ; |  |  |
| 340 | ; |  | , |
| 350 | ; |  |  |
| $3605000=$ |  | $C R=\$ 3 D$ | CARRIfGE RETLJRN. |
| 370 | ; |  |  |
| 380 DeDa= |  | $L F=$ DRA | LINE FEED. |
| 390 | ; |  |  |
| 400 | ; |  |  |
| $410 \mathrm{n} 07 \mathrm{~F}=$ |  | TEX $=$ \# 7 F | THIS CHARACTER MUST START |
| 420 | ; |  | ANY MESSAGE. |
| 430 | ; |  |  |
| 440 O3FF $=$ |  | $E T X=\$ F F$ | THIS CHARACTER MUST END |
| 45] | ; |  | ANY MESSAGE. |
| 460 | ; |  |  |
| 470 | ; |  |  |
| 450 | ; |  |  |
| 490 | ; |  |  |
| 500 | ; |  |  |
| 510 | ; |  |  |
| 520 | ; |  |  |
| 530 | ; |  |  |
| 540 | ; |  |  |
| 550 | ; |  |  |
| 550 | : |  |  |
| 570 | ; |  |  |


| 500 | ; |  |
| :---: | :---: | :---: |
| 590 | ; |  |
| 000 | ; |  |
| 610 | ; | EKTERNAL ADDRESSES |
| E20 | ; |  |
| 630 | ; |  |
| 640 | ; |  |
| 659 | ; |  |
| E60 | ; |  |
| 670 | ; |  |
| 656 | ; |  |
| E90 | ; |  |
| 700 | ; |  |
| 710 | ; |  |
| 720 | ; |  |
| 730 | ; |  |
| $7401100=$ |  | TVSUES $=$ \$110日 STARTING PAGE OF DISPLAY |
| 750 | ; | code. |
| 760 1100= |  | CLR.TU=TUSUSS |
| $7701186=$ |  | FSCII =TUSUES+\$BS |
| 760 | ; |  |
| 750 | ; |  |
| 600 1200= |  | MTPAGE= 1200 STARTING FAGE OF UISIBLE |
| 816 | ; | MOMITOR CODE. |
| $8201205=$ |  | SELECT=UMPAGE +5 |
| $8301207=$ |  | UISMON=UMPAGE +7 |
| $8401294=$ |  | GET.SL=UMPAGE+ \% $^{\text {S }} 4$ |
| $6591309=$ |  | INC. SL=UMPAGE + \$10D |
| 869 | ; |  |
| 870 | ; |  |
| $6091409=$ |  | PRFRGE $=\$ 1400$ STARTING PAGE OF PRINT |
| 830 | ; | UTILITIES. |
| 906 $1408=$ |  | TUT. ON=PRPAGE + S |
| 910 140E= |  | TUTOFF=PRPAGE + \#VE |
| 920 141.4= |  | PR. ON =PRPAGE+ +14 |
| 330 1410= |  | PR.OFF=FRPRGE+ + IR |
| $9401440=$ |  | PR. CHR=PRFAGE + \$ 40 |
| $9501472=$ |  | CR.LF =PRPAGE + \$72 |
| S56 147D= |  | SPACE $=$ FRPAGE + \$7D |
| 5781 149 $=$ |  | SPRCES $=$ PRPAGE + \$96 |
| SE5 1483= |  | PR.BYT=PRPRAGE + \$83 |
| SSO 14E4= |  | PRINT: =FRPAGE+\$E4 |
| $10001512=$ |  | PUSHSL=FRPFGE + \$112 |
| 1619 152B= |  | POP.SL=PRPAGE + \#12B |
| 1020 | ; |  |
| 1030 | ; |  |
| 1046 | ; |  |
| 1050 | ; |  |
| 10en | ; |  |
| 1079 | ; |  |
| 1080 | ; |  |
| 1090 | ; |  |
| 1180 | ; |  |
| 1110 | ; |  |
| 1120 | ; |  |
| 1130 | ; |  |
| 1140 | ; |  |
| 1150 | ; | UARIABLES |









```
415% 15EF 53
4150 15SQ 53
4150 1591 20
4150 1692 40
4150 15S3 45
4150 1594 53
4150 1635 53
4150 1EG5 20
4150 1ES7 54
4150 1598 45
4150 1ESS 41
4150 16G茵 45
4150 1ESE 20
4150 16SC 53
4150 1690 54
4 1 5 0 ~ 1 6 5 E ~ 4 1 ~
4150 1GGF 52
4150 1ERO 54
4150 1EN1 20
4150 16&2 41
415% 15A3 44
415G IEA4 44
4150 1EFS 52
415卬 LBHE 45
4150 15A7 53
4150 1EnO 53
4150 15Ag 2C
41EO 1Б^R 20
41GO 1EAB 5%
41ES IEAC 4B
41EO IEAD 43
41ER LERE 43
41EO LEFF 4B
41EO 1GEO 20
41ED IEEL 49
41EO 16E2 53
4160 18E3 20
41EQ 1EEA FF
41TO 15ES 20SE1E
4180 ;
4190 IEBO 4CICIS
4210 ;
4こご品 ;
4ごころ ;
4240] ;
4250l ;
42E日 ;
4270
4280 ;
4290 ;
4300 ;
4312 ;
4320 ;
4330 ;
4340 ;
4350 ;
43E0 ;
```

JSR PR．SA PRINT START ADDRESS．

JMP SET．EA RND LET THE USER SET A NEW END ADDRESS．



```
52501717 20
52501718 33
5250 1719 20
5250 171月 20
5250 1718 34
5250 171C 20
5250 171D 20
5250 171E 35
5250 171F 20
5250 1720 20
5250 1721 35
5250 1722 20
5250 1723 20
5250 1724 37
52501725 20
5250 1725 20
5250 1727 38
5260 1728 20
5250 1729 20
5260 172A 39
5260 1728 20
5260 1720 20
5260 172D 41
52E0 172E 20
52E0 172F 20
5250 173042
5260 1731 20
5250 1732 20
52601733 43
5250 1734 20
5260}17352
S2E6 173644
5260 1737 20
52501738 20
5260 173945
5260 173^ 20
5260 1738 20
5250 173C 46
5270 173D 0D
5270 173E DA
5270 173F GA
5270 1740 FF
5250 1741 60
5250 ;
5300 ;
5310 ;
5320 ;
5330 ;
5340 ;
5350 ;
53E0 ;
5370 ;
5350 ;
5330
54,0
5 4 1 0
5420
5430
```

| 5440 |  | ； |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 5450 |  | ； |  |  |
| 5460 |  | ； |  |  |
| 5470 |  | ； |  |  |
| 5480 |  | ； |  |  |
| 5490 | 174220.214 | PRLINE | JSR CR．LF |  |
| 5500 | 1745 AD0512 |  | LDA SELECT | DETERMINE STARTING COLUMN． |
| 5510 | 174848 |  | FHA | FOR THIS DUMP． |
| 5520 | 1749 298F |  | AND \＃\＄5FF |  |
| 5530 | 174 B 85515 |  | STA COLUMN | NOW COLUMN HOLDS NUMEER OF |
| 5540 |  | ； |  | HEX COLUMN IN WHICH WE DUMP |
| 5550 |  | ； |  | THE FIRST BYTE． |
| 5560 | 174E 68 |  | PLA | SET SELECT＝EEGINNING OF $A$ |
| 5570 | 174F 29Fb |  | AND \＃\＄FD | HEX LINE． |
| 5580 | 1751800512 |  | STA SELECT |  |
| 5590 | 175420 ค115 |  | JSR PR．ADR | PRINT LINE＇S START ADDRESS． |
| 5600 | 1757 A203 |  | LDX 昔3 | SPACE 3 TIMES－－TO THE |
| 5510 | 1759209614 |  | JSR SPACES | FIRST HEX COLUMN． |
| 5620 |  | ； |  |  |
| 5630 |  | ； |  |  |
| 5640 | 175C AD5615 |  | LDA COLUMA | DO WE DUPMP FROM THE FIRST |
| 5650 |  | ； |  | HEX COLUMN？ |
| 5EED | 175F FOOD |  | BEQ COL．OK | If SO，WERE AT THE CORRECT |
| 5670 |  | ； |  | COLLIN NOW． |
| 5580 |  | ； |  |  |
| 5690 | 1751 A203 | LOOP | LDK 共3 | IF NOT，SPACE 3 TIMES FOR |
| 5700 | 1763209614 |  | JSR SPACES | ERCH BYTE NOT DUPIPED． |
| 5710 | 1766200013 |  | JSR INC．SL |  |
| 5720 | 1769 CE5615 |  | DEC COLUMN |  |
| 5730 | 176C DOF3 |  | ENE LOOP |  |
| 5740 |  | ； |  |  |
| 5750 | 176E 209al5 | COL．OK | JSR DUMPSL | DUMP SELECTED BYTE． |
| 5760 | 1771207014 |  | JSR SPACE | SPACE ONCE． |
| 5770 | 1774208317 |  | JSR NEXTSL | SELECT NEXT BYTE |
| 5750 |  | ； |  |  |
| 5790 | 17773009 |  | BMI EXIT | MINLE MEANS WE＇UE DUMPED |
| 5850 |  | ； |  | THROUGH TO THE END ADDRESS． |
| 5810 |  | ； |  |  |
| 5820 |  | ； |  |  |
| 5830 | 1779 An0512 | NOT．EA | LDA SELECT | DUMPED EIfTIRE LIME？ |
| 5840 | 177C 290F |  |  | （4LSB OF SELECT＝0？） |
| 5850 | 177E CSOU |  | CMP \＃${ }^{\text {B }}$ | IF SO，WE＇UE DUMPED THE |
| 5850 |  | ； |  | ENTIRE LINE．IF NOT， |
| 5670 | 1780 日GEC |  | BNE COL．OK | SELECT NEXT BYTE AND DUMP IT． |
| 5880 | 1782 60 | EXIT | RTS | RETURN MINUS If EA DUMPED； |
| 5890 |  | ！ |  | RETURN PLUS IF EA NOT DUMPED． |
| 5900 |  | ； |  |  |
| 5910 |  | ； |  |  |
| 5920 |  | ； |  |  |
| 5530 |  | ； |  |  |
| 5340 |  | ； |  |  |
| 5950 |  | ； |  |  |
| 5560 |  | ， |  |  |
| 5976 |  | ； |  |  |
| 5980 |  |  |  |  |
| 5990 |  | ； |  |  |
| 6000 |  | ；＊＊＊ | ＊＊＊＊＊＊＊＊＊＊＊＊＊＊ |  |
| E010 |  | ； |  |  |


| 5020 |  |  | SELECT NEXT BYTE（IF＜END ADDRESS） |  |
| :---: | :---: | :---: | :---: | :---: |
| 6030 |  | ； |  |  |
| 6040 |  |  |  |  |
| 6850 |  | ； |  |  |
| 6060 |  | ； |  |  |
| 6070 |  | ； |  |  |
| 60su |  | ； |  | ！ |
| 5090 |  | ； |  |  |
| 6100 | 1783 38 | NEXTSL | SEC |  |
| 5110 | 1784 AD0612 | LDR SELECT +1 |  | MIGH BYTE OF SELECT LESS |
| 6120 | 1787 CD5515 |  | CMP EA＋1 | THAN HIGH BYTE OF EA？ |
| 6130 | 17EA 90DB |  | BCC SL．OK | IF SO，SELECT＜END ADDRESS． |
| 6140 | 178C DOAF |  | BNE NO．INC | If SELECT＞EA，DON ${ }^{\text {T }}$ |
| 5150 |  | ； |  | INCREMENT SELECT． |
| 6160 |  | ； |  |  |
| 5170 | 178E 38 |  | SEC | SELECT IS IN SRME PAGE RS EA． |
| 6180 | 178F ALES12 |  | LDA SELECT |  |
| 6198 | 1732 CD5415 |  | CMP ER |  |
| 6200 | 1755 B00s |  | BCS MO．InC |  |
| 6210 |  | ； |  |  |
| 6220 | 1797200513 | SL．OK | JSR INC．SL | SInce SElect $<=$ Ef，WE MAY |
| E230 |  | ； |  | INCREMENT SELECT． |
| 6240 |  | ； |  |  |
| 6250 | 179A ASED |  | LDA 故 | SET＂INCREMENTED＂RETURN |
| 5260 | 179C E日 |  | RTS | CODE AND RETURN． |
| 6270 |  | ； |  |  |
| 5280 | 179D RSFF | NO．INC | LDA 茾娩FF | SET＂NO INCREMENT＂RETURN |
| 6250 | 179F E0 |  | RTS | CODE AND RETURN． |
| 6300 |  | ； |  |  |
| 6310 |  | ； |  |  |
| 6320 |  | ； |  |  |
| 5330 |  | ； |  |  |
| 6340 |  | ； |  |  |
| 6359 |  | ； |  |  |
| $536 \square$ |  | ；＊＊＊＊＊ |  | ＊＊ |
| 6370 |  | ； |  |  |
| 6380 |  | ； | SELECT S | ART PIDDRESS |
| E390 |  | ； |  |  |
| 6460 |  | ；＊＊＊＊ | ＊＊＊＊＊＊＊＊＊＊＊＊＊＊ | ＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊ |
| 6410 |  | ； |  |  |
| 6420 |  | ； |  |  |
| 6430 |  | ； |  |  |
| 6440 |  | ； |  |  |
| 5450 |  | ； |  |  |
| 6460 | 17AD RDS215 | gotosa | LDf 5A | SET SELECT＝SA． |
| E476 | 17A3 8DO512 |  | STR SELECT |  |
| 5480 | 17AS ADS315 |  | LDA SA＋1 |  |
| 6490 | 17 AS 8DE612 |  | STA SELECT＋1 |  |
| 6500 | 17RC E6 |  | RTS | RETURN W？SELECT＝5A． |

## Appendix C6:

## Table-Driven Disassembler (Top

Level and Utility Subroutines)








| 3620 |  | ； |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 3630 |  |  |  |  |  | GET LOW BYTE OF Xth POINTER |
| 3640 | 1988 BD1b1B | MODE．$\times$ |  | SUBS，$X$ |  |  |
| 3650 | 19BE 8D0419 |  | STA | SUBPTR |  | IN TABLE OF SUBROUTIINE |
| 3650 |  | ； |  |  |  | POINTERS． |
| 3670 | 19BE ES |  | INX |  |  | ADJUST INDEX FOR NEXT BYTE． GET HIGH BYTE OF POINTER． |
| 3EEO | 19EF BDIBIB |  | LDA | SUBS，$X$ |  |  |
| 3690 | 1902800519 |  | STA | SUPPTR＋1 |  |  |
| 3700 | $19 C 5$ ECD419 |  | JMP | （SUSPTR） |  | JUMF TO SUBROUTINE SPECIFIED |
| 3710 |  | ； |  |  |  | BY SUBROUTINE POINTER． |
| 3720 |  | ； |  |  |  | THAT SUBROUTINE WILL RETURN |
| 3730 |  | ； |  |  |  | TO THE CALLER OF MODE．$X$ ， |
| 3740 |  | ； |  |  |  | NOT TO MODE．X ITSELF． |
| 3750 |  | ； |  |  |  |  |
| 3750 |  | ； |  |  |  |  |
| 3770 |  | ； |  |  |  |  |
| 3780 |  | ； |  |  |  |  |
| 3790 |  | ； |  |  |  |  |
| 3506 |  | ； |  |  |  |  |
| 3810 |  | ； |  |  |  |  |
| 3820 |  | ； |  |  |  |  |
| 3830 |  | ； |  |  |  |  |
| 3840 |  | ； |  |  |  |  |
| 3850 |  | ；＊＊＊＊＊＊ | ＊${ }^{\text {a }}$ 米为 |  | ＊＊＊＊＊＊ |  |
| 3860 |  | ； |  |  |  |  |
| 3870 |  |  |  | DISASSE | EMBL | LER UTILITIES |
| 3850 |  | ； |  |  |  |  |
| 3890 |  | ；＊＊＊＊＊＊ | 为为数常 |  | 粅为为为米 |  |
| 3900 |  | ； |  |  |  |  |
| 3916 |  | ； |  |  |  |  |
| 3920 |  | ； |  |  |  |  |
| 3930 |  | ； |  |  |  |  |
| 3940 |  | ； |  |  |  |  |
| 3950 |  | ； |  | PRINT O | ONE－ | －BYTE OPERAND |
| 3960 |  | ； |  |  |  |  |
| 3970 |  | ； |  |  |  |  |
| 3980 |  | ； |  |  |  |  |
| 3959 | 1scs 200013 | ONEBYT | JSR | INC．SL |  | ADUANCE TO BYTE FOLLOWING |
| 4000 |  | ； |  |  |  | OPCODE． |
| 4010 | 190E 209A15 |  | JSR | DUIPSL |  | DLITP THAT BYTE． |
| 4020 | 19CE E0 |  | RTS |  |  | RETUFN TO CRLLER． |
| 4030 |  | ； |  |  |  |  |
| 4040 |  | ； |  |  |  |  |
| 4050 |  | ； |  |  |  |  |
| 4000 |  | ， |  |  |  |  |
| 4070 |  | － |  |  |  |  |
| 4080 |  | ； |  | PRINT T | TWO－ | －BYTE OPERAND： |
| 40.90 |  | ； |  |  |  |  |
| 4100 |  | ； |  |  |  |  |
| 4110 |  | ； |  |  |  |  |
| 4120 | 190f 260013 | TWOEYT | JSR | INC．SL |  | ADUANCE TO FIRST BYTE OF |
| 4130 |  | ； |  |  |  | OFERAND． |
| 4140 | 1902203412 |  | JSR | GET．SL |  | LOAD THAT BYTE INTO ACC． |
| 4150 | 190548 |  | PHf |  |  | SAVE IT． |
| 4150 | 1905 20013 |  | JSR | INC．SL |  | ADUANCE TO 2ND BYTE OF |
| 4170 |  | ； |  |  |  | OFERAND． |
| 4180 | 19D9 209A15 |  | JSR | DUMPSL |  | DUMP IT． |
| 4150 | 19DC E3 |  | PLA |  |  | RESTORE FIRST BYTE TO ACC． |




## Appendix C7:

Table-Driven Disassembler (Addressing Mode Subroutines)


| 580 | ; |  |
| :---: | :---: | :---: |
| 590 | ; |  |
| 600 | ; |  |
| 510 | ; |  |
| E20 | ; |  |
| E30 | ; |  |
| 646 | ; |  |
| cera | ; |  |
| E日0 | ; |  |
| 670 | ; | EXTERNAL ADDRESSES |
| EEO | ; |  |
| ESa | ; |  |
| 700 | ; |  |
| 710 | ; |  |
| 720 | ; |  |
| 730 | ; |  |
| 740 | ; |  |
| 750 | ; |  |
| 760 | ; |  |
| 779 | ; |  |
| 780 | ; |  |
| 790 | ; |  |
| 806 | ; |  |
| 810 1200= |  | UMPAGE $=$ \#1200 STARTING PAGE OF UISIBLE |
| 820 | ; | MONITOR CODE. |
| 830 1205= |  | SELECT=UMPAGE +5 |
| $8401297=$ |  | UISMON=UMPAGE+7 |
| $8591294=$ |  | GET. SL=UMPRGE+\$34 |
| 600 130U= |  | INC. SL=UMPRGE + \$10D |
| $8701514=$ |  | BEC.SL=UMPAGE+\$11A |
| 8ed | ; |  |
| 890 | ; |  |
| $9001409=$ |  | PRPFGEE $=1400$ STARTING PAGE OF PRINT |
| 910 | ; | UTILITIES. |
| $9201440=$ |  | PR. CHR=PRPAGE + W, 40 |
| 930 1472= |  | CR.LF =PRPAGE+ +72 |
| S40 147D= |  | SPACE =FRPAG:+ + 7 |
| 9561 1496= |  | SPACES=PRPAGE + \$96 |
| $5601483=$ |  | PR. BYT=PRPAGE+\$83 |
| $57014 E 4=$ |  | PRINT: =PRPRGE+ + EL 4 |
| $9801512=$ |  | FUSHSL=PRPAGE + \$112 |
| 950 1528= |  | POP. SL=PRPRGE +12 B |
| 15100 | ; |  |
| 1010 | ; |  |
| $10201500=$ |  | HEX.PG=T1500 RDDRESS OF PAGE IN WHICH |
| 1030 | ; | HEXDUMP CODE STARTS. |
| 1040 | ; |  |
| $105015 \mathrm{Al}=$ |  | PR. $A D R=H E X . P G+\mathbb{\#} 1$ |
| 10E0 1783= |  | NEXTSL=HEX.PG + \$283 |
| 10 O | ; |  |
| 1080 | ; |  |
| 1090 1909 $=$ |  | DSPAGE $=\$ 1900$ START OF DISASSEMBLER CODE. |
| 1100 | ; |  |
| $11101908=$ |  | ONEBYT= DSFAGE + \#C8 |
| $112019 \mathrm{CF}=$ |  | TWOBYT $=$ DSPAGE + \#CF |
| $113019 E 1=$ |  | LPARET $=$ DSFAGE+\#E1 |
| $11451955=$ |  | RFAREN $=$ DSPAGEE + EE 5 |
| $115019 E B=$ |  | XINDEX $=$ DSPAGE + SEB |




| 2320 |  | ； |  | INDIRECT | mode |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 2330 |  | ； |  |  |  |
| 2340 |  | ； |  |  |  |
| 2350 |  | ； |  |  |  |
| 2360 | 1ATF 20E115 | INDRCT | JSR L | LFAREN | PRINT LEFT PARENTHESIS． |
| 2370 | 1AB2 204Eila |  | JSR A | ABSLUT | PRINT TWO－EYTE OPERAND． |
| 2360 | 1A8E 20E519 |  | JSR R | RPAREN | PRINT RIGHT PARENTHESIS． |
| 2390 | 1R8E A306 |  | LDA | 和6 | A HOLDS NUMBER GF CHARACTERS |
| 2400 |  | ； |  |  | －IN OPERAND． |
| 2410 | 1aba azgz |  | LDX \＃ | \＃2 | $\times$ HOLDS NUMPER OF BYTES IN |
| 2420 |  | ； |  |  | OPERAIMD． |
| 2430 | 1HSC 60 |  | RTS |  | RETURN TO CALLER． |
| 2440 |  |  |  |  |  |
| 2450 |  | ； |  |  |  |
| 2465 |  | ； |  |  |  |
| 2470 |  | ； |  |  |  |
| 2480 |  | ； |  |  |  |
| 24900 |  | ； |  | INDIRECT | ，$\times$ MODE |
| 2500 |  | ； |  |  |  |
| 2516 |  | ； |  |  |  |
| 2520 |  | ； |  |  |  |
| 2530 | 1ABD 20E119 | IND． X | J5R | LPAREN |  |
| 2540 | 1月90 20Esia |  | JSR | ZERO．X | PRINT A ZERO FAGE ADDRESS， |
| 2555 |  | ； |  |  | A COMMA，AND THE LETTER＂X＂． |
| 2550 | 1RS3 20E519 |  | JSE | RPRREN |  |
| 2570 | 1RSS AZDI |  | LDX | 41 | ONE BYTE IN OPERAND． |
| 2580 | 1 A98 R9GE |  | LDA | 非显 | 8 CHARACTERS IN OPERAND． |
| 2590 |  | ； |  |  | （C－IP OWIERS：AG BE，NOT |
| 2600 |  | ； |  |  | AS 08，FOR NARROU FORMAT． 3 |
| 2810 | 1R9A 50 |  | RTS |  |  |
| 2620 |  | ； |  |  |  |
| 2630 |  | ； |  |  |  |
| 2640 |  | ； |  |  |  |
| 2650 |  | ； |  |  |  |
| 2650 | ． | ； |  |  |  |
| 2670 |  | ； |  | MDIRECT，Y | MODE |
| 2680 |  | ； |  |  |  |
| 2690 |  | ； |  |  |  |
| 2700 |  | ； |  |  |  |
| 2710 | 1f9SE 20.119 | IND．Y | JSR | LFAREN |  |
| 2720 | 1ABE 20DE1A |  | J SR | ZEROPG | FRINT A ZERO PAGE ADDRESS． |
| 2730 | 1RA1 $20 E 519$ |  | JSR | RFAREN |  |
| 2740 | 1AR4 20F619 |  | JSR | YINDEX | PRINT A COMMA AND A＂Y＂． |
| 2750 | 1RAT A2D1 |  | LDK | \＃1 | OPERAND HAS 1 BYTE．．． |
| 2760 | lans ages |  | LDA | 薪 | ．．．AND E CHARACTERS． |
| 2770 |  | ； |  |  | （C－IP OWNERS：R9 EE，NOT |
| 2780 |  | － |  |  | AS［BE．FOR NARROW FORMAT． |
| 2790 | 1ARE ED |  | RTS |  |  |
| 2800 |  | ； |  |  |  |
| 2810 |  | ； |  |  |  |
| 2820 |  | ； |  |  |  |
| 2830 |  | ； |  |  |  |
| 2840 |  | ； |  |  |  |
| 2850 |  |  |  | RELATIUE | MODE |
| 2860 |  | ； |  |  |  |
| 2876 |  | ； |  |  |  |
| 2880 |  | ； |  |  |  |
| 2890 | 1ARC 200 D 13 | RELATU | JSR | INC．SL | SELECT NEXT BYTE． |


| 2900 | 1ARF | 201215 |  | JSR | PUSHSL | SAUE SELECT FOINTER ON STACK. |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2910 | 1AB2 | 293412 |  | JSR | GET.SL | GET OPERAND BYTE. |
| 2920 | 1AB5 | 48 |  | PHA |  | SAVE IT ON STACK. |
| 2930 | 1ABE | 200013 |  | JSR | IMC. SL | INCREMENT SELECT POINTER |
| 2940 |  |  | ; |  |  | SO IT FOINTS TO NEXT OPCODE. |
| 2950 |  |  | ; |  |  | (RELATIVE ERANCHES RRE |
| 2550 |  |  | ; |  |  | RELATIUE TO NEXT OFCODE. 3 |
| 2970 | $19 B 9$ | 68 |  | PLA |  | RESTORE OPERAND BYTE TO ACC. |
| 2380 | IABA | C900 |  | CMP | \# | IS IT PLUS OR MINUS? |
| 2950 | IABC | 1003 |  | BPL | FORWRD | IF PLUS, IT MEANS A FORWARD |
| 3000 |  |  | ; |  |  | BRANCH. |
| 3010 |  |  | ; |  |  |  |
| 3020 |  |  | ; |  |  | OPERAND IS MINUS, 50 WE' LL |
| 3030 |  |  | ; |  |  | BRANCH BACKUARD. |
| 3040 | IABE | CE0612 |  | DEC | SELECT +1 | BRANCHING BACKWRRD IS LIKE |
| 3050 |  |  | ; |  |  | BRANCHING FORWARED FROM ONE |
| 3060 |  |  | ; |  |  | PAGE LOWER IN MEMORY. |
| 3070 |  |  | ; |  |  |  |
| 3080 |  |  | , |  |  |  |
| 3090 | 1AC1 | 08 | FORWRD | PHP |  | SRUE CALLER' 5 decimal flrg. |
| 3100 | 1ACZ | D8 |  | CLD |  | CLEAR DECIMPL MODE, FOR |
| 3110 |  |  | ; |  |  | BINARY ADDITIION. |
| 3120 | IfC3 | 18 |  | CLC |  | FREPARE TO ADD. |
| 3130 | 1AC4 | 600512 |  | ADC | SELECT | ADD OPERAND BYTE TO SELECT. |
| 3140 | 1AC7 | 9003 |  | BCC | RELEMD |  |
| 3150 | IRCS | EED612 |  | INC | SELECT+1 |  |
| 3150 | IACC | 850512 | RELEND | STA | SELECT | NOW SELECT FOINTS TO ADDRESS |
| 3170 |  |  | ; |  |  | SPECIFIED BY RELATIUE |
| 3180 |  |  | ; |  |  | BRANCH INSTRUCTION. |
| 3190 | 1ACF | 28 |  | PLP |  | RESTORE CALLER' 5 DECIMAL |
| 3200 |  |  | ; |  |  | FLAG. |
| 3210 | 1RDU | 20A115 |  | JSR | PR.ADR | PRINT ADDRESS SPECIFIED |
| 3220 |  |  | ; |  |  | EY INSTRUCTION. |
| 3230 | IfD3 | 202B15 |  | JSR | POP.SL | RESTORE SELECT=ADDRESS OF |
| 3240 |  |  | ; |  |  | OPERAND. |
| 3250 | 1ADS | A201 |  | LDX |  | OPERAND HAD ONE BYTE... |
| 3260 | 1ADE | A904 |  | LDA | \#4 | AND FOUR CHARACTERS. |
| 3270 | 1ada | 60 |  | RTS |  | RETURN TO CALLER. |
| 3280 |  |  | ; |  |  |  |
| 3290 |  |  | ; |  |  |  |
| 3300 |  |  | ; |  |  |  |
| 3310 |  |  | ; |  |  |  |
| 3320 |  |  | ; |  | ZERO PA | E MODE |
| 3330 |  |  | ; |  |  |  |
| 3340 |  |  | ; |  |  |  |
| 3350 |  |  | ; |  |  |  |
| 3350 |  |  | ; |  |  |  |
| 3370 | 1RDB | ค900 | ZEROPG | LDA | 岍 | PRINT TWO ASCII ZERO'S TO |
| 3380 | 1ADD | 208314 |  | JSR | PR.BYT | ALL SELECTED EYTES. |
| 3390 |  |  | ; |  |  | (C-IP OWNERS: SUBSTITUTE NOPS |
| 3400 |  |  | ; |  |  | --EA EA EA--FOR JSR PR. BYT. |
| 3410 |  |  | ; |  |  | TO GET NARROW FORMAT. |
| 3420 | 1RED | 200819 |  | J5R | ONEEYT | PRINT ONE-BYTE OPERAND. |
| 3430 | 1 fie3 | A201. |  | LDX | \#1 | OFERAND HAS ONE BYTE... |
| 3440 | 1AES | R904 |  | LDA | \#4 | ...fild four characters. |
| 3450 |  |  | ; |  |  | (C-IP OWNERS:AS O2, |
| 3450 |  |  | ; |  |  | NOT AS E4, FOR NARROW FORIHAT.) |
| 3470 | 1PET | 60 |  | RTS |  |  |



| 4nEb |  | ；SICH STRING．TEX HAS A PSEUDO－ADDRESSING |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 4079 |  | ；MODE：TEXT MODE． |  | IN TEXT MODE，WE PRINT THE |
| 4050 |  | STRING AND RETURN |  | WITHOUT DUMPING THE LINE |
| 4090 |  | ；IN HEX．T |  | $G$ MAY be of fNY LENGTH． |
| 4100 |  | ； |  |  |
| 4110 |  | ； |  |  |
| 4120 |  | ； |  |  |
| 4130 |  | ； |  |  |
| 4148 |  | ； |  |  |
| 4150 |  | ； |  |  |
| 4160 |  | ； |  |  |
| 4178 |  | ； |  |  |
| 4180 |  | ； |  |  |
| 4150 |  | ； |  |  |
| 4こ00 |  | ； |  |  |
| 4210 | IPFE EG | TXMODE | PLA | POP RETURN ADDRESS TO |
| 4220 | 1FFF EB |  | PLA | OPERND． |
| 4230 |  | ； |  |  |
| 4240 | 150668 |  | PLA | POP RETURN ADDRESS TO |
| 4250 | 180168 |  | PLA | DSLIME． |
| 4260 |  |  |  |  |
| 4270 |  | ； |  | NOW DSLINE＇ 5 CALLER IS ON |
| 4280 |  | ； |  | THE STACK． |
| 4230 |  | ； |  |  |
| 43018 |  | ； |  |  |
| 4310 | 1802 20e317 |  | JGR NEXTSL | ADUANCE PAST TEX PSEUDO－OP． |
| 4320 | $1 \mathrm{BQ5}$ 3囚ดD |  | BMI TXEXIT | RETURN IF REACHED EA． |
| 4330 | 1EDT 203412 |  | JSR GET．SL | GET THE CHARACTER． |
| 4340 | 1EOFI C9FF |  | CMP 非ETX | IS IT END OF TEXT？ |
| 4350 | 1B6C FOBE |  | EEQ TXEXIT | IF SO，STRING ENDEN． |
| 4350 | 1 BOE 204014 |  | JER PR．CHR | IF NOT，PRINT CHARACTER． |
| 4370 | 1811 18 |  | CLC | BRANCH BACK TO GET MEXT |
| 4380 | 1512 SJEE |  | BCC TXMODE＋4 | CHARACTER． |
| 4350 |  | ； |  |  |
| 4480 |  | ； |  |  |
| 4410 | 1Ei4 207214 | TXEXIT | JSR CR．LF | ADUANCE TO A NEW LINE． |
| 4420 | 181？20e31？ |  | JSR NEXTSL | ADURNCE TO NEXT OPCODE． |
| 4430 | 1B1A EO |  | RTS | RETURN TO CALLER OF DSLINE． |
| 4440 |  | ； |  |  |
| 4450 |  | ； |  |  |
| 44E日 |  | ； |  |  |
| 4470 |  | ； |  |  |
| 4480 |  | ； |  |  |
| 4450 |  | ； |  |  |
| 4500 |  | ； |  |  |
| 4510 |  | ； |  |  |
| 45201 |  | ； |  |  |
| 4530 |  | ； |  |  |
| 4540 |  |  |  |  |
| 4550 |  | ； |  |  |
| 45E0 |  | ；TABLE OF ADDRESSING MODE SUBROUTINES |  |  |
| 4570 |  |  |  |  |
| 4580 |  |  |  |  |
| 4590 |  | ； |  |  |
| 4 EVO |  | ； |  |  |
| 4610 |  | ； |  |  |
| 4520 |  | ； |  |  |
| 4630 |  | ； |  |  |




# Appendix C8: <br> Table-Driven Disassembler (Tables) 



| 580 |  |  | ; |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 550 | $1900=$ |  | DSPAGE $=\$ 1900$ |  | STARTING PAGE OF DISASSEMELER |
| E00 1 |  |  |  |  |  |
| 610 |  |  | ; |  |  |
| E20 |  |  | ; |  |  |
| 630 |  |  | ; |  |  |
| 640 |  |  | : |  |  |
| E50 |  |  |  |  |  |
| 650 |  |  |  |  |  |
| 676 |  |  | ; LIST OF MNEMONICS |  |  |
| 680 |  |  |  |  |  |  |
| 690 |  |  |  |  |  |
| 760 |  |  |  |  |  |
| 710 |  |  | ; |  |  |
| 720 |  |  | ; |  |  |
| 730 |  |  | ; | * $=$ DSPAGE+\$250 |  |
| 74018 | 1850 |  |  |  |  |  |
| 750 |  |  | ; |  |  |
| 760 |  |  | ; |  |  |
| 77E |  |  | ; |  |  |
| 780 |  |  | ; |  |  |
| 790 |  |  | MnAmes .byte tex |  | SINCE THIS TABLE IS A STRING OF CHARACTERS, START IT WITH THE TEX PSEUDO-OP. |
| 600 | 18507 | 7F |  |  |  |  |  |
| 810 |  |  |  |  |  |  |  |
| 820 |  |  | ; |  |  |
| 850 |  |  | . BYTE ' BAD' |  |  |
| 840 | 185i | 42 |  |  |  |  |  |  |
| 840 | 1852 | 41 |  |  |  |
| 849 | 1553 | 44 | . BYTE ' ADC' |  |  |
| 850 | 1854 | 41 |  |  |  |  |  |  |
| 850 | 1855 | 44 |  |  |  |
| 858 | 1556 | 43 |  |  |  |  |  |  |
| 8 E | 1857 | 41 | . BYTE ' AND' |  |  |
| ES9 | 1856 | $4 E$ |  |  |  |
| 860 | 1E5G | 44 |  |  |  |  |  |  |
| 870 | 185A | 41 | .BYTE 'RSL' |  |  |
| 870 | 185E | 53 |  |  |  |
| 870 | 1 ESC | 4 C | . BYTE ' BCC' |  |  |
| 8 80 | 185D | 42 |  |  |  |  |  |  |
| 880 | 185E | 43 |  |  |  |
| 8 cog | 185F | 43 |  |  |  |  |  |  |
| 890 | 1EES | 42 | . EYTE ' ECS' |  |  |
| 896 | 1.261 | 43 |  |  |  |
| 896 | 1 LE 2 | 53 | . BYTE ' BEQ' |  |  |
| 900 | 1 1563 | 42 |  |  |  |  |  |  |
| 900 | 1EE4 | 45 |  |  |  |
| 900 | 1EES | 51 |  |  |  |  |  |  |
| 910 | 18EE | 42 | . EYTE ' BIT' |  |  |
| 910 | 1EE7 | 45 |  |  |  |
| 916 | 1568 | 54 | . BYTE ' BMI' |  |  |
| 520 | 1859 | 42 |  |  |  |  |  |  |
| 929 | 1EEA | 4 d | . BYTE ' Bryes |  |  |
| 928 | 1beb | 49 |  |  |  |  |  |  |
| 930 | 18EC | C 42 |  |  |  |  |  |  |
| 930 | 1560 | 4E |  |  |  |
| 930 | 18EE | 45 | . BYTE ' BPL' |  |  |
| 940 | - 1BEF | 42 |  |  |  |  |  |  |
| 540 | [1ET0 | 50 |  |  |  |



| 1140 | 1BAB 4C | . byte ' LDA' |
| :---: | :---: | :---: |
| 1140 | 1BAC 44 |  |
| 1140 | 1BAD 41 | . BYTE 'LDX' |
| 1150 | IEAE 4C |  |
| 1150 | 1EAF 44 |  |
| 1150 | 18EO 58 | - BYTE 'LDY' |
| 1160 | 1ESI 40 |  |
| 1160 | 1EB2 44 |  |
| 1160 | 18E3 59 | . BYTE 'LSR' |
| 11:0 | 1EB4 4C |  |
| 1170 | IRE5 53 |  |
| 1170 | 15ES 52 | . BYTE $^{\prime}$ NOP' |
| 1180 | 1Eb7 4E |  |
| 1180 | 1BEE 4F |  |
| 1120 | 18ES 50 | . BYTE 'ORA' |
| 1190 | 18BA $4 F$ |  |
| 1190 | 18BE 52 |  |
| 1190 | 1BEC 41 | . BYTE 'PHA' |
| 1200 | 1BED 50 |  |
| 1200 | 1BEE 48 |  |
| 1200 | 1EEF 41 | . BYTE 'PHP' |
| 1210 | 18CD 50 |  |
| 1210 | 1EC1 45 |  |
| 1210 | $1 \mathrm{BC2} 50$ | . BYTE 'PLA' |
| 1220 | 180350 |  |
| 1220 | 1EC4 4C |  |
| 1220 | 1 ECS 41 | . EYTE 'PLP' |
| 12301 | IECE 50 |  |
| 1230 | 1ECT 4C |  |
| 1230 | 18CS 50 | . BYTE 'ROL' |
| 1240 | 180952 |  |
| 1240 | 1ECA 4F |  |
| 1240 | 18CB 4C | .EYTE 'ROR' |
| 1250 | 18CC 52 |  |
| 1250 | 18CD 4F |  |
| 1250 | IECE 52 | . BYTE 'RTI' |
| 1260 | 1BCF 52 |  |
| 1260 | 1EDO 54 |  |
| 1260 | 1ED1 49 | . EYTE 'RTS' |
| 1270 | 1ED2 52 |  |
| 1270 | 1ED3 54 |  |
| 1270 | 18D4 53 | . BYTE $^{\prime}$ SEC' |
| 1280 | $1 E D 553$ |  |
| 1280 | 150542 |  |
| 1280 | 1 BD 743 | . BYte ' SEC' |
| 1290 | 150853 |  |
| 1290 | 1 E09 45 |  |
| 1290 | 1EDA 43 | . EYTE $^{\prime}$ SED' |
| 1300 | 18DE 53 |  |
| 1300 | 1BDC 45 |  |
| 1300 | 1EDD 44 | . BYTE 'SEI' |
| 1310 | 1BDE 53 |  |
| 1310 | 18DF 45 |  |
| 1310 | 1BEO 49 | . BYTE 'STR' |
| 1320 | 1RE1 53 |  |
| 1320 | 1EE2 54 |  |
| 1320 | 18E3 41 |  |
| 1330 | 1 BE 453 | . BYTE ' STX' |

```
1330 1EE5 54
1330 1BE5 58
1340 1BE7 53
1340 1BEG 54
1340 1BEG 59
1350 1EEA 54
1350 1BEB 41
1350 1BEC 5B
1350 1BED 54
13EQ 1EEE 41
1350 1BEF 5S
1370 1BFD }5
1370 1BF1 53
1370 1BF2 58
1380 1BF3 54
1380 1BF4 5%
13E0 1BF5 41
1390 1BFG }5
1390 1BF7 58
1390 1BFO 53
1400 1BFS 54
1400 1BFA 59
140Q 1BFB 41
1410 1BFC }5
1410 1BFD 45
1410 1BFE 58
1420
1430 1EFF FF
1440
1450
146\square
1470
1480
1490
1500
1510
1520
1530
1540
1550
1550
1570
1580
1590
160日
1610
1620
1630
1540
1650
1650
1E70
1589
1690
1700
1710
1720 1C00 22
1720 1C01 EA
```

```
1720 1002 01
1720 1003 01
1720 1C04 01
1720 1085 EA
1720 1006 ロA
1720 1007 01
1720 1C0日 70
1730 1C09 6?
17シ0 1CDA 日A
1730 100B 01
1730 1CQC 01
1730 1CDD 6A
1730 1CDE OA
1730 1COF D1
1740 1C1D 1F
1740 1Ci1 GA
1740 1012 01
1740 1C13 01
1740 1C14 01
1740 1C15 6A
1740 1C16 0A
1740 1C17 01
1750 1C18 2B
1750 1C19 6A
1750 1C1A 01
1750 101B01
1750 1C1C 01
1750 101D Ef
1750 1C1E QA
1750 101F 01
1760 1C20 5s
1760 1021 07
1750 1C22 01
17E0 1C23 01
17EE 1C24 IG
17E0 1C25 07
1760 1C25 79
17EQ 1C27 日1
1770 1C28 75
1770 1C29 07
1770 1C2A 79
1770 1C\angleB 01
1770 102C 15
1770 1C2D 07
1770 1C2E 79
1770 1C2F 日1
1780 1C30 19
17E0 1C31 07
1750 1032 01
1750 1033 ロ1
1780103401
1780 1035 07
1780 10.36 79
1786 1C37 01
1790 1C38 88
1790 1039 ロ7
1790 1C3A @1
1790 1CכE 01
```



```
.BYTE $1F,#6A,1,1,1,#6A,$0A,1
.BYTE $2E,$EA,1,1,1,#5A,$0A,1
.BYTE $58,7,1,1,$16,7,$79,1
.BYTE #76,7,#75,1,$16,7,$79,1
.BYTE $19,7,1,1,1,7,$79,1
.BYTE $&8,7,1,1,1,7,$79,1
```

```
1790 103C 01
1790 1C3D 07
1790 1C3E 79
1790 1C3F B1
1560 1C40 7F .BYTE $7F,$49,1,1,1,$49,$54,1
1800 1C41 49
1800 1C42 01
1200 1C43 01
1800 1C44 01
1800 1C45 49
1800 1C4E 64
1800 1C47 01
1819 1C4P 6D .BYTE $60,$49,$64,1,$55,$49,$64,1
1810 1C49 43
1815 1C4A 64
1810 1C4B Dl
1810 1C4C 55
1810 1C4D 49
1210 1C4E E4
1810 1C4F D1
18201C50 25
1820 1C51 49
1820 1052 01
1820 1053 01
1820 1C54 01
1820 1055 49
182% 1C55 64
1820 1057 @1
1830 1056 31
1830 105949
1830 1C5& 01
1830 105B D1
1830 1C5C 01
1830 1050 49
1830 1C5E 64
1830 1CSF E1
1840 1CE0 82
1840 1C61 04
1640 1CE2 D1
1840 1053 01
1840 1C54 01
1840 1CE5 04
1840 1CEE 7C
1840 1CG7 01
1850 1C58 73
1850 1059 04
1850 1C5A 7C
1E5Q 1CEB D1
1850 LCEC 55
1850 1CED 04
1850 1CSE 7C
1250 1CEF Ol
1850 1C75 28 . BYTE $28,4,1,1,1,4,$7C,1
1EEO 1C71 04
18E0 1C72 01
1860 1C73 01
1860 1074 D1
1850 1C75 04
.BYTE $25,$49,1,1,1,$49,$64,1
.BYTE $31,$49,1,1,1,$49,$64,1
.GYTE $82,4,1,1,1,4,$7C,1
.BYTE $73,4,$7C,1,$55,4,$7C,1
```

```
1850 1C76 7C
1560 1C77 01
1B7D 1C7B BE
1870 1C79 04
1870 1C7A 01
1870 1C7B 01
1870 1C7C 01
1070 1C7D 04
1870 1C7E 7C
1870 1C7F AC
1880 1080 D1
1880 1081 91
1880 1C82 B1
1880 1083 01
1880 1084 97
1880 1085 91
1880 1086 94
1880 1C87 01
1890 1088 46
1890 1C89 01.
1890 1C8A 月3
1890 108B 01
1850 1C8C 37
1890 1C8D 91
1890 1CBE 94
1890 1C8F 01
1900 1CSO 0D
1500 1091 91
1900 1C92 01
1900 1C33 01
1900 1054 97
1900 1095 91
15001 1095 94
1900 1097 01
1910 1098 AS
1910 1C99 91
1910 1CSA A3
1910 1css 01
1910 1CSC 01
1910 1C9D 91
1910 1CSE 01
1910 1CSF O1
1920 1CAD 51
1520 1CA1 5B
1920 1CAZ 5E
1920 1CA3 01
1920 1CA4 61
1920 1CF5 5B
1920 1CA6 5E
1920 1CA7 DI
1930 1CAB 90
1930 1CAG 5B
1830 1CAA 5A
1930 1CAB 01
1930 1CAC 51
1930 1CAD 5B
193D ICAE 5E
1930 1CAF 01
.BYTE $EE,4,1,1,1,4,.#FC,$RC
```



```
1940 1CB0 10
1940 1CB1 5B
1940 1CB2 D1
1940 1CB3 01
1940 1CB4 61
1940 1CE5 5B
1940 1CBG 5E
1940 1CB7 01
1950 1CB8 34
1950 1CB9 5B
1950 1CBA 9E
1550 1CBB O1
1950 1CBC 51
1950 1CBD 5B
1950 1CBE SE
1350 1CEF Q1
1560 10CD 30
19ED 1CC1 37
1960 1CC2 日1
1960 1CC3 01
1950 1CC4 3D
1960 1CC5 37
1560 1CCG 40
1360 1CC7 01
1970 10C8 52
157! 1CCS 37
1970 1CCA 43
1970 1CCB 01
1970 1CCC 3D
1970 1CCD 37
1970 1CCE 40
1970 1CCF O1
1580 1CDO 1C
1986 1CD1 37
1950 1002 01
1980 1003 01
1980 1CD4 01
1380 1CD5 37
1500 1CDG 40
1980 1CD7 01
1950 1CD8 2E
1990 1009 37
1990 1CDA 51
15S0 1CDB E1
1950 1CDC 01
1590 1CDD 37
1SSO 1CDE 40
1990 1CDF D1
20R0 1CEO 3A
2000 1CE1 85
2000 1CE2 01
2000 1CE3 01
2000 1CE4 3A
2000 1CES 85
2000 1CE6 4C
2000 1CE? 01
2010 1CES 4F
2010 1CES }8
. BYTE $10,$5B,1,1,$61,$5B,$5E,1
. BYTE $34,#5B; $9E,1,$61,$5B,$5E,1
.BYTE $3D,$37,1,1,$3D,$37,$40,1
. BYTE $52,$37,$43,1,$3D,$37,$40,1
.BYTE $1C,$37,1,1,1,$37,$40,1
.EYTE $2E,$37,1,1,1,$37.$40,1
    . BYTE $3A,$85, 1,1,$3A,$85,$4C,1
    .BYTE $4F,$85,#57,1,$3R,$85,$4C.1
```

```
2 0 1 0 ~ 1 C E A ~ 6 7 ~
2010 1CEB 01
2010 1CEC 3A
2010 1CED 55
2010 1CEE 4C
2010 1CEF R1
2020 1CFD 13
2020 1CF1 85
2020 1CF2 D1
2020 1CF3 D1
2020 1CF4 01
2020 1CF5 55
2020 1CFG 4C
2020 1CF7 D1
2030 1CF8 8B
2030 1CFG }8
2030 1CFA D1
2030 1CFB D1
2030 1CFC O1
2030 ICFD 55
2030 1CFE 4C
2030 1CFF O1
2040 ;
2050 ;
2060 ;
2070 ;
2080 ;
20.00 ;
2100 ;
2110 ;
2120 ;
2130 ;
2140 ;
2150 ;
21E0 ;
2170 ;
2180 ;
2190 ;
2200 ;
2210 ;
22200;
2230 ;
2240 ;
2250
2260
2270
2280
2290
2300
2310
2320
2330 1000 12
2330 1001 15
2330 1D62 00
2330 1D03 00
2330 1004 00
2330 1005 06
2330 1006 06
.BYTE $13,#85,1,1,1,$85,$4C,1
.BYTE $8B,$85,1,1,1,585,$4C,1
,
;
;
;
*
;
;
**************************************************
;
****************************************************
    RN ADDRESSING MODE'S CODE IS ITS OFFSET
INTO SUBS, THE TFBLE OF ADDRESSING MODE
SUBROUTINES.
;
;
;
;
;
;
    MODES . BYTE 18,22,0,0,0,6,6,0
```

```
2330 1007 00
23401050 12
2340 1005 04
23401D0A E2
23401DOE 00
2340 1DดC D0
2340 1D00 DC
2340 1D0E 日C
2340 1NOF 60
2350 1010 14
2350 1011 18
2350 1012 00
2350 1013 00
23501D14 00
2350 1015 0E
2350 1D16 0E
23501017 00
23601018 12
2360 1D19 10
23E0 1DIA DO
2350 1D1B 00
2350 1D1C D0
2360 1D1D 16
2360 1DIE IE
ZSSU 1DIF DG
2370 1n20 RC
2370 1021 15
2370 1022 00
2370 1023 00
237% 1024 06
2370 1D25 0E
2370 1D2E @5
2370 1027 口0
23E日 IDZG }1
2360 1D29 04
23E0 1D2A 02
2380 1D2E 00
2380 1D2C DC
2380 1D2D ロC
2300 1D2E 日C
2380 1D2F EO
2350103014
2390 1D31 18
2390 1032 00
2350103.3 00
23901034 00
2395 1035 00
2390 1D36 00
2330 1037 50
2400 1038 12
2406 10.35 10
2405 1D3A 00
2400 103E 00
2400 1D3C 00
2400 1D3D BE
2400 1D3E 0E
2400 1D3F 00
2410 1D40 12
```

```
.BYTE 18,4,2,0,0,12,12,0
```

.BYTE 18,4,2,0,0,12,12,0
.BYTE 20,24.0,0,0,14,14,0
.BYTE 20,24.0,0,0,14,14,0
.BYTE 18,16,0,0,0,22,22,0
.BYTE 18,16,0,0,0,22,22,0
.BYTE 12,22,0,0,6,6,5,0
.BYTE 12,22,0,0,6,6,5,0
.EYTE 18,4,2,0,12,12,12,0
.EYTE 18,4,2,0,12,12,12,0
.BYTE 20,24,0,0,0,8,B,0
.BYTE 20,24,0,0,0,8,B,0
.EYTE 15,16,0,0,0,14,14,0
.EYTE 15,16,0,0,0,14,14,0
.BYTE 18,Z2,0,0,0,6,E,0

```
.BYTE 18,Z2,0,0,0,6,E,0
```

```
2410 1041 15
2410 1D42 0% 
2410 1043 60
2410 1D44 00
2410 1045 05
2416 1D45 05
2410 1D-47 00
2420 1D4% 12
2420 1545 日C
2420 154月 E2
2420 1D4B 00
2420 1D4C 0C
2420 1D4D DC
2420 1D4E EC
2420 1D4F 00
2430 155014
24301051 18
2430 1052 00
2430 1053 20
2430 155400
2430 1D55 08
2430 1055 05
24301057 00
2440 1D5S 12
244ण 1D5910
2440 1D5A D0
2440 105E 00
2440 1050 B0
2446 1DST 位
2440 1DSE OE
2440 1DEF Q0
2450 10EE 12
2450 1DE1 16
2450 117E2 00
2450 1053 00
2450 1064 00
2450 10E5 05
2450 1DES 0G
2450 1DE? 账
2460 11588 12
24E0 108S B4
2450 1DEH G2
24E0 1DE昌 00
24ED 1DEC 1A
24EG 1DED DC
24E9 1DEE DC
24EO 1DEF DO
2470 1D70 14
2470 1D71 15
2470 1D72 60
2470 1D73 00
2470 1D7400
2470 1D75 08
2476 1DTE D8
247D 1D77 日G
2480 1078 12
2400 1079 10
2480 1D7A B0
.BYTE 20,24,0,0,0,8,8,0
.BYTE 18,1E,D,0,0,14,14,0
.BYTE 18,22,0,D,0,E,6,0
.EYTE 18,4,2,0,26,12,12,0
.BYTE 20,24,0,0,0,8,8,0
.BYTE 18,16,0,0,0,14,14,28
```

```
.BYTE 18,12,2,0,12,12,12,0
```

```
.BYTE 18,12,2,0,12,12,12,0
```

```
2480 107B 00
2400 1D7C 00
2490 1DTD DE
2480 1DTE DE
2480 107F 1C
24501
2500 1DED 00
2500 1081 15
2500 1082 00
2500 1083 00
2500 1D84 05
2505 1D85 65
2500 1D85 05
2500 1DET 00
2510 1088 12
2510 1D83 00
2510 1DBA 12
2510 1DEB EG
2510 1DEC 日C
2510 1D8D 0C
2510 1DSE 0C
2510 15SF 00
2529 12S0 14
2520 10S1 18
2520 1DS2 吅
2520 1093 00
2520 1n94 08
2520 1n95 88
2520 1D96 ロค
2520 1B9? 00
25301098 12
2530 1093 10
2530 10SA 12
2530 1DSE 00
2530 1DGC 00
2530 1D9D OE
2530 1DSE 00
2530 1DSF DO
2540 1DRD 04
2540 1DA1 16
2540 1DA2 04
2540 1nA3 00
2540 1DR4 05
2540 1EA5 D6
2540 1DAS DE
2540 1DAT DO
2550 10A8 12
2550 1DA9 04
2550 1DAR 12
2550 1DAB 00
2550 1DAC OC
2550 1DAD SC
2550 1DAE RC
2550 1DAF DO
25E0 1DEO 14
25E[ 1DE1 18
2560 1DEL B0
25E0 1023 00
.BYTE 18,0,18,0,12,12,12,0
.BYTE 20,24,0,0,8,8,10,0
.BYTE 1巴,15,15,0,0,14,0,0
.BYTE 4,22,4,0,5,6,6,D
.BYTE 1B,4,18,0,12,12,12,0
.BYTE 20,24,0,0,8,8,10,0
```

```
.BYTE R,Z2,D,D,G,G,6,0
```

```
.BYTE R,Z2,D,D,G,G,6,0
```

```
2550 1DB4 08
2560 1085 08
2560 1DEE DR
2550 10B7 日ด
2570 1DES 14
2570 1DE9 10
2570 1DEA 12
2570 1DEB DO
2570 1DEC DE
2570 1DBD EE
2570 1DEE 10
2570 1DBF [0
2580 1DC0 04
2500 1DC1 15
2580 1DC2 E0
2580 1DC3 DD
2560 1DC4 05
2530 1DC5 06
2500 1DC6 06
2580 10C7 00
2590 10C8 12
2590 1DC9 04
2550 1DCA 12
2550 1DCB DO
2550 1DCC DC
25S0 1DCD DC
2590 1DCE DC
2590 1DCF GD
2E00 1DDO 14
2600 10D1 18
2500 10D2.00
2500 1003 E0
2600 10D4 00
2600 1005 08
2500 1DDG 05
2600 1DD7 ロ0
2610 10n6 12
2610 1009 10
2610 10DA 狍
2610 1DDS g0
2610 10DC 00
2610 1DDD DE
2610 1DDE GE
2610 1DDF D0
2620 1DE0 04
2620 1DE1 16
2520 1DEZ Dด
2E2D 1DE3 ED
2520 1DE4 EO
2520 1DE5 06
2620 1DEG DG
2620 1DE7 00
2630 1DES 12
2630 1DES D4
2539 1DEA 12
2530 1DEB B0
2630 1DEC BC
2630 1DED EC
```

26301 DEE DC
2630 1DEF DO 2540 1DFO 14 2640 1DF1 18 26401 DFZ 日 26401 DF 3 DO 2640 1DF4 00 $26401 \mathrm{DF5} \mathrm{BB}$ 2640 1DF6 08 $26401 D F 700$ 26501 DFg 12 26501 DFG 10 2650 1DFA 0 2650 1DFE 0 26501 DFC CO 2550 1DFD DE 2650 1DFE 日E 2650 1DFF DO

```
.BYTE 20,24,0,0,0,8,8,0
```

．BYTE 18，16．0，日，0，14，14，0

Appendix C9:
Move Utilities

| 10 | ; | APPEMDIX CG: | ASSEMBLER LISTING OF |
| :---: | :---: | :---: | :---: |
| 29 | ; |  | MOVE UTILITIES |
| 30 | ; |  |  |
| 40 | ; |  | . |
| 50 | ; |  |  |
| 60 | ; | SEE CHAPTER | 10 OF EEYOND GAMES: SYSTEMS |
| 70 | ; | SOFTWARE FOR YOUR | 6502 PERSONAL COMPUTER.: |
| ED | ; |  |  |
| 90 | ; |  | $\therefore$ BY KEN SKIER |
| 100 | ; |  |  |
| 110 | ; |  |  |
| 120 | ; |  |  |
| 130 | ; |  |  |
| 140 | ; |  |  |
| 159 | ; |  |  |
| 160 | ; |  | ************************** |
| 170 | ; |  |  |
| 180 | ; | CONSTANTS |  |
| 190 | ; |  |  |
| 2001 | ; |  | ******************************** |
| 210 | ; |  |  |
| 220 | ; |  |  |
| 230 |  |  |  |
| 240 | ; |  |  |
| 250 | ; |  |  |
| 259 00ด】= |  | $C R=\$ 0 \mathrm{D}$ | CARRIAGE RETURN. |
| 270 D06F= |  | $\underline{L F}=$ TUA | LINE FEED. |
| $2800 \mathrm{GTF}=$ |  | TEX=年TF | START OF TEXT CHARACTER. |
| $290 \mathrm{BQFF}=$ |  | $E T X=\$ F F$ | END OF TEXT CHARACTER. |
| 300 | ; |  |  |
| 310 | ; |  |  |
| 320 | ; |  |  |
| 330 | ; |  |  |
| 340 | ; |  |  |
| 359 | ; |  |  |
| 350 | ; |  |  |
| 370 | ; |  |  |
| 380 | ; | EXTERNAL | ADDRESSES |
| 390 | ; |  |  |
| 400 | ; |  |  |
| 418 | ; |  |  |
| 420 | ; |  |  |
| 430 | ; |  |  |
| 440 | ; |  |  |
| 450 | ; |  |  |
| 469 | ; |  |  |
| 470 | ; |  |  |
| 480 | ; |  |  |
| $4501200=$ |  | UMPAGE $=1200$ | STARTING FAGE OF UISIBLE |
| 506 | ; |  | MONITOR CODE. |
| 510 | ; |  |  |
| $5201205=$ |  | SELECT=UMPAGE + |  |
| $5301207=$ |  | UISMOM= WMPAGE+ |  |
| 540 | ; |  |  |
| 550 | ; |  |  |
| 560 | ; |  |  |
| $5701409=$ |  | PRPAGE $=\$ 1400$ | STARTING FAGE OF PRINT COD |
| 580 | ; |  |  |


| 590 | 1408= | TUT. ON=PRPAGE+8 |  |  |
| :---: | :---: | :---: | :---: | :---: |
| E00 | 14E4= | FRINT: =PRPAGE+\#E4 |  |  |
| 610 | 1512= | FUSHSL=PRPAGE + + 112 |  |  |
| 620 | $152 \mathrm{~B}=$ | POP. SL=PRPRGE + \$12B |  |  |
| E30 |  | ; |  |  |
| 640 |  | ; |  |  |
| 650 | $1500=$ | HEX. $\mathrm{PG}=\$ 1500$ |  | ADDRESS OF PAGE IN WHICH |
| 6E0 |  | ; |  | HEXDUMP CODE STRRTS. |
| 670 |  | ; |  | (HEXDUMP CODE STARTS AT |
| E80 |  | ; |  | \$1550, BUT IT'S EASIER TO |
| 690 |  | ; |  | COUNT FROM \$1500.) |
| 700 |  | ; |  |  |
| 710 | 15ES= |  | SETADS $=$ HEX. PG + | Es |
| 720 |  | ; |  |  |
| 730 |  | ; |  |  |
| 740 |  | ; |  |  |
| 750 |  | ; |  |  |
| 760 |  | ; |  |  |
| 770 |  | ; |  |  |
| 780 |  | ; |  |  |
| 790 |  | ; |  |  |
| 800 |  | ** |  |  |
| 810 |  | ; |  |  |
| 820 |  | ; | UARIABLES |  |
| 830 |  | ; |  |  |
| 840 |  | ** |  | *************************** |
| 850 |  | ; |  |  |
| 860 |  | ; |  |  |
| 878 |  | , |  |  |
| 880 |  | ; |  |  |
| 890 |  | ; |  |  |
| 900 | 1780 |  | * $=$ \% 17 Ba |  |
| 910 |  | ; |  |  |
| 920 |  | ; |  |  |
| 930 | $1552=$ |  | SA=HEX. PG $+\$ 52$ | POINTER TO START ADDRESS |
| 940 |  | ; |  | OF ELOCK TO BE MOUED. |
| 950 |  | ; |  |  |
| 960 | 1554= |  | $E A=5 A+2$ | POINTER TO END OF BLOCK TO |
| 970 |  | ; |  | BE MOUED. |
| 1000 |  | ; |  |  |
| 1010 | 17B0 0000 | NUM | . WORD $\square$ | MUMBER OF BYTES IN BLOCK |
| 1020 |  | ; |  | TO BE MOUED. ZERO MERNS |
| 1030 |  | ; |  | BLOCK CONTAINS 1 BYTE. |
| 1040 |  | ; |  |  |
| 1856 |  | ; |  |  |
| 1060 | 17820000 | DEST | .WORD $\square$ | FOINTER TO ELOCK' 5 |
| 1079 |  | ;' |  | DESTINATION. |
| 1020 |  | ; |  |  |
| 1090 |  | ; |  |  |
| 1100 |  | ; |  |  |
| 1110 |  | ; |  |  |
| 1120 |  | ; |  |  |
| 1130 |  | ; |  |  |
| 1140 |  | ; |  |  |
| 1150 | Q000 $=$ |  | GETPTR=0 | THESE THO "PFige fointers |
| 1150 | -0002= |  | PUTPTR=GETPTR | 2 GET RIND PUT bYtes. |
| 1170 |  | ; |  |  |
| 1150 | . | ; |  |  |




| 2180 |  |  | ; |  |  | MOVE-DOWN. |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2150 | 1.807 | RDS215 |  | LDA | SA |  |
| 2280 | 180A | CDE217 |  | CMP | DEST |  |
| 2210 | $180 \square$ | 9036 |  | BCC | MOVEUP |  |
| 2220 | 180F | dage |  | BNE | MOUEDN | IF DEST=5R, |
| 2230 | 18.11 | F000 | OK.RTN | LDY | \# | RETURN BEARING "OKAY" CODE. |
| 2240 |  |  | ; |  |  | RESTORE ZERO PAGE BYTES |
| 2250 | 1513 | 58 | LOOP. 2 | PLA |  | THAT WERE CHANGED. |
| 2250 | 1814 | 590000 |  | STA | GETPTR, Y |  |
| 2270 | 1517 | C8 |  | INY |  |  |
| 2280 | 1816 | CEB4 |  | CPY | \#4 |  |
| 2290 | 181A | DGF7 |  | BNE | LOOP. 2 |  |
| 2300 | 131C | figff |  | LDA | \#OKAY | RETURN W/"OKAY" CODE. |
| 2310 | 161E | 60 |  | RTS |  |  |
| 2320 |  |  | ; |  |  |  |
| 2530 |  |  | ; |  |  |  |
| 2340 |  |  | ; |  |  |  |
| 2350 | 181F | 200418 | MOUEDM | JSR | LOPAGE | SET PAGE POINTERS TO LOWEST |
| 2350 |  |  | ; |  |  | PAGES IM ORIGIN, DESTINATION BLOCKS. |
| 2376 |  |  | ; |  |  |  |
| 2380 |  |  | ; |  |  |  |
| 2390 | 1822 | F000 |  | LDY | \# 0 | INITIALIZE PAGE INDEX TO |
| 2400 |  |  | ; |  |  | BOTTOM OF PRGE. |
| 2410 |  |  | ; |  |  |  |
| 2420 | 1824 | AEE117 |  | LDX | NUM+1 | USE $\times$ TO COUNT THE NUMBER |
|  |  |  | ; |  |  | OF PAGES TO MOUE. MORE THAM |
|  |  |  | ; |  |  | ONE PAGE TO MOVE? |
| 2430 | 1827 | fode |  | EEQ | LESSDM | IF NOT, MOUE LESS THRN A |
| 2440 |  |  | ; |  |  | PAGE. |
| 2450 |  |  | ; |  |  |  |
| 2460 |  |  | ; |  |  | IF 50, |
| 2470 | 1829 | B100 | PAGEDM | LDA | (GETPTR), Y | move a page down, |
| 2483 | 1828 | 9102 |  | STA | (PUTPTR), Y | STARTING AT THE BOTTOM. |
| 2490 | 122D | Ca |  | INY |  | INCREMENT PAGE INDEX. |
| 2500 | 162E | DOF9 |  | BNE | PAGEDN | If PAGE NOT MOUED, MOVE |
| 2510 |  |  | ; |  |  | NEXT BYTE... |
| 2520 |  |  | ; |  |  |  |
| 2530 | 1530 | E6D1 |  | INC | GETPTR+1 | INCREMENT PAGE POINTERS. |
| 2545 | 1832 | ES93 |  | INC | PUTPTR+1 |  |
| 2550 | 1834 | CA |  | DEX |  | DECREMENT PAGE COUNT. |
| 2550 | 1835 | Defz |  | BNE | Pagedn | If a page left to move, |
| 2570 |  |  | ; |  |  | MOVE IT RS A PRGE. |
| 25E0 |  |  | ; |  |  |  |
| 2590 | 1837 | EB | LESSDN | DEY |  |  |
| 2690 | 1838 | C8 |  | INY |  | MOUE LESS THAN A PAGE |
| 2510 | 1839 | E100 |  | LDA | (GETPTR),Y | DOWN. STARTING AT THE |
| 2620 | 183B | 9102 |  | STA | (PUTPTR), Y | BOTTOM. |
| 2530 | 1835 | cced 17 |  | CPY | NuM | MOVED LAST BYTE? |
| 2640 | 1840 | DEF6 |  | BHE | LESSDN+1 | IF NOT, MOVE NEXT BYTE... |
| 2650 | 1842 | $4 \mathrm{Cl118}$ |  | JMP | OK. RTN | IF SO, RETURN BEARING |
| 2650 |  |  | ; |  |  | "OKAY" CODE. |
| 2670 |  |  | ; |  |  |  |
| 2 ESO |  |  | ; |  |  |  |
| 2690 |  |  | ; |  |  |  |
| 2700 | 1545 | PDB117 | MOUEUP | LDA | $\mathrm{NuM+1}$ | MORE THAN A PAGE TO MOUE? |
| 2710 | 1848 | r048 |  | BEO. | LESSUP | If NOT, MOUE LESS THAN A |
| 2720 |  |  | : |  |  | PAGE. |



```
2850 184A ACB117
2890 184D ADD017
2900 1850 38
2910 1851 EgFF
2920 1853 B001
2930 1855 88
2540 1856 AA
2950 ;
2960 ;
2970 ;
2980 ;
2950 :
3000 1857 6403 STY PUTPTR+1
3010 1859 8f
3020 185A 18
3030 185B 605215
TXA
    CLC
    ADC 5A
    STA GETPTR
    BCC NEXT.Z
    INY
3060 1862 c8
3070 ;
3080 ;
3090 1863 98 NEXT.2 TYA
3100 1864 5D5315 RDC 5A+1
3110 1857 8501 STA GETPTR+1
3120 ;
3130 ; PTR=5A+NUM-$FF.
3140 ;
3150 ;
3150 ;
3170 1869 BA TXA
3180 185A 18 CLC
3190 186B 6DB217 ADC DEST
3200 186E 5502 STA PUTPTR
3210 1E70 9002 BCC NEXT.3
3220 1872 E603 INC PUTPTR+1
3230 ;
3240 ;
3250 1874 A503 NEXT.3 LDA PIJTPTR+1
3250 1876 5DB317 ADC DEST+1
3270 1879 5503 STA PUTPTR+1
3280 ;
3250 ;
3300 ;
LDY NUM+1
```

LDA NUM SEC
SBC \#\#FF
BCS NEXT. 1
BEY
NEXT. 1 TAX
;
;
;
;
STY PUTPTR+1
TXA
CLC
STA GETPTR
BCC NEXT. 2 INY

## ;

NEXT. 2 TYA
RDC $5 A+1$ STA GETPTR+1
$P T R=5 A+N U M-\$ F F$.

TXA
CLC
ADC DEST STA PUTPTR BCC MEXT. 3 INC PUTPTR+1

TO MOUE MORE THAN A PRGE, SET PRGE FOINTERS TO HIGHEST PFGGES IN ORIGIN, DESTINATION BLOCKS.

TO DO THIS, FIRST
SET $(X, Y)=$ NUM - \$FF. (RELATIUE ADDRESS OF HIGHEST PAGE IN A ELGCK.)

NOW ( $X, Y$ ) - NUM - \$FF. $X$ IS LOW BYTE, $Y$ IS HIGH BYTE
(LAST PAGE IN SOLRCE ELOCK.)

NOW PUTFTR=DEST+MUM-\$5FF. CLAST PRGE IN DEST BLOCK. 3

| 3310 |  |  | ； |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 3320 |  |  | ； |  |  |  |
| 3330 | 18TB | AEE117 |  | LDX | NuM +1 | LOAD $\times$ WITH NUMBER OF |
| 3340 |  |  | ； |  |  | FAGES TO MOVE． |
| 3350 |  |  | ； |  |  |  |
| 3360 | 187E | AGFF | PAgelup | LDY | \＃\＄FF | SET PAGE INDEX TO TOP OF |
| 3370 |  |  | ； |  |  | PAgE． |
| 3380 | 1800 | B100 | LOOP． 3 | LDA | （GETPTR），Y | MOVE A PAGE UP，STARTING |
| 3390 | 1882 | 5162 |  | STA | （PUTPTR），Y | AT THE TOP OF THE BLOCK． |
| 3400 | 1854 | 88 |  | DEY |  | DECREMENT FAGE INDEX． |
| 3410 |  |  | ； |  |  | ABOUT TO MOUE LAST BYTE |
| 3420 |  |  | ； |  |  | IN PRGEE？ |
| 3430 | 1865 | DEFS |  | BME | LOOP． 3 | IF NOT，HANDLE NEXT BYTE． |
| 3448 |  |  | ； |  |  | AS BEFORE． |
| 3450 |  |  | ； |  |  |  |
| 3450 |  |  | ； |  |  |  |
| 3470 |  |  | ； |  |  |  |
| 3400 | 1897 | 8100 |  | LDA | （GETPTR），Y | IF SO，MOUE THIS BYTE FROM |
| 3490 | 1869 | 9182 |  | STA | （PUTPTR），Y | SOURCE TO DESTINATION． |
| 3506 | 1888 | c601 |  | DEC | GETPTR＋1 |  |
| 3510 | 18SU | CED3 |  | DEC | PUTPTR＋1 | DECREMENT PAGE POINTERS． |
| 3520 | 188F | CA |  | DEX |  | DECREMENT PAGE COUNTER． |
| 3530 | 1890 | DGEC |  | BINE | PAGEİP | If a page left to moue． |
| 3540 |  |  | ； |  |  | MOVE IT AS A PAGE．．．． |
| 3550 |  |  | ； |  |  |  |
| 3560 |  |  | ； |  |  |  |
| 3578 | 1892 | $20 \cap 418$ | LESSUP | JSR | LGPRGE | MOUE LESS THAN A PRGE UP， |
| 3580 | 1895 | Ficeol |  | LDY | NUM | StARTING GIT THE TOP． |
| 3590 |  |  | ； |  |  |  |
| 3610 | 1898 | 8150 | MOVE． 6 | LDA | （GETPTR），Y | COPY A BYTE FROM ORIGIN |
| 3620 | 185f | 9102 |  | STA | （PIUTPTR），Y | TO DESTINATION． |
| 3630 | 1890 | 8 E |  | DEY |  | DECREMENT PAGE INDEX． |
| 3546 | 1890 | COFF |  | CPY | \＃\＃FF | COPIED THE LAST EYTE？ |
| 3650 | 188F | DQF？ |  | BNE | HOVE． 5 | IF NOT，HANDLE AS BEFORE．．． |
| 3650 | 1891 | $4 \mathrm{Cl118}$ |  | JMP | OK．RTN | IF SO，RETURM BEARING |
| 3670 |  |  | ； |  |  | ＂OKAY＂CODE． |
| $3 E E 0$ |  |  | ； |  |  |  |
| 3696 |  |  | ； |  |  |  |
| 3700 |  |  | ； |  |  |  |
| 3710 |  |  | ； |  |  |  |
| 3720 |  |  | ； |  |  |  |
| 3730 |  |  | ； |  |  |  |
| 3749 |  |  | ； |  |  |  |
| 3750 |  |  | ； |  |  |  |
| 3760 |  |  | ； |  |  |  |
| 3770 |  |  | ； |  |  |  |
| 3780 |  |  | ； |  |  |  |
| 3750 |  |  | ；枵为米 | 絲为米米 |  |  |
| 3800 |  |  | ； |  |  |  |
| 3815 |  |  | ； |  | T FAGE POIN | TEES TO BOTTOM OF |
| 3825 |  |  | ； |  | ORIGIM，DES | TINATION BLOCKS． |
| 3830 |  |  |  |  |  |  |
| 3840 |  |  | ；＊＊＊＊ |  |  | ＊＊＊ |
| 3850 |  |  | ； |  |  |  |
| 3850 |  |  | ； |  |  |  |
| 3870 |  |  | ； |  |  |  |
| 3580 |  |  | ； |  |  |  |
| 3890 |  |  | ； |  |  |  |



4320 18D0 4E
4320160120
4320180241
$432018034 E$ 43ट0 180444 4320180520 4320180550 4320180752 4320180845 4320180953 4320 15DA 53 4320180820 4320 1200 51 4320 16DD 2 E 4330 18DE FF . EYTE ETX 4349 16DF 200712 JSR UIGMON LET USER SET AN RDDRESS. $435018 E 2$ ADES 12 DAHERE LDA SELECT SET DEST=SELECT.
4360 LBES EDB217 STA DEST 4370 10ES RDOS 12 LDA SELECT 12 430 CB 18EE EDB317 STA DEST +1 4330 ;
44013 1SEE SO RTS

## Appendix CIO:

## Simple Text Editor (Top Level and Display Subroutines)

| 10 | ; | APPENDIX CID: ASSEMBLER LISTING OF |
| :---: | :---: | :---: |
| 20 | ; | A SIMPLE TEXT EDITOR |
| 30 | ; | tof Level and display subroutines |
| 40 | ; |  |
| 50 | ; |  |
| ED | ; |  |
| 70 | ; |  |
| 80 | ; |  |
| 50 | ; | SEE CHRPTER 11 OF BEYOND GAMES: SYSTEMS |
| 100 | ; | SOFTWRRE FOR YOUR 6502 FERSONAL COMPUTER |
| 110 | ; |  |
| 129 | ; |  |
| 130 | ; | BY KEN SKIER |
| 140 | ; |  |
| 159 | ; |  |
| 160 | ; |  |
| 179 | ; |  |
| 180 | ; |  |
| 190 | ; |  |
| 200 | ; |  |
| 210 | ; |  |
| 220 | ; |  |
| 230 | ; |  |
| 240 | ; |  |
| 250 | ; |  |
| 260 | ; |  |
| 270 | ; | CONSTRNTS |
| 280 | ; |  |
| 250 | ; | ********************************************* |
| 300 | ; |  |
| 310 | ; |  |
| 320 | ; |  |
| 330 | ; |  |
| 346 | ; |  |
| $350 \mathrm{0ab1}=$ |  | $C R=\$ 0 D$ CARRIAGE RETURN. |
| 360 | ; |  |
| 370 009i= |  | $L F=\Phi \boxtimes A \quad$ LINE FEED. |
| 380 | ; |  |
| 350 | ; |  |
| $400907 F=$ |  | TEX $=$ \$7F $\quad$ THIS CHARACTER MUST START |
| 410 | ; | finy Messhge. |
| 420 | ; |  |
| 430 $\mathrm{nGFF}=$ |  | ETX $=$ \$FF $\quad$ THIS CHARACTER MUST END |
| 440 | ; | ANY MESSAGE. |
| 450 | ; |  |
| 400 50.49= |  | INSCHR $=$ I GRAPHIC FOR INSERT MODE |
| 470004F= |  | OURCHR=' 0 GRAPHIC FOR OUERSTRIKE MODE. |
| 480 | ; |  |
| 496 | ; |  |
| 509 | ; |  |
| 510 | ; |  |
| 520 | ; |  |
| 530 | ; |  |
| 540 | ; |  |
| 550 | ; |  |
| 560 | ; | ******************************************** |
| 570 | ; |  |
| 580 | ; | EXTERNAL ADDRESSES |


| 550 | ; |  |
| :---: | :---: | :---: |
| 600 |  |  |
| 510 | ; |  |
| 628 | ; |  |
| E30 | ; |  |
| $6400000=$ |  | TU. PTR=0 POINTER TO A SCREEN ADDRESS. |
| 650 1000= |  | PARAMS $=1000$ SYSTEM DATA BLOCK. |
| 6ED | ; |  |
| 670 | ; |  |
| Ee0 1093= |  | TUCOLS=PARAMS+3 |
| 690 1004= |  | TUROWS=PARAMS+4 |
| $7001007=$ |  | ARROLS=PARAMS+7 |
| 710 | ; |  |
| 720 | ; |  |
| 730 | - |  |
| $7401100=$ |  | TUSUBS $=\$ 1100$ |
| $7501113=$ |  | CLR. $\mathrm{XY}=$ TUSUSS + \$13 |
| $760112 \mathrm{~B}=$ |  | TUHOME=TUSUBS + \$2B |
| 770 113C= |  | TUTOXY=TUSUBS+\#3C |
| 780 1175= |  | TUDOWM=TUSUBS + \$56 |
| $790117 \mathrm{~F}=$ |  | TUSKIP = TUSUBS + \$57F |
| $5001181=$ |  | TUPLUS=TUSUBS + \$81 |
| $8101198=$ |  | TU. PUT=TUSUS + + $\$ 98$ |
| 82011 A3 $=$ |  | ULSYTE=TUSUBS+\#A3 |
| $8301104=$ |  | TUPUSH=TUSUBS + \#C4 |
| 840 1103= |  | TU. POP=TUSUBS + \$D3 |
| 850 | ; |  |
| 860 | ; |  |
| $8701200=$ |  | $\begin{aligned} & \text { UMPAGE }=\$ 1200 \quad \text { STFRTING PAGE OF UISIBLE } \\ & \text { MONITOR CODE. }\end{aligned}$ |
| 880 | ; |  |
| $8901205=$ |  | SELECT=UMPAGE +5 |
| 980 1234= |  | GET. SL=UMPAGE+\$94 |
| $9101300=$ |  | INC. SL=UMPRGE + \#10D |
| $520131 \mathrm{~A}=$ |  | DEC. SL=UMPAGE + \# 11 A |
| 938 | ; |  |
| 940 | ; |  |
| S50 1400 $=$ |  | PRPFGGE $=\$ 1400 \quad$ STARTING PAGE OF PRINT UTILITIES. |
| 960 | ; |  |
| 970 1408= |  | TUT. ON=PRPAGE+8 |
| $580140 \mathrm{E}=$ |  | TUTOFF $=$ PRPAGE + \$RE |
| 990 1414= |  | PR.ON =PRPAGE + \$ 14 |
| $1000141 \mathrm{~A}=$ |  | PR. OFF $=$ PRFAGE + \$1R |
| 1010 1440= |  | PR. CHR=PRPAGE + \$ 40 |
| 1020 14E4= |  | PRINT: =PRPAGE+ + E 4 |
| $10301512=$ |  | PUSHSL=PRPRGE + \$112 |
| $1040152 \mathrm{~B}=$ |  | FOF. SL=PRFAGE+\$12E |
| 1050 | ; |  |
| 1060 | ; |  |
| $10701500=$ | ; | HEX.PG=\#150D ADDRESS OF PAGE IN WHICH HEXDUMP CODE STARTS. |
| 1090 | ; |  |
| 1100 1552= |  | SA=HEX. PG+\$52 |
| $11101554=$ |  | $E R=5 f+2$ |
| 1120 15E9= |  | SETADS=HEX.PG+\$ES |
| $11301783=$ |  | NEXTSL $=$ HEK.FG+\$2B3 |
| $114017 A 8=$ |  | GOTOSA $=\mathrm{HEX} \cdot \mathrm{PG}+\Phi 2 \mathrm{RO}$ |
| 1150 |  | ; |
| 11EO |  | ; |




| 2100 | 1E3A | 202811 |  | JSR | TUHOME | SET HOME POSITION OF EDIT |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2110 |  |  | ; |  |  | DISPLFAY. |
| 2120 |  |  | ; |  |  |  |
| 2130 |  |  | ; |  |  |  |
| 2140 | 1E3D | AE0310 |  | LIX | TUCOLS | CLEAR THREE ROWS FOR |
| 2150 | 1E40 | f003 |  | LDY |  | THE EDIT DISPLAY. |
| 2160 | 1E42 | 201311 |  | JSR | CLR.XY |  |
| 2170 |  |  | ; |  |  |  |
| 2180 |  |  | ; |  |  |  |
| 2190 | $1 E 45$ | 202811 |  | JSR | TUHOME | RESTORE TU.PTR TO HOME |
| 2200 |  |  | ; |  |  | POSITION OF EDIT DISPLAY. |
| 2210 | 1 E 4 B | 297611 |  | JSR | TUDOWN | SET TU.PTR TO EEGINNING |
| 2220 | 1E4B | 20C411 |  | JSR | TUPUSH | OF LINE TWO AND SfUE IT. |
| 2230 | 1E4E | 205E1E |  | JSR | LINE. 2 | display text in line two. |
| 2240 |  |  | ; |  |  |  |
| 2250 |  |  | ; |  |  |  |
| 2260 | $1 E 51$ | 200311 |  | JSR | TU.POP | SET TU.PTR TO BEGINNING OF |
| 2270 | 1E54 | 207611 |  | JSR | TUDOWN | OF THIRD LINE OF EDIT |
| 2280 |  |  | ; |  |  | DISPLAY. |
| 2290 | $1 E 57$ | $20891 E$ |  | JSR | LINE. 3 | DISPLAY THIRD LINE OF EDIT |
| 2300 |  |  | ; |  |  | DISPLAY. |
| 2310 |  |  | ; |  |  |  |
| 2320 | 1ESA | 200311 |  | JSR | TU.POP | RESTORE ZERO PAGE BYTES USED. |
| 2330 | 1ESD | 60 |  | RTS |  | RETURN TO CRLLER, WITH EDIT |
| 2340 |  |  | ; |  |  | OISPLAY ON SCREEN, REST OF |
| 2350 |  |  | ; |  |  | SCREEM UNCHRNGED, RNDD ZERO |
| 2350 |  |  | ; |  |  | FRGE FRESERUED. |
| 2370 |  |  | ; |  |  |  |
| 2380 |  |  | ; |  |  |  |
| 2390 |  |  | ; |  |  |  |
| 2400 |  |  | ; |  |  |  |
| 2410 |  |  | ; |  |  |  |
| 2420 |  |  | ; |  |  |  |
| 2430 |  |  | **** | ******** | ******** |  |
| 2440 |  |  | ; |  |  |  |
| 2450 |  |  | ; |  | DISPLA | TEXT LINE |
| 2450 |  |  | ; |  |  |  |
| 2470 |  |  | * ${ }_{\text {为 }}$; | ****** | ******** |  |
| 2480 |  |  | ; |  |  |  |
| 2490 |  |  | ; |  |  |  |
| 2500 |  |  | ; |  |  |  |
| 2516 |  |  | ; |  |  |  |
| 2520 |  |  | ; |  |  |  |
| 2530 | IESE | 201215 | LINE. 2 | JSR | PUSHSL | SAUE SELECT POIMTER. |
| 2540 | 1E61 | ค 40310 |  | L.DA | TUCOLS | SET $\times$ EQUPL TO |
| 2550 | $1 E 64$ | 4A |  | LSR | ค | HALF THE WIDTH |
| 2560 | 1E65 | AA |  | TAX |  | OF THE ECREEM. |
| 2570 | $1 E 65$ | $C A$ |  | DEX |  |  |
| 2580 | 1E67 | CA |  | DEX |  |  |
| 2590 |  |  | ; |  |  |  |
| 2500 | 1 ESB | 201A13 | LOOF. 1 | JSR | DEC.SL | DECREMENT SELECT... |
| 2610 | 1E6B | CA |  | DEX |  |  |
| 2520 | LEGC | LDFA |  | BPL | LOOP. 1 | . . X TIMES. |
| 2630 |  |  | ; |  |  |  |
| 2640 | 1ESE | AD0310 |  | LDA | TUCOLS | INITIRLIEE COUNTR. |
| 2650 | $1 E 71$ | 80001E |  | STA | COUNTR | (WE'LL DISPLAY TUCOLS |
| 2660 |  |  | ; |  |  | CHPRACTERS.) |
| 2670 | $1 E 74$ | 209412 | LOOP. 2 | JSR | GET. SL | GET A CHARACTER FROM BUFFER. |



## Appendix CII:

## Simple Text Editor (EDITIT Subroutine)

| 10 | ; | APFENDIX C11: ASSEMELER LISTING OF |
| :---: | :---: | :---: |
| 20 | ; | A SIMPLE TEXT EDITOR |
| 30 | ; | EDITIT SUBROUTINE |
| 45 | ; |  |
| 50 | ; | , |
| EG | ; |  |
| 70 | ; |  |
| 50 | ; |  |
| 90 | ; | SEE CHRPTER 11 OF BEYOND GAMES: SYSTEMS |
| 100 | ; | SOFTINARE FOR YOUR 6502 PERSONFL COMPUTER |
| 115 | ; |  |
| 120 | ; |  |
| 130 | ; | BY KEN SKIER |
| 140 | ; |  |
| 150 | ; |  |
| 150 | ; |  |
| 175 | ; |  |
| 180 | ; |  |
| 190 | ; |  |
| 200 | ; |  |
| 210 | ; |  |
| 220 | ; |  |
| 230 | ; |  |
| 240 | ; |  |
| 250 | ; |  |
| 250 | ; |  |
| 270 | ; | CONSTANTS |
| 290 | ; |  |
| 250 | ; | *********************************************** |
| 300 | ; |  |
| 310 | ; |  |
| 320 | ; |  |
| 330 | ; |  |
| 340 | ; |  |
| 350 009n= |  | $C R=$ IBD $\quad$ CARRIAGE RETURN. |
| 360 | ; |  |
| $3700008=$ |  | $L F=\$ 0 \mathrm{~A}$ ( LINE FEED. |
| 380 | ; |  |
| 390 | ; |  |
| 400 007F= |  | TEX $=$ \$7F $\quad$ THIS CHARACTER MLUST START |
| 410 | ; | finy message. |
| 420 | ; |  |
| $430 \mathrm{OQFF}=$ |  | ETX $=$ \$FF $\quad$ THIS CHARACTER MUST END |
| 44010 | ; | ANY MESSAGE. |
| 450 | ; |  |
| 4E0 | ; |  |
| 470 | ; |  |
| 480 | ; |  |
| 490 | ; |  |
| 500 | ; |  |
| 510 | , |  |
| 520 | ; |  |
| 530 | ; |  |
| 540 | ; |  |
| 550 | ; |  |
| 5 E 0 | ; | EXTERIAL ADDRESSES |
| 570 | ; |  |


| 580 |  | ; |  |
| :---: | :---: | :---: | :---: |
| 500 |  | ; |  |
| 600 |  | ; |  |
| E10 |  | ; |  |
| 620 |  | ; |  |
| EGO |  | ; |  |
| 640 | $1200=$ |  |  |
| 650 |  | ; | MONITOR CODE. |
| EED | 1205= |  | SELEET $=1$ MPAGE+5 |
| ETa | $1207=$ |  | UISMON=UMPAGE + ? |
| 650 | $1294=$ |  | GET. SL= UMPAGE + \$ 34 |
|  | 12E日 $=$ |  | GETKEY=UMPAGE+\$E® |
| 780 | $1301=$ |  | INC. SL= UMPAGE + \# 10 D |
| 710 | $131 \mathrm{~A}=$ |  | DEC. SL=UMPAGE+\#11A |
| 720 | $132 \mathrm{n}=$ |  | FUT. SL=UMPFGE + \$120 |
| 730 |  | ; |  |
| 740 |  | ; |  |
| 750 | $1400=$ |  | FRPAGE $=1406$ STARTING PAGE OF PRINT |
| 760 |  | ; | UTILITIES. |
| 765 | 1414= |  | PR. ON =PRPAGE + \$14 |
| 767 | 141A= |  | PR. OFF=PRPAGE+\#IA |
| 770 | 144B= |  | PR. CHR=PRPAGE +340 |
| 780 | 14E4= |  | FRINT: =PRPAGE+\$E4 |
| 736 | 1512= |  | PUSHSL=FRPRGE + \$112 |
| 809 | $152 \mathrm{E}=$ |  | POP.SL=PRPAGE+\#12B |
| 810 |  | ; |  |
| 820 |  | ; |  |
| 830 | $1500=$ |  | HEX.PG=\$150 ADDRESS OF PAGE IN WHICH |
| 849 |  | ; | HEXDUMP CODE STARTS. |
| 856 |  | ; |  |
| 850 | $1552=$ |  | SA=HEX.FG+\$52 |
| E76 | 1554= |  | $E R=S A+2$ |
| 838 | 16ET $=$ |  | SAHERE $=$ HEX.PG+\#1ET |
| 836 | 1783= |  | NEXTSL $=$ HEX.PG $+\$ 283$ |
| 50.6 | $17 \mathrm{AD}=$ |  | GOTOSfi $=\mathrm{HEX} . \mathrm{PG}+\$_{\text {F }}$ 2AD |
| Sib |  | ; |  |
| 920 |  | ; |  |
| 930 | 17EQ= |  | MOUERS $=$ \$ 1780 STRRT OF MOUE OBJECT CODE. |
| 940 | $17 \mathrm{BZ}=$ |  | DEST = MOUERS+2 |
| 950 | $17 \mathrm{DE}=$ |  | MOU. EA=MOUERS + \$26 |
| 960 | 18EE= |  | DAHERE=MOUERS+\$132 |
| 970 |  | ; |  |
| 980 | 1EOS= |  | EDFAGE = \#LEQD STARTING PAGE OF EDITOR. |
| 390 | $1 E C O=$ |  | EDKEYS=EDPRGE + SCO |
| 1600 |  | ; |  |
| 1016 |  | ; |  |
| 1020 |  | ; |  |
| 1030 |  | ; |  |
| 1640 |  | ; |  |
| 1050 |  | ; |  |
| 1080 |  | ; |  |
| 1090 |  | ; |  |
| 1080 |  | ; | UARIABLES |
| 1055 |  | ; |  |
| 1100 |  | ; |  |
| 1110 |  | ; |  |
| 1129 |  | ; |  |
| 1130 |  | ; |  |




| 2300 |  | ； |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 2319 | IEET CDCIIE | DO．KEY | CMP | MODEKY | IS IT THE＂CHANGE MODE＂KEY？ |
| 2320 | IEEA DDOE |  | BHE | IFHEXT | IF NOT，PERFORM NEXT TEST． |
| 2330 | 1EEC CEDIIE |  | DEC | EDMODE | If SO，Chinge the editor＇ 5 |
| 2340 | 1EEF 1095 |  | BPL | DO．END | MODE． |
| 2350 | 1EF 1 ASSI |  | LDA | $\# 1$ |  |
| 2360 | 1EF3 EDU11E |  | STA | Edmode | ： |
| 2370 | 1EFE 69 | DO．END | RTS |  | RETIJRN TO CALLER． |
| 2300 |  | ； |  |  |  |
| 2350 |  | ； |  |  |  |
| 2400 | 1EFT CdCzle | IFNEXT | CMP | NEXTKY | IS IT THE＂NEXT＂KEY？ |
| 2410 | IEFA DQQ4 |  | ENE | IFPREU | IF NOT，PERFORM NEXT TEST． |
| 2420 |  | ； |  |  |  |
| 2430 | 1EFC 2日TSiF |  | JSR | NEXTCH | If SO，RDUANCE TO NEXT |
| 2440 |  | ； |  |  | CHARACTER．．． |
| 2450 | 1EFF 60 |  | RTS |  | ．．．AND RETURN． |
| 2460 |  | ； |  |  |  |
| 2470 |  | ； |  |  |  |
| 2489 | 1FOU CDC3IE | IFFRES | CMP | PREUKY | IS IT THE＇FREUIDUS＇KEY？ |
| 2450 | 1 F 03 D 004 |  | ENE | IF．RUB | IF NOT，PERFORM MEXT TEST． |
| 2500 | 1F95 25871F |  | JSR | PREUSL | IF SO，BECK UP TO PREUIOLS |
| 2510 | 1 FQE E |  | RTS |  | CHARACTER AND RETURN． |
| 2525 |  | ； |  |  |  |
| 2530 |  | ； |  |  |  |
| 2543 | $1 F 09$ cocsie | IF．RUE | CMP | RUSKEEY | IS IT THE＊RUBOUT＊KEY？ |
| 2550 | $1 F \mathrm{COC}$［004 |  | ENE | IF．PRT | IF NOT，FERFORM NEXT TEST． |
| 25E0 | 1FGE 20DE1F |  | JSR | DELETE | IF SO，DELETE CURRENT |
| 2570 | $1 F 1160$ |  | RTS |  | CHARACTER AND RETURN． |
| 2580 |  | ； |  |  |  |
| 2589 |  | ； |  |  |  |
| 2E00 | 1F12 CDC4IE | IF．PR，T | CMP | FRTKEY | IS IT THE＂PRINT＂KEY？ |
| 2810 | 1F15 DL04 |  | BHE | IFFLSH | IF NOT，PERFORM NEXT TEST． |
| 2520 | $1 F 17$ 2Đcsif |  | JSR | PRTSUF | IF SO，PRINT THE BUFFER．．． |
| 2535 | IFIF 69 |  | RTS |  | ．AND RETLRN． |
| 2640 |  | ； |  |  |  |
| 2650 |  | ； |  |  |  |
| 2 EEO |  | ； |  |  |  |
| 26.5 | 1FIE CDCDIE | IFFLSH | CMP | FLSHKY | IS IT THE＂FLUSH＂KEY＂？ |
| 2680 | 1F1E［004 |  | BNE | CHARKY | If NOT，IT MUST BE A CHARACTER |
| 2630 |  | ； |  |  | KEY． |
| 2700 | 1F20 20B41F |  | JSR | FLUSH | IF SO，FLUSH THE BUFFER． |
| 2710 | $1 F 23 \mathrm{ED}$ |  | RTS |  | FiNd RETURN． |
| 2720 |  | ； |  |  |  |
| 2730 |  | ； |  |  |  |
| 27.4 |  | ； |  |  |  |
| 2750 |  | ； |  |  |  |
| 2763 |  | ； | OK． | IT＇S NOT Fill | N EDITOR FUNCTION KEY， 50 It |
| 27アロ |  | ；Mus | T BE | a CHARPACTER | R KEY．DEPENDING ON THE |
| 2750 |  | ；CuR | FEAT | MODE，WE＇LL | EITHER INEERT OR OUERSTRIKE |
| 27.90 |  | ； |  |  | THE CURRENT CHARACTER． |
| 2800 |  | ； |  |  |  |
| 2810 | $1 F 24$ AE011E | CHAREX | LDX | EDMODE | ARE WE IN OUERSTRIKE MODE？ |
| 2825 | $1 F 27$ FGU4 |  | EEQ | STRIKE | IF SO，OUERSTRIKE THE CURRENT |
| 2830 |  | ； |  |  | CHARACTER． |
| 28.40 | 1F29 20341F |  | JSR | INYSERT | IF MOT，INSERT THE CHARACTER． |
| 2850 | 1F2C E0 |  | RTS |  | RETURN． |
| 28E0 |  | ； |  |  |  |
| 2870 | 1F20 202013 | STRIKE | JSR | FUT．SL | REFLACE CURRENT CHARACTER |


| 2880 |  | ; |  | WITH NEW CHARACTER. |
| :---: | :---: | :---: | :---: | :---: |
| 2890 | 1730202317 |  | JSR NEXTSL | SELECT NEXT CHARACTER. |
| 2900 | 153360 |  | RTS | RETURN. |
| 2910 |  | ; |  |  |
| 2920 |  | ; |  |  |
| 2930 |  | ; |  |  |
| 2940 |  | ; |  |  |
| 2950 |  | ; |  |  |
| 2950 | $1 F 3448$ | INSERT | PHA | SAUE THE CHARACTER TO BE |
| 2970 |  | ; |  | INSERTED, WHILE WE MAKE ROOM |
| 2980 |  | ; |  | FOR IT IN THE BUFFER... |
| 2990 | $1 F 35201215$ |  | JSR FUSHSL | SFUE THE CURRENT ADDRESS. |
| 3000 | $1 F 38$ AD5315 |  | LDA 5A+1 | SAUE THE EUFFER'S RDDRESS. |
| 3010 | 1F3E 48 |  | PHA |  |
| 3020 | 1F3C An5215 |  | LDA 5A |  |
| 3030 | 1F3F 45 |  | PHA |  |
| 3040 |  | ; |  |  |
| 3050 |  | ; |  |  |
| 3060 | 1F40 AD5515 |  | LDA EA +1 | SAUE BUFFER'S END ADIRESS. |
| 3070 | 1F43 48 |  | PHA |  |
| 3080 | 1F44 AD5415 |  | LDA EA |  |
| 3090 | $1 F 4748$ |  | PHA |  |
| 3100 |  | ; |  |  |
| 3110 |  | ; |  |  |
| 3120 | $1 F 48206716$ |  | JSR SAHERE | SET SA=SELECT, SO CIJRRENT |
| 3130 |  | ; |  | LOCATION WILL BE START OF |
| 3140 |  | ; |  | the block he' ll move. |
| 3150 |  | ; |  |  |
| 3160 |  | ; |  |  |
| 3170 |  | ; |  |  |
| 3180 | 1F4B 208317 |  | JSR NEXTSL | ADUANCE TO NEXT CHARACTER |
| 3190 |  | ; |  | POSITION IN THE BUFFER. |
| 3200 | 1F4E 3011 |  | EMI ENDINS | IF WE' RE AT THE END OF THE |
| 3210 |  | ; |  | BUFFER, WE' LL OUERSTRIKE |
| 3220 |  | ; |  | INSTEAD OF INSERTING. |
| 3230 |  | ; |  |  |
| 3240 |  | ; |  |  |
| 3250 | 1F50 20E218 |  | JSR DRHERE | SET DEST=SELECT. |
| 3260 |  | ; |  | DESTINATION OF BLOCK MOUE |
| 3270 |  | ; |  | WILL EE ONE BYTE RBOUE |
| 3280 |  | ; |  | BLOCK'S INITIAL LOCPTION. |
| 3290 |  | ; |  |  |
| 3300 |  | ; |  |  |
| 3310 | $1 F 53 \mathrm{AD} 5415$ |  | LDA EA | DECREMIENT END ADDRESS |
| 3320 | 1F5S D004 |  | BNE *+5 |  |
| 3330 | $1 F 58$ CE5515 |  | DEC EA+1 |  |
| 3340 | 1F5B CES415 |  | DEC EA |  |
| 3350 |  | ; |  |  |
| 3360 |  | ; |  |  |
| 3376 |  |  |  |  |
| 3380 | 1F5E 200617 | OPENUP | JSR MOU.EA | OPEN UP OME EYTE OF SFACE |
| 3390 |  | ; |  | AT CURRENT CHARACTER' 5 |
| 3400 |  | ; |  | LOCATION, EY MOUING TO DEST |
| 3410 |  | ; |  | THE BLOCK SPECIFIED EY SA, EA. |
| 3420 |  | ; |  |  |
| 3430 |  | ; |  |  |
| 3440 | $1 F 6168$ | ENDINS | PLA | RESTORE EA 50 IT POINTS |
| 3450 | $1 F 62855415$ |  | Sta ef | TO EINE OF BUFFER. |


| 3460 | $1 F 6568$ |  | PLA |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 3470 | 1F66 805515 |  | STA | $E A+1$ |  |
| 3480 |  | ; |  |  |  |
| 3490 |  | ; |  |  |  |
| 3500 | 1FES 58 |  | PLA |  | RESTORE SA SO IT POINTS TO |
| 3510 | IFGA SD5215 |  | STA | SA | START OF BUFFER. |
| 3520 | 1F6D 68 |  | PLA |  | ! |
| 3530 | IFGE 8D5315 |  | STA | $\mathrm{SA}+1$ |  |
| 3540 |  | ; |  |  |  |
| 3550 |  | ; |  |  |  |
| 3560 | 1F71 202615 |  | JSR | POP.SL. | RESTORE SELECT SO IT POINTS |
| 3570 |  | ; |  |  | TO CURRENT CHARRACTER POSITION. |
| 3580 |  | ; |  |  |  |
| 3590 |  | ; |  |  |  |
| 3500 | 1F74 68 |  | PLA |  | RESTORE NEW CHARACTER TO |
| 3510 |  | ; |  |  | ACCUMULATOR. WE VE CREATED |
| 3620 |  | ; |  |  | A ONE-BYTE SPACE FOR IT, SO |
| 3630 | 1F75 20201F |  | JSR | STRIKE | WE NEED ONLY OUERSTRIKE IT |
| 3640 | $1 F 7860$ |  | RTS |  | AND RETURN. |
| 3650 | 1579209412 | NEXTCH | JSR | GET.SL | GET CURRENT CHARACTER. |
| 36E0 | 1F7C C9FF |  | CMP | \#ETX | IS IT END OF TEXT CHARACTER? |
| 3670 | 1F7E FOQ4 |  | BEQ | AN.ETX | IF SO, RETURN TO CALLER, |
| 3580 |  | ; |  |  | berring a negative return code. |
| 3690 |  | ; |  |  |  |
| 3700 | 1580208317 |  | JSR | NEXTSL | If NOT, SELECT NEXT BYTE IN |
| 3710 |  | ; |  |  | BLIFFER. |
| 3720 | 178360 |  | RTS |  | RETURN PLUS IF WE INCREMENTED |
| 3730 |  | ; |  |  | SELECT; MINUS IF SELECT |
| 3740 |  | ; |  |  | RLREADY EQUALLED EA. |
| 3750 |  | ; |  |  |  |
| 3760 | $1 F 84$ fgrf | AN.ETX | LDA | \#\#FF | SINCE WE' RE ON RN ETX, WE |
| 3770 | 1F8G 6 |  | RTS |  | WILL RETURN MINUS, WITHOUT |
| 3780 |  | ; |  |  | INCREMENTING SELECT. |
| 3790 |  | ; |  |  |  |
| 3300 |  | ; |  |  |  |
| 3810 |  | ; |  |  |  |
| 3820 |  | ; |  |  |  |
| 3830 | 158738 | PREUSL | SEC |  | PREPARE TO COMPARE. |
| 3840 | $1 F 68$ RD5315 |  | LDA | SR+1 | IS SELECT IN A HIGHER PAGE |
| 3850 | 1FBB CDOEI2 |  | CMP | SELECT+1 | THAN START OF BUFFER? |
| 3860 | IFSE GIDC |  | ECC | SL.OK | If SO, SELECT MRY EE DECREMENTED |
| 3870 | 1 FGO DOLO |  | ENE | NOT.OK | IF SELECT IS IN A LOWER |
| 3880 |  | ; |  |  | PAGE THAN SA, IT'S NOT OK. |
| 3896 |  | ; |  |  |  |
| 3960 |  | ; |  |  | SELECT IS IN SRME PAGE AS SA. |
| 3910 | 1F92 A 55215 |  | LDA | SA | IS SELECT $>5$ A? |
| 3920 | 1595 CD0512 |  | CMP | SELECT |  |
| 3330 | 1 FGE F017 |  | BEQ | NO. DEC | IF SELECT=SA, DON' T DECREMENT |
| 3940 |  | ; |  |  | SELECT. |
| 3950 | IFGA BDES |  | BCS | NOT. OK | IF SELECT<SA, DON T DECREMENT |
| 3960 |  | ; |  |  | SELECT. |
| 3570 | 1FSC 201A13 | SL.OK | J5R | DEC. SL | SELECT $>5 \mathrm{SA}$, 50 WE MAPY |
| 3980 |  | ; |  |  | DECREMENT SELECT AND IT |
| 3990 |  | , |  |  | WILL REMAIN IN THE EUFFER. |
| 4000 | 1F9F ASOO |  | LDA | \# | SET A POSITIUE RETURN CODE... |
| 4010 | 1FR1 50 |  | RTS |  | ...RND RETURN. |
| 4020 |  | ; |  |  |  |
| 4030 |  | ; |  |  |  |


| 4040 | $1 F A 2$ | AD5215 | NUT.OK | LDA | 5A | SINCE SELECT<SA, IT IS NOT |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 4050 | 1FAS | 600512 |  | STA | SELECT | EUEN IN THE EDIT EUFFER. SO |
| 4060 | 1FAB | AD5315 |  | LDA | Sf+1 | MAKE SELECT LEGAL, by Setting |
| 4070 | $1 F A E$ | 800512 |  | STA | SELECT+1 | IT EQUAL TO SA. |
| 4080 | 1FRE | F300 |  | LDA | \# | SET A FOSITIUE RETURN CODE. |
| 4090 | 1FED | 60 |  | RTS |  | . . And return. |
| 4100 |  |  | ; |  |  |  |
| 4110 |  |  | ; |  |  |  |
| 4120 | 1FBI | RSFF | NO.DEC | LDA | \#\$FF | SELECT=SA, SO CHANGE |
| 4130 | LFE3 | E日 |  | RTS |  | NOTHING. RETURN WITH |
| 4140 |  |  | ; |  |  | NEGATIUE RTURN CODE. |
| 4150 |  |  | ; |  |  |  |
| 4160 |  |  | ; |  |  |  |
| 4170 |  |  | ; |  |  |  |
| 4180 | 1FE4 | 20f017 | FLUSH | JSR | gotosa | SET SELECT=5\%. |
| 4190 | 1FB7 | AGFF | FLOOP | LDA | \#ETX | PUT AN ETX CHARACTER |
| 4200 | 1FES | 202013 |  | JSR | PUT.SL | INTO THE BUFFER. |
| 4210 | IFBC | 208317 |  | JSR | MEXTSL | ADUANCE TO NEXT POSITION IN |
| 4220 |  |  | ; |  |  | BUFFER. |
| 4230 | 1FBF | 10F6 |  | BPL | FLOOP | If We haven' $T$ reached end |
| 4240 |  |  | ; |  |  | OF BUFFER, PUT RH ETX INTO |
| 4250 |  |  | ; |  |  | THIS POSITION, TOO. |
| 4260 |  |  | ; |  |  |  |
| 4270 | $1 F C 1$ | 2 ARO 17 |  | JSR | gotosa | HAUING FILLED BUFFER WITH |
| 42 Ca |  |  | ; |  |  | ETC CHARACTERS, RESET SELECT |
| 4230 |  |  | ; |  |  | TO BEGINNING OF BUFFER. |
| 4300 | 1FCA | 60 |  | RTS |  | RETURN. |
| 4310 | 1FCS | 20AD17 | FRTEUF | JSR | GOTOSA | SET SELECT TO START OF BUFFER |
| 4320 | ifCE | 201414 |  | JSR | PR.ON | SELECT PRINTER FOR OUTPUT. |
| 4330 | $1 F C B$ | 203412 | FRLOOP | JSR | GET. SL | GET CURRENT CHARACTER. |
| 4340 | IFCE | CSFF |  | CMP | \#ETX | IS IT ETX? |
| 4350 | $1 F D 0$ | F0id |  | BEO | ENDPRT | IF SO, WE' RE DONE. |
| 4360 | 1FDC | 204014 |  | JSR | PR.CHR | IF NOT, PRINT IT. |
| 437 D | $1 F D 5$ | 206317 |  | JSR | NEXTSL | SELECT NEXT CHARACTER |
| 4300 | IFDS | 10F1 |  | EFL | PRLOOF | IF WE HAUEN' T REACHED THE |
| 4390 |  |  | ; |  |  | END OF THE BuFfer, HAMdLE |
| 4400 |  |  | ; |  |  | THE CURRENT CHARACTER RS BEFORE. |
| 4410 | 1FDA | 4C1A14 | ENDFRT | JMP | PR.OFF | HAUING REACHED END OF MESSAGE |
| 4420 |  |  | ; |  |  | OR END OF BUFFER, RETURN TO |
| 4430 |  |  | ; |  |  | CALLER OF EDITIT, DESELECTING |
| 4440 |  |  | ; |  |  | THE PRIHTER AS WE DO 50. |
| 4450 |  |  | ; |  |  |  |
| 44E0 |  |  | ; |  |  |  |
| 4475 | 1FnD | 201215 | DELETE | JSR | PUSHSL | SAUE CURRENT ADDRESS. |
| 4480 | 1FED | AD5315 |  | LDA | SA+1 | SRUE BUFFER' S START ADDRESS. |
| 4450 | 1FE3 | 48 |  | PHA |  |  |
| 4500 | 1FE4 | AD5215 |  | LDA | SA |  |
| 4510 | IFET | 48 |  | FHA |  |  |
| 4520 |  |  | ; |  |  |  |
| 4530 | $1 F E B$ | 20E218 |  | JSR | DAHERE | SET DEST=SELECT, BECAUSE |
| 4540 |  |  | ; |  |  | WE' LL MOUE A ELOCK OF TEXT |
| 4550 |  |  | ; |  |  | DOWIN TO HERE, TO CLOSE UP |
| 4560 |  |  | ; |  |  | THE EUFFER AT THE CURRENT |
| 4570 |  |  | ; |  |  | CHARACTER. |
| 4580 | 1FEE | 208317 |  | JSR | MEXTSL | ADUANCE BY ONE BYTE THROUGH |
| 4590 |  |  | ; |  |  | BUFFER, IF POSSIBLE. |
| 4800 | 1FEE | 206716 |  | JSR | SRHEFE | SET SA=SELECT, EECGUSE THIS |
| 4610 |  |  | ; |  |  | IS THE START OF THE BLOCK WE' LL |


| 4820 |  | ; |  | MOVE DOWIN. |
| :---: | :---: | :---: | :---: | :---: |
| 4630 |  | ; |  | NOTE: THE ENDING ADDRESS OF |
| 4546 |  | ; |  | THE BLOCK IS THE END ADDRESS |
| 4850 |  | ; |  | OF THE TEXT BUFFER. |
| 4EET | 1FF1 20D617 |  | JSR MOU.EA | MOUE BLOCK SPECIFIED BY |
| 467 C |  | ; |  | SA, EA TO DEST. |
| 4 ESO |  | ; |  |  |
| 4690 |  | ; |  |  |
| 4700 | 1FF4 68 |  | PLA | RESTORE INITIAL SA (WHICH |
| 4710 | 1FFE EDS215 |  | STA SA | IS THE START ADDRESS OF THE |
| 4720 | 1 FFE 69 |  | PLA | TEXT BUFFER, NOT OF THE BLOCK |
| 4730 | 1FFF EL5315 |  | STA SA+1 | WE JUST MOUED.) |
| 4740 | 1FFC 202315 |  | JSR POP. SL | RESTIORE CURRENT ADDRESS. |
| 4750 | 1 FFF 50 |  | RTS | RETURM TO CALLER. |

## Appendix Cl 2 :

## Extending the Visible Monitor




| 1160 | 1RES DODD |  | BNE | IF.T | IF NOT, PERFORM NEXT TEST. |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1170 | 10E7 ADOE14 |  | LDA | PRINTR | IS THE PRINTER SELECTED? |
| 1180 | 16EA DOD4 |  | ENE | NEXT. 2 | If SO, PRINT A DISASSEMBLY. |
| 1190 | 10EC 200919 |  | JSR | TU.DIS | IF. NOT, DISASSEMBLE TO THE |
| 1200 | 10EF S0 |  | RTS |  | SCREEN AND RETURN. |
| 1210 | 10Fb 202619 | NEXT. 2 | JSR | PR.DIS | FRINT A DISASSEMBLY... |
| 1220 | 10F3 60 |  | RTS |  | GIND RETURN. |
| 1230 |  | ; |  |  |  |
| 1240 | 10F4 CS54 | IF.T | CMP | ${ }^{\text {\# }}$ T | IS IT THE ' $T$ ' KEY? |
| 1250 | $10 F 6$ DgO4 |  | BNE | EXIT | IF NOT, RETURN. |
| 1260 | 10F8 20621E |  | JSR | EDITOR | If SO, CALL THE SIMPLE |
| 1270 | 10FE 60 |  | RTS |  | TEXT EDITOR RND RETURN. |
| 1280 |  | ; |  |  |  |
| 1250 | 16FC 60 | EXIT | RTS |  | EXTEND THE UISIBLE MONITOR |
| 1300 |  | ; |  |  | EUEN FLSTHER EY REPLAGING |
| 1310 |  | ; |  |  | THIS 'RTS' WITH A 'JMP' TO |
| 1320 |  | ; |  |  | MORE TEST-AND-BRANCH CODE. |

## Appendix CI3:

System Data Block for the Ohio Scientific C-IP


| 580100503 | HIFAGE | . BYTE \$D3 | HIGHEST PRGE IN SCREEN MEMORY. |
| :---: | :---: | :---: | :---: |
| 590100620 | BLANK | - BYTE \$2, | OSI DISFLAY CODE FOR A BLANK. |
| 600 100710 | ARROW | . BYTE \$10 | OSI DISPLAY COLE FOR AN UP-ARROW |
| E10 | ; |  |  |
| 620 | ; |  |  |
| 630 | ; |  |  |
| 640 | ; |  |  |
| 650 | ; |  |  |
| 660 | ; |  |  |
| 670 | ; |  |  |
| 680 | ; |  |  |
| 690 | ; |  |  |
| 700 | ; ***** |  | ************************** |
| 710 | ; |  |  |
| 720 | ; | INPUT | STPUT UECTORS |
| 730 | ; |  |  |
| 740 | ; ***** |  |  |
| 750 | ; |  |  |
| 750 | ; |  |  |
| 770 | ; |  |  |
| 780 | ; |  |  |
| 790 | ; |  |  |
| 800 | ; |  |  |
| 8101608 EDFE | ROMKEY | . WORD SFEED | POINTER TO ROUTINE THAT GETS PN ASCIT CHARACTER FROM THE |
| 820 | ; |  | AM ASCII CHARACTER FROM THE |
| 830 | ; |  | KEYBOARD. (NOTE: PFFEB IS |
| 840 | ; |  | THE GENERAL CHARACTER-INPUT |
| 850 | ; |  | ROUTINE FOR OSI BRSIC-IN-ROM |
| 860 | ; |  | COMPUTERS. 3 |
| 870 | ; |  |  |
| 880 | ; |  |  |
| 850 100A 2DBF | ROMTUT | . WORD \$BF2D | POINTER TO ROUTINE TO PRINT AN ASCII CHARACTER ON THE SCREEN |
| 900 | ; |  | AN ASCII CHARACTER ON THE SCREEN |
| 910 | ; |  | (NOTE: \$FFEE IS THE |
| 920 | ; |  | CHARRACTER-OUTPUT ROUTINE FOR |
| 930 | ; |  | OSI BASIC-IN-ROM COMPUTERS.) |
| 940 | ; |  |  |
| 950 | ; |  |  |
| $960100 C$ B1FC | ROMPRT | . WORD \$FCBI | POINTER TO ROUTINE TO SEND AN ASCII CHARACTER TO THE PRINTER |
| 970 | ; |  | (ACTUALLY, TO THE CRSSETTE.PORT. |
| 980 | ; |  |  |
| 590 | ; |  |  |
| 1000 | ; |  |  |
| 1010100 E 1010 | USROUT | . WORD DUMMY | POINTER TO USER-WRITTEN OUTPUT |
| 1020 | ; |  | ROUTINE. (SET HERE TO DUMMY |
| 1030 | ; |  | UNTIL YOU SET IT TO POINT |
| 1040 | ; |  | TO YOUR OWN CHARACTER-OUTPUT |
| 10.50 | ; |  | ROUTINE.) |
| 1060 | ; |  |  |
| 1070 | ; |  |  |
| 10ee 1010 50 | DUMMY | RTS | THIS IS A DUMMY SUBROUTINE. |
| 1090 | ; |  | IT DOES NOTHING BUT RETURN. |
| 1100 | ; |  |  |
| 1110 | ; |  |  |
| 1120 | ; |  |  |
| 1130 | ; |  |  |
| 1140 | ; |  |  |
| 1150 | ; |  |  |



## Appendix CI4:

System Data Block for the PET 2001



| 1190 |  | ； |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1200 |  | ； |  |  |  |
| 1210 |  | ； |  |  |  |
| 1220 | 1011 297F | FIXCHR | FND | 为 7 7F | CLEAR BIT 7，TO MAKE IT |
| 1230 |  | ； |  |  | A LEGAL ASCII CHARACTER． |
| 1240 | 101338 |  | SEC |  | PREPARE TO COMPARE． |
| 1250 | 1014 C940 |  | CMP | \＃\＄40 | IS IT LESS THAN \＄40？（IS |
| 1280 |  | ； |  |  | IT A NUMBER OR PUPCTUATION |
| 1270 |  | ； |  |  | MARK， |
| 1280 | 10169811 |  | BCC | FIXEMD | IF SO，NO CONUERSION NEEDED． |
| 1290 |  | ； |  |  |  |
| 1300 | 1918 c960 |  | CMP | 并あEワ | IS IT BETWEEN \＄40 AND \＄60？ |
| 1310 |  | ； |  |  |  |
| 1320 | 101a 900a |  | BCC | SUB． 40 | IF＇ 50, SUBTRACT $\$ 40$ TO |
| 1330 |  | ； |  |  | CONUERT FROM．ASCII TO PET． |
| 1340 |  | ； |  |  |  |
| 1350 |  | ； |  |  | IT＇S＞＝\＄50，SO WE MUST |
| 1370 | 101C A2BE |  | LDX | \＃14 | SET PET DISPLAY MODE FOR |
| 1350 | 101E ED4CES |  | STA | 59468 | CHARACTER SET THAT INCLUDES |
| 1390 |  | ； |  |  | LOWER CASE ALPHA CHARACTERS． |
| 1400 | 1021 E920 |  | SBC | \＃\＃20 | SUBTRFACT \＄20 TO CONUERT |
| 1410 |  | ； |  |  | LOWER CASE ASCII TO PET CODE． |
| 1420 | 102318 |  | CLC |  |  |
| 1430 | 10249093 |  | BCC | FIXEND |  |
| 1435 |  | ； |  |  |  |
| 1440 | 162638 | SUB． 40 | SEC |  | PREPARE TO SUBTRACT． |
| 1450 | 1027 E940 |  | SBC | \＃\＄4］ | SUBTRACT \＄40 TO．CONUERT ASCII |
| 1460 |  | ； |  |  | UPPER CASE CHAR TO FET CODE． |
| 1470 | 1629 E0 | FIXEND | RTS |  | RETURN，WITH A HOLDING |
| 1480 |  | ； |  |  | PET DISPLAY CODE FOR ASCII |
| 1499 |  | ； |  |  | ORIGINALLY IN A． |
| 1506 |  | ； |  |  |  |
| 1510 |  | ； |  |  |  |
| 1520 |  | ； |  |  |  |
| 1530 |  | ； |  |  |  |
| 1540 |  | ； |  |  |  |
| 1550 |  | ；＊＊＊＊＊ | ＊＊＊＊＊ |  |  |
| 1563 |  | ； |  |  |  |
| 1576 |  | ； | GET A | AN RSCII | CHARACTER FROM THE KEYBOARD |
| 1580 |  | ； |  |  |  |
| 1590 |  | ；＊＊＊ | ＊＊＊＊＊＊＊ | ＊考＊＊＊＊＊＊＊＊＊＊ |  |
| 1506 |  | ； |  |  |  |
| 1610 |  | ； |  |  |  |
| 1620 |  | ； |  |  |  |
| 1630 |  | ； |  |  |  |
| 1640 | 102A 20E4FF | PETKEY | JĐR | \＄FFE4 | SCAN THE PET KEYBOARD |
| 1650 | 10202975 |  | AMD | \＃${ }^{\text {d }}$ 7F | CLEAR BIT 7，TO BE SURE |
| 1660 |  | ； |  |  | IT＇ 5 A LEGAL ASCII CHARACTER． |
| 1670 | 102F FGFs |  | BEO | PETKEY | ZERO MEANS NO KEY，SO |
| 1686 |  | ； |  |  | SCAN AGAIN． |
| 1 ESO |  | ； |  |  |  |
| 1700 | 1031 E0 |  | RTS |  | RETURN WITH RSCII CHARACTER |
| 1716 |  | ； |  |  | FROM THE KEYBOARD． |

## Appendix CI5:

System Data Block for the Apple II


| 500 |  | ； |  | NORMAL DISPLAY MODE（WHITE |
| :---: | :---: | :---: | :---: | :---: |
| 590 |  | ； |  | CHARACTERS ON A DARK |
| E00 |  | ； |  | BACKGROUND． 3 |
| 610 | 1007 DE | GRRON | ．BYTE \＄DE | APPLE II DISFLAY CODE FOR |
| 629 |  | ； |  | a CARAT（USED dechuse mfple |
| E30 |  | ； |  | II HAS NO UP－ARROW．） |
| 640 |  | ； |  |  |
| E50 |  | ； |  |  |
| 650 |  | ； |  |  |
| 679 |  | ； |  |  |
| 689 |  | ； |  |  |
| 690 |  | ； |  |  |
| 700 |  | ； |  |  |
| 710 |  | ； |  |  |
| 720 |  | ； |  |  |
| 730 |  | ；＊＊＊＊＊ |  | 考＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊ |
| 740 |  | ； |  |  |
| 750 |  | ； | INPUT | ITPUT VECTORS |
| 760 |  | ； |  |  |
| 770 |  | ；＊＊＊＊ | 漓＊＊＊＊＊＊＊＊＊＊＊＊＊ |  |
| 750 |  | ； |  |  |
| 790 |  | ； |  |  |
| 850 |  | ； |  |  |
| 810 |  | ； |  |  |
| 820 |  | ； |  |  |
| 830 |  | ； |  |  |
| 840 | 10081410 | ROMKEY | ．WORD PPLKEY | POINTER TO ROUTINE THAT GETS |
| 850 |  | ； |  | AN ASCII CHARACTER FROM THE |
| 860 |  | ； |  | KEYBOARD．（NOTE：APLKEY |
| 870 |  | ； |  | CALLS A ROM SUbroutine，But |
| 889 |  | ； |  | APLKEY IS NOT AN APPLE ROM |
| 890 |  | ； |  | SUBROUTINE． 3 |
| 900 |  | ； |  |  |
| 910 |  | ； |  |  |
| 920 | 10日A 1alo | ROMTUT | ．WORD APLTUT | POINTER TO ROUTIINE TO FRINT |
| 939 |  | ； |  | AN ASCII CHARACTER ON THE SCREEN |
| 940 |  | ； |  |  |
| 950 |  | ； |  |  |
| 960 | 10001010 | ROMPRT | ．WORD DUMMY | POINTER TO ROUTINE TO SEND AN |
| 970 |  | ； |  | ASCII CHARACTER TO THE PRINTER |
| 980 |  | ； |  | （SET TO DUMMY UNTIL YOU MAKE |
| 999 |  | ； |  | IT POINT TO THE CHARACTER－ |
| 1000 |  | ； |  | OUTPIUT ROUTINE THAT DRIVES |
| 1010 |  | ； |  | YOUR PRINTER． 3 |
| 1020 |  | ； |  | YOU MAY WISH TO |
| 1030 |  | ； |  | SET ROMPRT 50 IT FOINTS TO |
| 1040 |  | ； |  | \＄FDED，THE GPFLE II＇S |
| 1050 |  | ； |  | GEMERAL CHARACTER OUTPUT |
| 10 E |  | ； |  | ROUTINE，SFDED WILL PRINT TO |
| 1070 |  | ； |  | A PRINTER IF YOU TELL |
| 1080 |  | ； |  | YOUR APFLE II ROM SOFTWARE |
| 1090 |  | ； |  | TO SELECT YOUR PRINTER RS |
| 1100 |  | ； |  | AN OUTPUT DEUICE．DO THAT |
| 1110 |  | ； |  | IN EASIC BY TYPING＊PR 期＂， |
| 1120 |  | ； |  | WHERE $N$ IS THE NUMEER OF THE |
| 1130 |  | ； |  | SLOT HOLDING THE CIRCUIT CARD |
| 1140 |  | ； |  | THAT DRIUES YOUR PRINTER． |



| 1730 | ; |  |  |
| :---: | :---: | :---: | :---: |
| 1740 |  |  |  |
| 1750 | ; |  |  |
| 1760 | PRINT AN ASCII CHARACTER ON THE SCREEN |  |  |
| 1770 | ( |  |  |
| 1780 |  |  |  |
| 1790 | ; |  |  |
| 1800 | ; |  |  |
| 1810 | ; |  |  |
| 1820 | ; |  |  |
| 1830 | ; |  |  |
| 164010180980 | APLTUT ORA \#\#BG |  | SET BIT 750 CHARACTER WILL PRINT IN NORMAL MODE. |
| 1850 |  |  |  |
| 1860 101C 20FDFB |  | JSR SFBFD | CALL APPLE II ROM ROUTINE TO |
| 1870 | ; |  | PRINT A CHARFGCTER TO SCREEN. |
| $1880101 F 60$ |  | RTS | RETURN TO CPALLER. |

## Appendix Cl 6 :

## System Data Block for the Atari 800

| 19 |  | ; | APPEINIX CIE: ASSEMBLER LISTING OF |
| :---: | :---: | :---: | :---: |
| 20 |  | ; | SYSTEM DATA ELOCK |
| 30 |  | ; | FOR THE ATARI 80] |
| 40 |  | ; |  |
| 53 |  | ; |  |
| E0 |  | ; | , |
| 7 B |  | ; |  |
| 80 |  | ; | SEE RPFEMDIX B4 OF BEYOND GAMES: SYSTEM |
| 90 |  | ; | SOFTWRE FOR YOUR G502 PERSOMAL COMPUTER |
| 16.4 |  | ; |  |
| 116 |  | ; |  |
| 120 |  | ; | BY KEN SKIER |
| 130 |  | ; |  |
| 148 |  | ; |  |
| 150 |  | ; |  |
| 168 |  | ; |  |
| 178 |  | ; |  |
| 180 |  | ; |  |
| 190 |  | ; |  |
| 2010 |  | ; |  |
| 216 |  | ; |  |
| 220 |  | ; |  |
| 239 |  | ; |  |
| 240 |  | ; |  |
| 250 |  | ; |  |
| 26! |  | ; |  |
| 270 |  | , |  |
| 280 |  | ; | EXTEFNAL ADDRESSES |
| 259 |  | ; |  |
| 3015 |  | ; |  |
| 310 |  | ; |  |
| 320 |  | ; |  |
| 330 |  | ; |  |
| 340 |  | ; |  |
| 359 |  | ; |  |
| 350 |  | ; |  |
| 370 | 900] $=$ |  | TU. $P T R=0$ |
| 3601 |  | ; |  |
| 393 | 1120] |  | TUSUES $=$ 等1100 |
| 400 | $1113=$ |  | CLR. $X$ Y = TUSUBS + \$13 |
| 410 | 1125= |  | TUHOME=TUSUBS + H2B |
| 4 C | $113 C=$ |  | TUTOXY=TUSUBS+53C |
| 430 | 1176= |  | TUDOWM=TUSUBS+47E |
| 440 | $1104=$ |  | TUPUSH=TUSUBS+\#C4 |
| 450 | $1183=$ |  | TU. FOP=TUSUES + \$D3 |
| 400 | $1170=$ |  | UUCHAR=TUSUBS + WTC |
| 479 |  | ; |  |
| 480 | $1500=$ |  | HEX. PG= 1500 |
| 490 | $1552=$ |  | SA=HEX. PG + \#5 52 |
| 500 | $1554=$ |  | $E A=5 A \div 2$ |
| 510 |  | ; |  |
| 523 | $1700=$ |  | MOU. $\mathrm{FG}=\$ 1700$ |
| 533 | $17 \mathrm{BL}=$ |  | DEST=MOU.FG+ \#B2 |
| 540 | 17DS= |  | MOU. $E$ A $=$ MOU. $P G+\Phi D G$ |
| 550 |  | ; |  |
| 569 |  | ; |  |
| 579 |  | ; |  |




| 1740 |  | ； |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 1750 |  | ； |  |  |
| 1760 |  | ； |  |  |
| 1770 |  | ； |  |  |
| 1780 | 1811 297F | FIXCHR | AND 4 \＃ 7 F | CLEFIR BIT 7 SO CHARACTER IS |
| 1790 |  | ； |  | A LEGITIMATE ASCII CHARACTER． |
| 1800 | 101338 |  | SEC | PREPARE TO COMPARE． |
| 1810 | 1014 C920 |  | CMP \＃\＄20 | IS CHARACTER＜\＄20？ |
| 1820 | 1016 9008 |  | BCC BADCHR | If SO，IT＇ 5 NOT A UIEMAELE |
| 1830 |  | ； |  | ASCII CHARACTER， 50 RETURN |
| 1848 |  | ； |  | A BLANK． |
| 1850 |  | ； |  |  |
| 18 ED | 1018 Csso |  | CMP \＃\＄50 | IS CHRRACTER＜\＄60？ |
| 1870 | 101A 9006 |  | ECC SUB．20 | IF SO，SUETRACT £2g AND RETURN． |
| 18 CO | 101C c97b |  | CMP 部7B | CHARACTER＜\＄7B？ |
| 1890 | 101E 9007 |  | BCC FIXEND | IF SO，NO CONUERSION IS NEEDED． |
| 15E0 |  | ； |  |  |
| 1510 | 1020 AD0610 | BADÇHR | LDA BLANK | THE CHARACTER IS NOT A |
| 1920 |  | ； |  | UIENAELE ASCII CHARACTER， |
| 1930 | 1023 E0 |  | RTS | SO RETURN A BLANK． |
| 1940 | 102438 | 5UB． 20 | SEC | FREPARE TO SUBTRACT． |
| 1950 | 1025 E920 |  | SEC 来弗20 | SUBTRACT 520 TO COMUERT ASCII |
| 1950 |  | ； |  | TO ATARI DISFLAY CODE． |
| 1970 | 102760 | FIXEND | RTS | RETURY WITH ATARI DISPLAY |
| 1980 |  | ； |  | COEE FOR ORIGINAL ASCII |
| 1590 |  | ； |  | CHARACTER． |
| 2500 |  | ； |  |  |
| 2010 |  | ； |  |  |
| 2020 |  | ； |  |  |
| 2030 |  | ； |  |  |
| 2040 |  | ； |  |  |
| 2050 |  | ； |  |  |
| 20E0 |  | ； |  |  |
| 2070 |  | ； |  |  |
| 2089 |  | ；＊＊＊＊ |  |  |
| 2090 |  | ； |  |  |
| 2100 |  |  | GET AN ASCII | CHARACTER FROM THE KEYBOARD |
| 2110 |  | ； |  |  |
| 2120 |  | ；＊＊＊＊ |  |  |
| 2130 |  | ； |  |  |
| 2140 |  | ； |  |  |
| 2150 |  | ； |  |  |
| 2150 |  | ； |  |  |
| 2170 |  | ； |  |  |
| 2180 |  | ； |  |  |
| 2150 |  | ； |  |  |
| 2200 |  | ； |  |  |
| 2210 | 1028 ADFC02 | ATRKEY | LDf \＄D2FC | HAS A KEY BEEN IEPRESSED？ |
| 2220 | 102B CSFF |  | CMP \＃\＃FF | \＄FF MEANS NO KEY． |
| 22301 | 1020 F0F9 |  | EEO ATRKEY | IF NOT，LOOK AGAIN． |
| 2240 |  | ； |  |  |
| 2250 |  | ； |  | A KEY HAS GONE DOWN． |
| 2260 |  | ； |  | ACCUMULRTOR HOLDS ITS |
| 2275 |  | ； |  | HARDMARE KEY－CODE． |
| 2280 | 102 F A8 |  | TAY | PREFARE TO USE THAT CODE AS |
| 2290 |  | ； |  | AS AN INDEX． |
| 2306 |  | ； |  |  |
| 2310 |  | ； |  |  |


| 2350 | 1050 | EOU6GF |  | LIA | ATEKYS，Y |  | LOOK UP CHARACTER FOR THAT |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2330 |  |  | ； |  |  |  | KEY AND SHIFT STATE． |
| 2346 | 1633 | 50 |  | RTS |  |  | RETUJRN WITH ASCII CHARACTER |
| 2300 |  |  | ； |  |  |  | FOR THAT KEY AND SHIFT STATE． |
| 25eb |  |  | ； |  |  |  |  |
| 23\％0 |  |  | ； |  |  |  |  |
| 23¢ |  |  | ： |  |  |  | ．： |
| 2330 |  |  | ； |  |  |  |  |
| 2490 |  |  | ； |  |  |  |  |
| 2410 |  |  | ； |  |  |  |  |
| 2429 |  |  | ； |  |  |  |  |
| 24303 |  |  | ；＊＊＊ | 为为为家 | ＊＊＊＊＊＊＊＊＊＊＊ | ＊＊＊ | ＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊ |
| 2449 |  |  | ； |  |  |  |  |
| 2450 |  |  |  | PRI | HT AN ASC | II | CHARACTER ON THE SCREEN |
| 2－160 |  |  | ； |  |  |  |  |
| 2476 |  |  | ＊＊＊＊ | ＊＊＊＊ |  | ＊＊＊ | ＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊＊ |
| 2480 |  |  | ； |  |  |  |  |
| 2490 |  |  | ； |  |  |  |  |
| 2 CEO |  |  | ； |  |  |  |  |
| 2510 | B00n＝ |  |  | $\mathrm{CR}=$ | \＄00 |  | ASCII CARPIAGE RETURN． |
| 2523 | 006A＝ |  |  | LF： | \＄0A |  | ASCII LIMEFEED CHRRACTER． |
| 2530 |  |  | ； |  |  |  |  |
| 2540 |  |  | ； |  |  |  |  |
| 2550 |  |  | ； |  |  |  |  |
| 2560 | 1834 | 05 | TUCHAR | ． BY | TE 0 |  | THIS BYTE HOLDS CHARACTER |
| 2570 |  |  | ； |  |  |  | TO BE DISPLAYED．（ALSO， |
| 2580 |  |  | ； |  |  |  | CHARACTER MOST RECENTLY |
| 2550 |  |  | ； |  |  |  | DISPLAYED，USING TUTSIM．） |
| 2EQC | 1035 | 븐 | TU．COL | ．BY | TE D |  | THIS BYTE HOLDS COLUIN IN |
| 2610 |  |  | ； |  |  |  | WHICH CHARACTER WILL NEXT |
| 2620 |  |  | ； |  |  |  | APPPEAR．WE MAY THINK OF IT |
| 2630 |  |  | ； |  |  |  | AS THE POSITION OF AN |
| 2540 |  |  | ； |  |  |  | ELECTRONIC＂PRINT－HEAD＂． |
| 2E50 |  |  | ； |  |  |  |  |
| 26EO |  |  | ； |  |  |  |  |
| 2670 |  |  | ； |  |  |  |  |
| 2685 | 1035 | Cs20 | TUTSIM | CIIP | \＃CR |  | IS CHARACTER AN ASCII |
| 2696 |  |  | ； |  |  |  | CARRIAGE RETURN？ |
| 2769 | 1038 | D005 |  | BNE | LFTEST |  | IF NOT，PERFORM NEXT TEST． |
| 2710 | 1639 | ficel | RESET | LDA |  |  | RESET TU COLUMN TO |
| 2720 | $103 C$ | 803510 |  | STA | TU．COL |  | LEFT MARGIN AND |
| 27.30 | 163F | EO |  | RTS |  |  | RETURN． |
| 2748 |  |  | ； |  |  |  |  |
| 2750 | 1040 | C90a | LFTEST | CMP | \＃LF |  | Is IT A LINEFEED CHARACTER？ |
| 2760 | 1042 | D0e3 |  | BNE | CHSRVE | IF | NOT，HANDLE IT AS A CHARACTER |
| 2770 | 10.44 | 4CBGOE |  | JMP | SCROLL |  | SCROLL TEXT UP FOR A LINEFEED． |
| 27E日 |  |  | ； |  |  |  |  |
| 2790 |  |  | ； |  |  |  | SINCE IT＇S NOT CR OR LF， |
| 2860 | 1047 | 803410 | CHSAVE | STA | TUCHAR |  | LET＇S SAUE IT． |
| 2315 | 164A | 20C411 |  | JSR | TUPUSH |  | SAVE ZERO PRGE BYTES WE＇LL USE． |
| 2820 |  |  | ； |  |  |  |  |
| 2830 | 104D | ACE410 |  | LDY | TUROWS |  | SET TU．PTR TO CURRENT |
| 2645 | 1050 | AESE1G |  | LDX | TU．COL |  | FOSITION OF＂PRINT－HEAD＂． |
| 2850 | 1053 | 203C11 |  | JSR | tutoxy |  |  |
| 28E0 |  |  | ； |  |  |  |  |
| 2870 | 1056 | AD3410 |  | LDA | TUCHAR |  | GET CHARACTER TO BE DISPLAYED． |
| 2880 | 1059 | 207611 |  | JSR | UUCHAR |  | SHOW IT． |
| 2890 | 105C | EE3510 |  | INTC | TU．COk |  | fduance＂print－head＂to mext |




LDA TU.PTR+1 MEMORY TOWARDS THE HOME STA UEST+1 ADDRESS.

JSR TUDOWN SET SA=ADDRESS OF SCREEN LDA TU.PTR POSITION AT COLUMN D, ROW 1. STA SA THAT MARKS THE START OF LDA TU.PTR+1 OF THE BLOCK TO BE MOUED. STA SA+1

LDX TUCOLS SET EA=ADDRESS OF POSITION LDY TUROWS IN BOTTOM RIGHT CORNER OF JSR TUTOXY THE SCREEN.
LDA TU.PTR
STA EA
LDA TU.PTR+1
EA WILL MARK THE END OF THE BLOCK TO BE MOVED.

NOW SA, EA, AND DEST SPECIFY THE BLOCK TO BE MOUED, RND ITS DESTINATION.
JSR MOU.EA MOUE THE BLOCK.
LDY TUROWS SET TU.PTR TO BOTTOM LEFT
LDX 新
JSR TUTOXY
LDX TUCOLS CLERR THIS ROW.
LDY ${ }^{\text {\# }} 1$
JSR CLR.XY
PLA
STA 5A PARAMETERS: SA, EA, AND DEST.
PLA
STA SA+1
PLA
STA EA
PLA
STA ER+1
PLA
STA DEST
PLA
STA DEST+1 JSR TU.POP

RTS
RESTORE ZERO PAGE BYTES WE USED. RETURN.

```
4060 ;
4070 ;
4बg0 ;
4090 ;
4100 ;
4110 OFOO * * # # OFDO
4120 ;
4130 ;
4140 ;
4150 ;
4150 ;
4170 0027=
4180 205E=
4190 001B=
4200 0020=
4210 0005=
4220 005 E=
4230 0008=
4240 005%=
4250 0651=
4250 EQ7F=
4270 ;
42B0 ;
4250 ;
430日 DF00 5C
43ZO OFD1 EF
430日 0FG2 3B
4300 RFGJ GO
43DG DFO4 00
4300 0FO5 ES
43GO BFDE 2B
430日 EFET 2F
4300 OFQS EF
4300 DFDS 0N
4300 [FBA 70
43E见 DFOR 75
43EO GFGC OD
4300 GFDD EG
4300 OFEE 2D
43GG 0FOF 3D
4310 0F10 TG
4310 DF11 EJO
4310 日F12 E3
4310 EF13 00
4310 OF1400
4310 EF15 EZ
4310 DF15 78
4310] OF17 7A
431D EFID 34
4310 DF19 B0,0
4310 DF1A 33
4310 OF1B 36
4310 OF1C 1B
4310 OF1D 35
431日 OF1E 32
4310 OF1F 31.
4 3 2 0 ~ B F 2 0 ~ 2 C ~
4320 DF2.1 20
```

```
43こ0.0F22 2E
4320 OF23 EE
4ここ5 OF24 00
43ズ0 日F25 6n
43こg 0F26 2F
4320 OF27 00
4320 0F28 72
4320 0%23 D0
4320] 0F2A 65
4320 0F2B 79
4320 0F2C 09
4320 0F2D 74
4320 DF2E 77
4320 EF2F 71
4330 0F30 39
4330 0F31 00
4330 6F32 30
4330 0F33 37
4330 0F34 日8
4338 0F35 38
4330 0F3E 3C
4330 0FS7 3E
4330 0F38 E5
4330 0F39 ES
4330 BF3月 64
4330] 0F3E BC
4330 [FJC [0
4330 0F3D 67
4330 6F3E 73
433\ [FFF 61
4340
4350
4300
;
4370
43e0 ;
4330 ;
44BOD BF 40 4C
44G0 0F41 4f
4400 BF42 39
4400 EF43 00 
4400 0F44 00
44,00 DF45 4B
4406 DF45 5B
4400 0F47 5E
4410 0F48 4F
4410 OF4500
4410 BF4A 50
4410 GF4E 55
4410 BF4C OD
441906F4D 45
441日 GF4E 2D
4410 OF4F 3D
4420 0F50 55
4420 0F51 50
4420 EF52 43
442000553 00
4420 0F54 E0
4420 0F55 42
                    FOLLOWING E4 EYTES CONTAIIY
                                    FASCII CODES FOR SHIFTED KEYS.
                                    .BYTE 'LJ:',0,0,'K', BACKSL,CARAT
                                    .BYTE 'O',D,'PU',CR,'I-='
```



```
.BYTE 'G',D,'Q7', ERCKSP,'BK>fhd',D,0,'gsa'
```

```
.BYTE 'G',D,'Q7', ERCKSP,'BK>fhd',D,0,'gsa'
```

```
4420 DF5658
4420 0F57 5A
4420 OF5B 34
4420 [F55 D0
4420 OF5A 33
4420 DF5B 36
4420 0F5C 1B
4420 DF5D 25
4420 0F5E 22
4420 0F5F 21
4430 0FGO 5A
```



```
4430 DFG2 5D
4430 0FG3 4E
4430 BFG4 DO
4430 DFG5 4D
4430 DFGG 3F
4430 DFG7 DO
4440 DFGB 52
4440 DF6S DO
4440 DFGA 45
4440 EFGB 59
4440 OFGC OG
4440 QFGD 54
4440 OFGE 57
4440 0FGF.51
4450 QF7B 28
4450 DFF1. 00
4450 6F72.29
4450 0F73'27
4450 0F74 7F
4450 OF75 40
4450 0F76 BO
4450 DF77, D\
4460 0F7B 46
4450 BF7G 4B
4460 DF.7A }4
44G0 0F7B 00
44B0 OF7C OD
4460 EF7D 47
44EO GF.TE 53
4460 DFTF 41
447
4480
4500 ;
4510 ;
4520 ;
4530 ;
4540 0FB0 B0
4540 0F81 00
4540 DFB2 OD
4540 0F83 00
4540 0F84 D0
4540 0F85 DO
4540 DFB6 DO
4540 0F87 DO
4540 OFgB DO
```

- BYTE LBRAKT, SFACE, FERAKT, ' $N$ ', $\mathrm{B}_{3}^{\prime} M ?^{\prime}, ~ D$
- BYTE 'R', 0 ; EY' , TAB,' TWQ'
. BYTE ' (', 日, ' ', AFOSTR, DELETE,' @' , $\square, \square$ BYTE 'FHD', D, D, GSR'

THE FOLLOWING 128 BYTES CONTAIN CHPRACTER CODES FOR COMTROL SHIFTED KEYS. EDITOR FUNCTION KEYS ARE DEFINED.
. BYTE $\theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \$ 1 \theta, \theta, \theta, \theta, \theta, \theta$

4540 6FSS 010
4540 ©FBA 10
4540 GFBE DD
4540 OFBC DD
4540 OFBD 00
4540 BFBE 0 O
4540 DFSF 10
4550 OFSO DO
$4550 \mathrm{FFO1} 00$
4550 EFG2 03
4550079300
4550 0FG4 00
4550675500
4550 OFOS 20
4550 OFS7 00
4550 BFGS BO
4550 日F99 日
4550 RFSA 00
4550 OFGB 日 0
4550 OFSC BO
4559 EFSD OD
4550 GFGE 00
4550 OFGF 00
455 Ej OFAO DO
45EO BFAI 00
$45 E 0$ braz 20
45 SO DFAS BO
45ED GFA4 $0 \square$
4560 ＠FA5 $0 \square$
4550 OFFS BO
4560 GFAT OE
4500 BFAS 00
4565 EFAS 00
$45 E 0$ GFAR 80
4560 GFRB 00
4560 OFRC $0 \square$
4550 EFAD 0
45ES DFAE 00
4560 EFAF 00
4570 OFBO 00
4570 OFEI 0
4570 0FB2 00
4570 OFEB 00
4570 DFE 4 00
4570 0FB5 OD
4570 OFEE DO
4570 BFB7 00
4570 EFBS 06
4570 日FES 00
4570 OFEA OD
4570 OFEB 00
4570 OFBC 00
4570 EFED 0
4570 OFBE 00
4570 OFEF DO 4580 BFCD 4580 CFCl 1 BO 45EQ DFC2 DO
．BYTE $0,0,3,0,0,0,8, B, 0,0,0,0,0,0,0,0$
．BYTE $\theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta, \theta$
．BYTE 日， $0,0,0,0,0,0,0,0,0,0,0,0,0,0,0$

```
4500 DFES RO
45ED OFC4 00
4SED SFC5 DO
450% LFCE 00
45OQ EFC7 00
45EO GFCS DO
45#LS EFCS DC
45502 EFCA 00
45SO EFCB 00
4500 EFCC BO
45EO GFCD OD
4580 EFCE OU
45ED OFCF OO
4590 EFDO OD
4590 OFD1 OQ
45SO BFD2 BD
4590 0FD3 00
4590 0FD4 00
4590 EFD5 00
4550 DFDG DO
4550 0FD7 00
4590 OFDE ED
4590 DFDS DO
45SO OFDA 00
4590 DFDB NO
45S0 OFDC EO
4590 BFDD DO
45ST GFDE OD
4590 BFDF DO
4EBD EFED EO
40GE OFED OD
400D DFEZ OD
4EOQ RFE3 DV
4E\O EFE4 GO
4EOM BFES DO
4500 DFEG ED
46@O OFET OD
45DOD BFES DO
4BOD DFEG DO
4EGO DFEA ED
4EDG EFEB DO
4EGO BFEC OD
4EEOS MFED DG
4EDS DFEE OR
4EDO DFEF BD
4610 OFFO DO
4E10 EFF1 00
4510 OFFZ DD
4510 EFF3 nG
4610}
4B10 EFFS DO
4B1G DFFG DO
451.D EFFT EO
4EIE EFFE QO
4B10 DFFG DD
4E10 DFFA DD
4ELD OFFE OD
4610 UFFC DO
```

4690 OFEl $0 \square$
460 CD QFEZ OO 4EJU DFES DG 460 B BFE4 CO 4E日G BFES DG 460 OFEG ED 4 BOU EFET DO $450 \square$ DFES 20 46 OD OFEG OD 4 EGO RFEA GO 4 EDG EFEB DO 4EGO GFEC DO 4EDi GFED DO 4EDS OFEE OR 4650 GFEF 日D 4510 OFFE D $4 E 10$ EFF1 00 4510 OFFZ 0 D 4610 BFF3 DO 4610 OFF4 0 4610 EFFS CD 4610 DFF6 DO 4510 QFFT 00 $4 E 10$ EFFE DO $4510 \mathrm{OFF9} 0$ 4610 DFFA DD $4 E 10$ GFFE 00 4610 GFFC DO
. BYTE $\square, 0,0,0,0,0,0,0,0,0,0,0,0,0,0,0$
. BYTE $0, \square, 0, \square, \square, \square, \square, D, 0,0, \square, \square, 0,0,0,0$
. BYTE $0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0$

4E1月 EFFD DO
4EDGEFE ND
4EJQ BFFF DG

## Appendix DI:

## Screen Utilities

SEE CHAPTER 5 OF BEYOND GAMES: SYSTEM SOFTWARE FOR YOUR 6502 PERSONAL COMPUTER.

DUMPING $\$ 1100-\$ 11 F F$

$$
\begin{array}{lllllllllllllll} 
& 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & G & A & C & D & E & F
\end{array}
$$



## Appendix D2：

Visible Monitor（Top Level and Display Subroutines）

FPPEMDIX D2： THE UIEIBLE MONITOR（TOP LEUEL AND DISPLAY SUEROUTINES） SEE CHAPTER 6 OF EEYOND GRMES：SYSTEM SOFTWARE FOR YOUR GSDZ PERSOMAL COMPUTE
［UMPIMG $31200-* 1205$

| 1200 | 00 | 8 | C | 9000 | 31 | 05 | 12 | 08 | DS | 20 | 12 | 12 | 20 | E3 | 12 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1210 | 90 |  | 620 | 20.64 | 11 | 20 | 25 | 12 | 29 | 34 | 12 | 20 | C | 12 | 20 | AF |
| 1223 | 12 |  | 208 | －3 11 | E0 | A2． | 02 | 2 AD | 02 | 220 | $3 C$ | 11 | R2 | 19 | AD | 3 |
| 1236 | 20 |  | 1311 | 11 EG | A2 | 10 | A0 | B2 | 20 | 3 C | 11 | AO | 00 | BC | 51 | 2 |
| 1249 | S |  | 5212 | 1220 | 7 C | 11 | E | 51 | 12 | 2 AC | 51 | 12 | CO | QA | D0 | 0 |
| 12 | ED |  | Q4 41 | 4129 | 20 | 59 | 2.0 | 20 | 59 | 20 | 20 | 50 | A2 | 02 | RD | 83 |
| 12 Eb | 21 |  | $3 C 1$ | 11 AD | 05 | 12 | － | ， | 11 | 1 AD | 05 | 12 | 20 | A3 |  | 120 |
| 1275 | 75 |  | 112 | 2094 | 12 | 45 | 20 | A3 | 11 | 120 | 7F | 11 | E6 | 0 | 70 |  |
| 1280 | 2 | ？ | 7F 1 | 11 AZ | $\square 5$ | ED | 1 | 112 | 20 | 0 |  | 20 | 7 | 11 | Ea |  |
| 1299 | － | 4 D | 呵 F | F2 60 | AS | 02 | 48 | A AE | 03 | 3 AD | 05 | 12 | 85 | 02 | AD |  |
| 12月0 | 2 | 2 E | 85 | 93 7 10 | 50 | B1 | 02 | 2 AS |  | 885 | 32 | 285 | 03 | 38 | Sb |  |
| 12 Eiz | 92 | f | fig 0 | 3429 | 3 C | 11 | Ac | C 60 | 12 | 238 | C0 | 07 | 90 | 05 | AD | bu |
| 1200］ |  | BL | 00 1 | $12 \mathrm{E9}$ | CD | 12 | ค8 | A A | 7 | 710 | 1 | 1 ®0 | EO | D3 | DE | 5 |
| 12 L |  |  |  | 1114 | 00 | 00 | 0 | 0 D | 0 | 0 0日 | 0 | ］ |  | －0 | 0日 |  |

## Appendix D3:

# Visible Monitor (Update Subroutine) 

APPEIIDIX D3: THE UISIBLE MONITOR (UPDATE SUBROUTINE)

SEE CHAPTER E OF BEYOMD GAMES: SYSTEM SOFTWARE FOR YOUR 6502 PERSONAL COMPUTER.

DUMPING

$$
\begin{array}{lllllllllllllllll}
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & A & B & C & D & E & F
\end{array}
$$



## Appendix D4：

## Print Utilities

PRINT UTILITIES
SEE CHAPTER 7 OF BEYOND GAMES：SYSTEM SOFTWARE FOR YOUR 6502 PERSONAL COMPUTE

DUMPING $\$ 1400-\$ 154 F$

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | $B$ | 9 | $A$ | $B$ | $C$ | $D$ | $E$ | $F$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |


| 1400 | FF |  | FF | $0 \square$ | 20 | 00 | 句 | DC | 15 | A9 |  | FF ED | $\square 1$ | 14 | 60 | As | 昒 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 81 |  | 1 | 14 | 50 | fis | FF | 81 | 0 | 14 |  | 5 AS | D | 8 | 00 |  | 0 |
| 1420 | AS |  | FF | 吅 | 02 | 14 | 60 | AS | 日 | 81 |  | 214 | Eם | 20 | 08 | 14 | 20 |
| 30 | 14 | 1 | 142 | 20 | 20 | 14 | ES | 20 | DE | 14 |  | 2019 | 14 | 20 | 25 | 1 | 60 |
| 40 | CS | 0 | 牱 F | FD 2 | 24 | 8 D | Q3 | 14 | AD | 01 |  | 14 Fb | －6 | AL | 83 | 1 | 2 |
| 1450 | 55 | 1 | 14 A | AD | － | 14 | FO | 06 | AD | 03 |  | 20 | 6C | 14 | AD | 02 | 14 |
| 1480 | F | － | 06 A | AD | $\square$ | 14 | 20 | 6F | 14 | 6 B |  | 6C DA | 10 | EC | OC |  | 6C |
| － | OE | 1 | 10 A | A9 | 日D | 20 | 40 | 14 | AS | DP |  | 2040 | 14 | 6b | AS | 20 | 20 |
| 14 | 48 | 1 | 14 | 60 | 48 | 4R | 4A | 4A | 4R | 20 |  | E6 11 | 20 | 40 | 14 | ¢ | 20 |
| 1490 | BG | 1 | 112 | 20 | 40 | 14 | 60 | A9 | 20 | 8 E |  | 0414 | 48 | AE | 4 |  | － |
| 14 RO | Q | C | CE | 04 | 14 | 20 | 40 | 14 | 68 | 18 |  | 30 Fb | 68 | 60 | 8E | 0 | 4 |
| 14 BD |  | － | 041 | 14 | Fb | 05 | CE | －4 | 14 | 20 |  | 14 | 18 | 90 | F2 | 60 | E |
| 1400 | 95 | 1 | 14 | B5 | 01 | 48 | E5 | Q0 | 43 | AE |  | 0514 | A1 | 0 | C9 |  | b |
| 14 | ロ | F | F6 | 00 | D0 | 02 | FS | Q1 | $\square$ | 40 |  | 18 | 90 | EB | 68 |  | BQ |
| 0 |  | 89 | 95 | 01 | 60 | 68 | AR | 68 | A8 | 20 |  | 15 | BE | $\square 5$ |  |  |  |
| 14 F |  | 2 | 20 | 01 | 13 | 20 | $0 \square$ | 13 | 20 | S |  | C9 | FF | F | － |  | 40 |
| 1500 |  | 14 | 18 | 90 | FO | RE | 05 | 12 | AC | Q6 |  | 1220 | 2B | 15 | 㫜 |  |  |
| 1510 |  | E | E0 | 6 | a | － | 14 | 68 | QD | 07 |  | 14 AD | ט6 | 12 | 4 | AD | D 0 |
| 1520 |  | 12 | 46 | AD | 07 | 14 | 8 | AD | 是 | 14 |  | 60 | 68 | BD | ט日 |  | 4 6 |
| 1530 |  | D | 日7． | 14 | 68 | 8D | 05 | 12 | 68 |  |  | D6 12 | AD | －7 | 14 |  |  |
| 1540 |  | 061 | 14 | 48 | 60 | 00 | ］ | Q0 | 日 | －0 | 0 | 00 | － | 日 | 日0 |  |  |

# Appendix D5： Two Hexdump Tools 

RPPENDIX $\mathrm{D5}:$
TWO HEXDUMP TOOLS
SEE CHAPTER 8 OF BEYOND ．GRMES：SYSTEM SOFTWARE FOR YOUR GS02 PERSONAL COMPUTER

DUMPING \＄1550－\＄17AF

$$
\begin{array}{llllllllllllllll}
0 & 1 & 2 & 3 & 4 & 5 & 5 & 7 & B & 9 & A & B & C & D & E & F
\end{array}
$$

| 1550 | 000 |  |  |  |  | 17 |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1550 | AD | 05 | 12 | 29 | Fs | 8 D | 05 | 12 | 20 | 72 | 14 | 20 | 72 | 14 | 20 | A1 |
| 70 | 152 | 20 | 72 | 14 | 20 | 71 | 14 | 20 | 9A | 15 | 20 | 0 D | 13 | AD | 65 | 12 |
| 1580 | 29 | 0 | D9 | F | 20 | 7 | 14 | AD | 05 | 12 | 29 | UF | 吅 | 83 | 28 | 72 |
| 1590 | 14 C | C | 50 | 15 | D® | D8 | 20 | QE | 14 | 60 | 20 | 94 | 12 | 20 | 83 | 14 |
| 15AR | 50 A | A | 05 | 12 | 20 | 83 | 1 | AD | 05 | 12 | 20 | 83 | 14 | B0 | 2 | 9 |
| 15BD | 5 | 20 | E | 15 | 20 | A | 17 | 20 | 14 | 14 | 20 | EB | 16 | 20 | 4 | 17 |
| 15 CD | 10 F | F | 20 | 72 | 14 | 20 | 1A | 14 | 60 | 20 | 00 | 11 | 29 | 08 |  | 20 |
| 15D0 | E4 | 14 | 7 F | QD | 50 | 52 | 49 | 4 | 54 | 49 | 4 E | 47 | 20 | 48 | 45 | 58 |
| 15Eb | 44.5 | 55 | 4D | 50 | 01 | 日ค | Q | FF | 60 | 2 | D8 | 14 | 20 | E4 |  | 7F |
| 15 Fb | Q1 | ถค | 53 | 45 | 54 | 20 | 53 | 54 | 41 | 52 | 5 | 49 | 4E | 47 | 20 | 41 |
| 1600 | 4 | 44 | 52 | 45 | 53 | 53 | 20 | 41 | 4E | 44 | 20 | 50 | 52 | 45 |  |  |
| 1610 | 2 | 22 | 51 | 22 | $2 E$ | FF | 20 | 07 | 12 | 20 | 67 | 16 | 20 | 08 |  | 28 |
| 1620 | E4 1 | 14 | 7 | 01 | OA | 53 | 45 | 54 | 20 | 45 | 4E | 4 | 20 | 41 |  |  |
| 1630 | 4 | 45 | 53 | 53 | 20 | 41 | 4E | 4 | 20 | 50 | 52 | 45 | 53 | 53 |  |  |
| 16 | 512 | 22 | 2E | FF | 20 | 07 | 12 | 38 | AD | 65 | 12 | CD | 53 | 15 | 90 | 24 |
| 16 | 吅 | 83 | AD | 05 | 12 | CD | 5 | 15 | 90 | 18 | AD | 06 | 12 | 9D |  |  |
| 16ED | － | － | 12 | 81 | 54 | 15 | ED | AD | 06 | 12 | 8D | 53 | 15 | AD | 05 | 12 |
| 16 | 8 B 5 | 52 | 15 | ED | 20 | E | 14 | 7 | $0 \square$ | 日A | A | － | 20 | 45 |  |  |
| 1680 | 5 | 52 | 21 | 21 | 21 | 20 | 45 | 4E | 44 | 29 | 41 | 44 | 44 | 52 |  |  |
| 16 | 2 | 20 | 4 C | 45 | 53 | 53 | － | 5 | 48 | 41 | 4 L | 20 | 53 | 54 |  |  |
| A | 54 | 20 | 41 | 44 | 44 | 5 | 45 | 53 | 53 | 2C | 20 | 57 | 48 | 49 |  |  |
| 16BD | 204 | 49 | 53 | 20 | FF | 20 |  | 16 | 4 | 1 C | 16 | AS | 24 | 2 |  |  |
| 15C0 | AD 5 | 53 | 15 | 20 | 83 | 14 | AD | 52 | 15 | 20 | 83 | 14 | E0 | AS | 24 | 20 |
| D | 4 D 1 | 14 | AD | 55 | 15 | 20 | 83 | 14 | AD | 54 | 15 | 20 | 83 | 14 | 51 |  |
| IGED | BB | 15 | A9 | 2D | 20 | 40 | 14 | 20 | CD | 16 | 50 | 20 | E4 | 14 | 7 | F 0 L |
| 15 | 日R | －A | 44 | 55 | 4D | 50 | 49 | 4E | 47 | 20 | FF | 20 | DF | 15 | 20 | －72 |
| 1700 | 14 | 20 | E4 | 14 | 7F | ตR | ® | 20 | 20 | 20 | 20 | 20 | 20 | 20 | 20 | 30 |
| 710 | 20 | 20 | 31 | 20 | 20 | 32 | 20 | 20 | 33 | 29 | 0 | 34 | 20 | 20 | 35 | 520 |



## Appendix D6:

# Table-Driven Disassembler (Top Level and Utility Subroutines) 

APPEMDIX DS: table-driven disassembler (top level and utility subroutines)

SEE CHAPTER 9 OF EEYOND GAMES: SYSTEM SOFTWARE FOR YOUR ESD 2 PERSOMAL COMPUTER

DUNPING $\$ 1900-\$ 1$ B3F

|  |  |  |  |
| :---: | :---: | :---: | :---: |



## Appendix D7:

# Table-Driven Disassembler (Addressing Mode Subroutines) 

FPFERHIX DT:<br>TAELE-DRIUEN DISASSEMBLER<br>(ADDRESSING MODE SUBROUTINES)

SEE CHIPTER 9 OF BEYOND GAMES: SYSTEM SOFTWFRE FOR YOUR ES0Z PERSONAL COMPUTE

DUHTPING E1A4B-\#154F

| $B$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | $B$ | $G$ | A | B | $C$ | $D$ | $E$ | $F$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |



## Appendix D8:

# Table-Driven Disassembler (Tables) 

APPENDIX DS:
TRELE-DRIUEN BISRSSEMBLER (TAELES)
SEE CHAPTER 9 OF BEYOND GAMES: SYSTEM SOFTHARE FOR YOUR G502 PERGONAL COMPUTER

DUMPING \#1E50-第1DFF
$\begin{array}{llllllllllllllll}\square & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & A & B & C & D & E & F\end{array}$
















## Appendix D9:

## Move Utilities

APFENDIX DS:
MOVE UTILItiES
SEE CHAPTER IO OF BEYOND GAMES: SYSTEM SOFTWARE FOR YOUR ESOZ PERSOMAL COMPIJTER

DIMPTNG $\# 17 E G-\$ 18 F F$

| $\square$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | $A$ | $B$ | $C$ | $D$ | $E$ | $F$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |



## Appendix DIO:

## Simple Text Editor

APPENDIX DID:

A SIMPLE TEXT EDITOR
SEE CHAPTER 11 OF BEYOND GAMES: SYSTEM SOFTWARE FOR YOUR GSD2 PERSONAL COMPI BY KEN SKIER

DUMPING \$1EOD-\$1FFF

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | $A$ | $B$ | $C$ | $D$ | $E$ | $F$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |



# Appendix DII: Extending the Visible Monitor 

APPENDIX DI1:
EXTENDING THE UISIBLE MONITOR
SEE CHAPTER 12 OF BEYOND GAMES: SYSTEM SOFTWARE FOR YOUR 6502 PERSONAL COMPUTER.

DUMPING : 10 ED - F 10FF

$$
\begin{array}{lllllllllllllll}
0 & 1 & 3 & 4 & 5 & 5 & 7 & 9 & A & B & C & D & F
\end{array}
$$

| 1080 | $\mathrm{C9}$ | 50 | DE | 09 | AD | 00 | 14 | 49 | FF | 8 D | 0 D | 14 |  |  | 5 | 5 D0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 10 CD | 09 | AD | 02 | 14 | 49 | FF | 81 | 02 | 14 | 60 | Cs | 48 | D | 01 | A | D00 |
| 1009 | 14 | DD | 04 | 20 | 57 | 15 | 60 | 20 | AE | 15 | 60 | C9 | 4I | D0 | 0 | 420 |
| 10EC | B4 | 17 | E ${ }^{\text {c }}$ | Cs | 3 F | D0 | 0 B | AD | 80 | 14 | D | 04 | 20 | 09 | 1 | 0 |
| 10\%0 | 28 | 26 | 15 | 69 | Cg | 54 | DC |  | 20 | 02 | 15 | 69 | ED | - | 入曰 | 0 |

# Appendix EI: <br> <br> Screen Utilities 

 <br> <br> Screen Utilities}

THE FOLLONING DATA STATEMENTS CONTAIN DECIMAL OBJECT CODE AND CHECKSUMS FOR MEMORY FROM 4352 TO 4607 SUITRBLE FOR LOADING WITH THE BASIC OBJECT CODE LOADER.

| 1000 DATA |
| :--- |
| 1001 DATA | 4352,32,196,17,32,43,17,174,3,4866

```
1029 DATA 4584, 日, 日, В, ロ, 日, 日, 日, 日, 4584
1030 DATA 4592, 0, 日, 日, 日, 日, 日, 日, 日, 4592
1031 DATA 4600, 0, 0, D, 日, 0, 0, 日, B, 460日
1032 END
```

OK

## Appendix E2：

## Visible Monitor（Top Level and Display Subroutines）

APPENDIX E2 UISIBLE MONITOR（TOP LEUEL \＆DISPLAY SUBS）

THE FOLLOWING DATA STATEMENTS
CONTAIN DECIMAL OBJECT CODE AND CHECKSUMS FOR MEMORY FROM 450 S TO 4831
SUITABLE FOR LOADING WITH
THE EASIC OBJECT CODE LOADER．

| D | 4608，日，12，日，日，49，177，252，日， |
| :---: | :---: |
| 1101 DATA | 4616，216，32，18，18，32，227，18， |
| 1102 DATA | 4524，144，246，32，195，17，32，37，18， 5345 |
| 1103 DATA | 4632，32，52，18，32，52，18，32，175， 5083 |
| 1164 DATA | 4640，16，32，211，17，36，162，2，160， 5338 |
| 1105 DATA | $4648,2,32,60,17,162,25,160,3,5109$ |
| 1106 DATA | 4556，32，19，17，96，162，13，160，2，5157 |
| 1107 DATA | 4564，32，60，17，160，0，140，81，18， 5172 |
| 1108 DRTA | 4672，185，82，18，32，124，17，238，81，54 |
| 1109 DRTA | 4680，18，172，81，18，132，10，208，240， |
| 1110 DATA | 4688，96，10，65，32，32，88，32，32，5075 |
| 1111 DATA | 4656，89，32，32，80，152，2，160，3， 5256 |
| 1112 DATA | 4704，32，60，17，173，6，18，32，163，5205 |
| 1113 Dfita | 4712，17，173，5，18，32，163，17，32， 5169 |
| 1114 DATA | 4720，127，17，32，148，18，72，32，163， 5329 |
| 1115 DATA | 4728，17，32，127，17，104，32，124，17， 51 |
| 1116 DATA | 4736，32，127，17，162， $0,189,1,18,5282$ |
| 1117 DATA | 4744，32，163，17，32，127，17，232，224，5563 |
| 1118 DATA | 4752，4，208，242，96，165，2，72，165， 5707 |
| 1119 DATA | 4760，3，173，5，18，133，2，173，5，5273 |
| 1120 DATA | 4768，18，133，3，150，0，177，2，168， 5429 |
| 1121 DATA | 4776，104，133，2，134，3，152，35，162，5562 |
| 1122 DATR | 4784，2，160，4，32，60，17，172，0，5231 |
| 1123 DATA | 4792，18，56，192，7，144，5，160，0， 537 |
| 1124 DATA | 4800，140，日，18，185，205，18，158，173，5707 |
| 1125 DATA | 4808，7，16，145，©，96，З，6，8，5089 |

```
1126 DATA 4815, 11, 14, 17, 20, ロ, D, ロ, ロ, 4878
1127 DATA 4824, ロ, 0, 日, D, 日, 0, D, D, 4824
1128 END
```


## Appendix E3：

# Visible Monitor（Update Subroutine） 

FiPFENDIX E3
UISIELE MONITOR（UPDATE SUBROUTINE）

THE FOLLOWING DATA STATEMENTS CONTAIN DECIMAL OBJECT CODE RMD CHECKSLIMS FOR MEMORY FROM 4832 TO 5119 SUITRBLE FOR LOFDING WITH THE BGSIC OBJECT CODE LOADER．

| 1200 DATA | 4832，108，8，16，32，224，18，201，62，5501 |
| :---: | :---: |
| 1201 DATA | 4840，208，16，238，0，18，173，日，18， 5511 |
| 1202 DATA | 4848，201，7，208，5，169，0，141，0， 5579 |
| 1203 DATA | 4856，18，96，201，60，200，11，206，ロ，5656 |
| 1204 DATA | 4564，18，16，5，169，6，141，D，18， 5237 |
| 1205 DATA | 4872，56，201，32，208，5，238，5，18， 5679 |
| 1205 DATA | 4850，208，3，238，5，15，9E，201，13， 5663 |
| 1207 DATA | 4683，208，12，173，5，18，203，3，206， 5721 |
| 1208 DATA | 4896，5，18，20E，5，18，96，174，0，5419 |
| 1209 DATA | 4994，18，224，2，208，27，156，165，0，5716 |
| 1210 DATfi | 4912，72，166，1，173，5，13，133，0，5400 |
| 1211 DATA | 4920，173，6．18，133，1，152，160，口， 5553 |
| 1212 BATA | 4328，145，日，134，1，104，133，ロ，96， 5541 |
| 1213 DATA | 4936，201，71，208，35，172，3，18，174， 5816 |
| 1214 DATA | 4944，2，1e，173，4，18，72，173，1，5485 |
| 1215 DRT | 4952，10，40，32，108，19，8，141，1， 5319 |
| 1216 DATA | 4950，18，142，2，18，140，3，18，104， 5405 |
| 1217 DAT | 4969，141，4，18，96，165，5，18，72，5430 |
| 1218 DATA | 4975，32，213，19，45，75，158，104，152，5787 |
| 1215 DATA | 4984，174，0，18，200，20，162，3，24， 5593 |
| 1220 DFTA | 4992，14，5，18，46，6，18，292，15， 5317 |
| 1221 пATA | 5000，246，152，13，5，10，141，5，18， 5588 |
| 1222 DATA | 5008，96，224，1，205，24，41，15，72，5683 |
| 1223 DATA | 5016，32，148，18，10，10，10，10，41， 5295 |
| 1224 DRTA | 5624，240，141，172，19，104，13，172，19，5904 |
| 1225 DATA | 5032，32，45，19，95，15，202，202，202， 5845 |
| 1226 DATA | 5040，160，3，24，30，1，18，136，16，5425 |
| 1227 DATA | 5049，249，29，1，18，157，1，18，96， 5617 |
| 1228 DATA | 5056，104，201，127，200，4，32，0，17， 5749 |

1229 DATA 5064， $96,201,81,208,4,104,104,40,5902$
1230 DATA $5072,95,32,16,16,56,56,233,48,5665$
1231 DATA $5030,144,15,201,10,144,14,233,7,5848$
1232 DATA 5088，291，16，176，5，56，201，10，176， 5929
1233 DATA 5056，3，169，255，96，162，日，95，0， 5877
1234 DATA $5104, \square, \square, \square, \square, 0, \square, \square, \square, 5104$
1235 DATA 5112，日，日，ロ，ロ，日，ロ，日，日， 5112
1236 END

## Appendix E4：

## Print Utilities

RPPENDIX E4
PRINT UTILITIES

THE FOLLOWING DATA STATEMENTS
CONTAIN DECIMRL OEJECT CODE AND CHECKSUMS FOR MEMORY FROM 5120 TO 5455
SUITABLE FOR LOADING WITH
THE ERSIC OBJECT CODE LOADER．

| 1300 DATA | 5120，日，255，日，日，日，日，日，日， |
| :---: | :---: |
| 1301 DATA | 5128，169，255，141，1，20，96，169，0， 5979 |
| 1302 DATA | 5135，141，1，20，96，169，255，141，0， 5959 |
| 1303 DATA | 5144，20，96，169，日，141，0，20，56， 5666 |
| 1304 DATA | 5152，169，255，141，2，20，96，169，日，ЕИ04 |
| 1365 DATA | 5160，141，2，20，56，32，8，20，32， 5511 |
| 1305 DATA | 5168，20，20，32，32，20，96，32，14， 5434 |
| 1367 DATA | 5176，20，32，26，20，32，38，20，96，5460 |
| 1306 DATA | 5184，201，Q，240，36，141，3，20，173， 5998 |
| 1309 DATA | 5192，1，20，240，6，173，3，20，32，5687 |
| 1310 DAT | 5200，105，20，173，0，20，240，6，173， 5937 |
| 1311 DATA | 5200，3，20，32，108，20，173，2，20，5596 |
| 1312 DA | 5216，240，6，173，3，20，32，111，20， 5821 |
| 1313 DATA | 5224，96，108，10，16，108，12，16，106， 5698 |
| 1314 DATA | 5232，14，15，169，13，32，64，20，169，5729 |
| 1315 DRTA | 5240，10，32，64，20，96，169，32，32，5695 |
| 1316 DAT | 5248，64，20，56，72，74，74，74，74，5796 |
| 1317 DATA | 5255，32，182，17，32，54，20，104，32， 5739 |
| 1318 DRTA | 5264，182，17，32，64，20，96，159，32， 5876 |
| 1319 DATA | 5272，142，4，20，72，174，4，20，240，5948 |
| 1320 DATA | 5280，10，206，4，20，32，64，20，104，5740 |
| 1321 DATA | 5268，24，144，240，104，96，142，4，20，6062 |
| 1322 DATA | 5296，174，4，20，240，9，206，4，20， 5973 |
| 1323 IATA | 5304，32，114，20，24，144，242，96，142，6118 |
| 1324 DA | 5312，5，20，181，1，72，181，0，72， 5844 |
| 1325 DATA | 5320，174，5，20，161，日，201，255，240， 6376 |
| 1326 DRTA | 5328，12，246，9，208，2，246，1，32，6075 |
| 1327 DATA | 5336，64，20，24，144，235，104，149，0，6076 |
| 1328 DATR | 5344，124，149，1，96，104，170，104，1ED，E240 |


| 1329 DATA | 5352，32，18，21，142，5，18，140，6， 5734 $5360,18,32,13,19,32,13,19,32,5538$ |
| :---: | :---: |
| 1331 DATA | 5358，148，18，201，255，240，6，32，64， 6332 |
| 1332 DATA | 5376，20，24，144，240，174，5，18，172， 6173 |
| 1333 DATA | 5384，6，18，32，43，21，152，72，138， 5866 |
| 133.4 DATA | 5392，72，96，104，141，6，20，104，141，6076 |
| 1335 DATA | 5400，7，20，173，6，18，72，173，5， 5874 |
| 1336 DATA | 5408，18，72，173，7，20，72，173，5， 5949 |
| 133 D DATA | 5415，20，72，96，104，141，6，20，104， 5979 |
| 1335 DATA | 5432，141，6，18，173，7，20，72，173， 6042 |
| 1345 DATFI | 5440，6，20，72，96，0，日，日，日，5634 |
| 1341 DATA | 5448，D，ロ，日，D，日，日，ロ，日， 5448 |

## Appendix E5:

## Two Hexdump Tools

APPEMDIX E5
TWO HEXDUMP TOOLS

THE FOLLOWING DATA STATEMENTS CONTAIN DECIMAL OBJECT CODE AND CHECKSUMS FOR MEMORY FROM 5456 TO 6063 SUITABLE FOR LOADING WITH THE BASIC OBJECT CODE LOADER.

```
1400 DATA
1401 DATA
1402 DATA
1403 DRTA
1404 DATA
1405 DATA
1406 DATA
1407 DRTR
1408 DATA
1409 DATA
1410 DATA
1411 DATA
1412 DATA
1413 DRTA
1414 DATA
1415 DATA
1416 DATA
1417 DATA
1418 DATA
1419 DATA
1420 BATA
1421 DATA
1422 DATA
1423 DRTA
1424 DATA
1425 DATA
1426 DATA
1427 DATA
1428 DATA
5456, 0, 4, 0, 0, 255, 255, 0, 32, 5002
5464, 8, 20, 173, 81, 21, 141, 80, 21, 6009
5472, 173, 5, 18, 41, 248, 141, 5, 18, 6121
5480, 32, 114, 20, 32, 114, 20, 32, 161, 6005
5488, 21, 32, 114, 20, 32, 125, 20, 32, 5654
5496, 154, 21, 32, 13, 19, 173, 5, 18, 5931
5504, 41, 7, 208, 240, 32, 114, 20, 173, E339
5512, 5, 18, 41, 15, 208, 3, 32, 114, 5948
5520, 20, 206, 80, 21, 208, 216, 32, 14, 6317
5528, 20, 96, 32, 148, 18, 32, 131, 20, 6025
5536, 96, 173, 6, 18, 32, 131, 20, 173, 6185
5544, 5, 18, 32, 131, 20, 95, 32, 201, 6079
5552, 21, 32, 233, 21, 32, 160, 23, 32, 6106
5560, 20, 20, 32, 235, 22, 32, 66, 23, 5010
5568, 16, 251, 32, 114, 20, 32, 26, 20, 6079
5576, 96, 32, Ø, 17, 32, 8, 20, 32, 5813
5564, 228, 20, 127, 13, 80, 82, 73, 78, 6285
5592, 84, 73, 78, 71, 32, 72, 69, 88, 6153
5500, 68, 85, 77, 80, 13, 10, 10, 255, 6198
5608, 96, 32, 8, 20, 32, 228, 20, 127, 6171
5E16, 13, 10, 83, 69, 84, 32, 83, 84, 6074
5624, 65, 82, 84, 73, 78, 71, 32, 65, 6174
5532, 68, 68, 82, 69, 83, 83, 32, 55, 6182
5640, 78, 68, 32, 80, 82, 69, 83, 83, 6215
5648, 32, 34, 81, 34, 46, 255, 32, 7, 6159
5655, 18, 32, 103, 22, 32, 8, 20, 32, 5923
5664, 228, 20, 127, 13, 10, 83, 63, 84, 5298
5572, 32, 69, 78, 68, 32, 65, 68, 68, 6152
5680, 82, 69, 83, 83, 32, 65, 78, 58, 6240
```

| 1429 DATA | 5688，32，80，82，69，83，83，32，34，6183 |
| :---: | :---: |
| 430 DATA | 5696，81，34，46，255，32，7，18，56， 6225 |
| 1431 DA | 5704，173，6，16，205，83，21，144，36，6390 |
| 1432 DATA | 5712，208，8，173，5，18，205，82，21，E432 |
| 1433 DA | 5720，144，26，173，E，18，141，85，21， 6334 |
| 1434 DATA | $5728,173,5,18,141,84,21,96,173,5439$ |
| 1435 Df | $5736,5,18,141,83,21 ; 173,5,18,5201$ |
| 36 DATA | 5744，141，82，21，96，32，228，20，127，E491 |
| 1437 DA | 5752，13，10，10，10，32，69，82，82，G060 |
| 1438 DATA | 5760，79，82，33，33，33，32，69，78， 6199 |
| DA | 5768，68，32．65，68，68，82，69，83， 5303 |
| 1440 DATA | 5776，83，32，76，69，83，83，32，84，6318 |
| 1441 DATA | 5734；72，65，78，32，83，84，65，82， 6345 |
| 1442 DR | 5792，84，32，65，68，68，82，69，83，6343 |
| 1443 DA | 5800，83，44，32，87，72，73，67，72，6330 |
| 1444 DATA | 5808，32，73，83，32，255，32，187，22， 6524 |
| 1445 DA | 5816，76，28，22，169，36，32，64，20， 6263 |
| 1445 DATA | 5824，173，83，21，32，131，20，173，82， 6539 |
| 1447 DA | $5832,21,32,131,20,56,169,36,32,6369$ |
| 1448 DATA | 5840，64，20，173，85，21，32，131，20，6386 |
| 1449 DA | 5848，173，54，21，32，131，20，96，32， 6437 |
| 1450 DATA | 5656，187，22，169，45，32，64，29，32，E427 |
| 1451 DA | 5664，205，22，96，32，228，20，127，13，6507 |
| 1452 DRTA | 5872，10，10，68，85，77，80，73，78， 6353 |
| 1453 DA | 5880，71，32，255，32，223，22，32，114，6651 |
| 1454 EATA | 5688，20，32，228，20，127，10，10，32， 6367 |
| 1455 Df | 5896，32，32，32，32，32，32，32，48，6158 |
| 1456 DATA | 5904，32，32，49，32，32．50，32，32，6195 |
| DA | 5912，51，32，32，52，32，32，53，32，6228 |
| 1458 DATA | 5920，32，54，32，32，55，32，32，56，6245 |
| 1459 DATA | $5928,32,32,57,32,32,65,32,32,6242$ |
| 1460 DA | $5936,66,32,32,67,32,32,68,32,6297$ |
| DAT | $5944,32,69,32,32,70,13,10,10,6212$ |
| 1462 DATR | 5952，255，96，32，114，20，173，5，18，6665 |
| 1463 DRTA | 5950，72，41，15，141，85，21，184，41， 6481 |
| 1464 DATA | 5960，240，141，5，18，32，161，21，162，6748 |
| 1465 DATA | 5976，3，32，150，20，173，86，21，240， 6701 |
| 1456 DATA | $5584,13,162,3,32,150,20,32,13,6403$ |
| 1467 DAT | 5992，19，206，86，21，208，243，32，154，6961 |
| 1468 DATA | 6000，21，32，125，20，32，131，23，48， 6432 |
| 1469 DA | 6005，9，173，5，18，41，15，201，9， 5470 |
| 147 D DATA | E01E，208，236，96，55，173，5，18， 205,7014 |
| 1471 DA | 6024，E5，21，144，11，208，15，56，173， 6737 |
| 1472 DATA | 6032，5，18，205，84，21，176，6，32， 6579 |
| 1473 DATA | 6040，13，19，169，ロ，96．169，255，96， 6857 |
| 1474 DATA | 6048，173，82，21，141，5，18，173，83，E744 |
| 1475 DATA | 6056，21，141，6，18，96，日，日，Д，6336 |
| 5 EN |  |

## Appendix E6:

## Table-Driven Disassembler (Top Level and Utility Subroutines)

APPENDIX EG

DISASSEMBLER (TOP LEUEL \& UTILITY SUBS)

THE FOLLOWING DATA STATEMENTS CONTAIN DECIMAL OBJECT CODE GIND CHECKSUMS FOR MEMORY FROM 6400 TO ET19 SUITABLE FOR LOADING WITH THE BASIC OBJECT CODE LORDER.

| D |  |
| :---: | :---: |
| 1 DAT | 6408, 16, 32, 8, 20, 173, 0, 25, 141 |
| 2 DAT | 6416, 1, 25, 169, 255, 141, 84, 21, 141, 72 |
| 1503 DRTA | 6424, 85, 21, 32, 114, 20, 32, 125, 25, 6878 |
| 04 DATA | 6432, 205, 1, 25, 206, 248, 36, 32, 26, 7274 |
| 1505 DATA | 6440, 20, 32, 8, 20, 32, 228, 20, 127, 6927 |
| 1505 DATA | 6448, 13, 10, 32, 32, 32, 32, 32, 30, 6711 |
| 1507 DATA | 6456, 82, 73, 78, 84, 73, 78, 71, 32, 7027 |
| 1508 DATA | 6464, 68, 73, 83, 65, 83, 83, 69, 77, 7065 |
| Df | 6472, 66, 76, 69, 82, 46, 13, 10, 255, 7089 |
| 1510 DATA | 6460, 32, 233, 21, 32, 20, 20, 32, 225, 7038 |
| DA | 6488, 20, 127, 13, 10, 68, 73, 83, 65, 6947 |
| 1512 DATA | 6496, 83, 83, 69, 77, 66, 76, 73, 78, 7101 |
| 3 DA | 6504, 71, 32, 255, 32, 223, 22, 32, 160, 7331 |
| $15: 4$ IATA | 6512, 23, 32, 114, 20, 32, 125, 25, 16, 6893 |
| DA | 6520, 251, 32, 26, 20, 96, 32, 148, 18, 7143 |
| 5 DATA | 6528, 72, 32, 146, 25, 32, 125, 20, 104, 7084 |
| 7 D | 6536, 32, 175, 25, 32, 1, 26, 32, 131, 6990 |
| 1 DAT | 6544, 23, 96, 162, 3, 142, 2, 25, 170, 7167 |
| 19 DATA | 6552, 189, 0, 28, 170, 189, 80, 27, 142, 7377 |
| 1520 | 6560, 3, 25,.32, 54, 20, 174, 3, 25, 6506 |
| 1521 DRTA | 6568, 232, 206, 2, 25, 208, 238, 36, 170, 77 |
| 22 Df | 6576, 189, 0, 29, 170, 32, 184, 25, 56, 730 |
| 1523 DATA | 6584, 189, 27, 27, 141, 4, 25, 232, 189, 7418 |
| 24 DAT | 6592, 27, 27, 141, 5, 25, 108, 4, 25, 595 |
| 1525 DATA | 6600, 32, 13, 19, 32, 154, 21, 56, 32, 6999 |


| 1526 DATA | 6608，13，19，32，148，18， $72,32,13,6955$ |
| :---: | :---: |
| 1527 DATA | $6616,19,32,154,21,104,32,131,20,7129$ |
| 1529 DATA | 6524，55，169，40，298，2，169，41，32， 7381 |
| 1530 DATA | 6640，169，88，32，64，20， $96,169,44,7327$ |
| 1531 DRTA | 6648，32，64，20，169，89，32，64，20， 7138 |
| 1532 DATA | 6656，96，141，7，25，142，6，25，202， 7300 |
| 1533 DA | 6664，48，6，32，26，19，202，16，250， 7263 |
| 1535．DRTA | 6672，8，216，56，173，8，25，233，4， 7335 658D，237，7，25，40，170， 32 |
| 1536 DATA | 6688，32，161，21，32，125，20，32，154， 72 |
| 1537 DRTA | 6696，21，32，13，19，206，6，25，16， 7034 |
| 1538 DATA | 6704，242，32，26，19，32，114，20，36， 7285 |
| 1540 END | ถile，0，0，0，0，日，0，日，日，6712 |

## Appendix E7:

## Table-Driven Disassembler (Addressing Mode Subroutines)

THE FOLLOWING DATA STATEMENTS CONTAIN DECIMAL OBJECT CODE AND CHECFSLUMS FOR MEMORY FROM E720 TO EG91 SUITAELE FOR LOADING WITH
THE BASIC OEJECT CODE LOADER.

| DA | 6720, 32, 207, 25, 162, 2, 169, 4, 96, 7417 |
| :---: | :---: |
| 1601 DATA | 6720, 32, 54, 26, 32, 235, 25, 152, 2, 7306 |
| 02 DA | 6736, 163, 6, 96, 32, 64, 26, 32, 246, 7407 |
| 1603 DATA | 6744, 25, 162, 2, 169, 5, 96, 169, 65, 7438 |
| 16014 DATA | 6752, 32, 64, 20, 152, 0, 169, 1, 96, 7295 |
| 1505 DATA | 6760, 162, 日, 169, $\square, 96,169,35,32$, |
| IEDE DATA | 6768, 64, 20, 169, 36, 32, 64, 20, 32, 7205 |
| 1507 DATA | 6776, 200, 25, 162, 1, 169, 4, 96, 32, 7465 |
| 160 DATA | 6784, 225, 25, 32, 64, 26, 32, 229, 25, 7442 |
| DR | 6792, 169, 6, 162, $2,96,32,225,25,7509$ |
| 1610 DATA | 6800. 32, 232, 26, 32, 229, 25, 162, 1, 7539 |
| 1611 data | 550, 159, 8, 96, 32, 225, 25, 32, 219, |
| 1612 DATA | 6816, 26, 32. 229, 25, 32, 246, 25, 152, 7593 |
| 1613 DAT | 6824, 1, 163, 8, 56, 32, 13, 19, 32, 7134 |
| 1 1614 Dfta | 5832, 18, 21, 32, 148, 18, 72, 32, 13, 7185 |
| 1515 DA | 6840, 19, 104, 201, D, 16, 3, 206, 6, 7395 |
| 1516 DATA | 6848. 18, 8, 216, 24, 109, 5, 18, 144, 7390 |
| 1517 DATA | 6856, 3, 238, 6, 18, 141, 5, 18, 40, 7325 |
| 3 DAT | 5864, 32, 161, 21, 32, 43, 21, 162, 1, 7337 |
| 1619 DATA | 6072, 169, 4, 96, 159, D, 32, 131, 20, 7493 |
| 1620 DA | 6880, 32, 200, 25, 162, 1, 169, 4, 96, 7569 |
| 1621 DATA | 685e, 32, 219, 26, 32, 235, 25, 162, 1, 7620 |
| 1622 DA | 685E, 169, 5, 95, 32, 219, 26, 32, 246, 7722 |
| 1523 DATA | 6904, 25, 152, 1, 16S, 6, 96, 104, 104, 7571 |
| 1524 DA | 6512, 104, 104, 32, 131, 23, 48, 13, 32, 7393 |
| $1 E 25$ DAT | 6920, 148, 18, 201, 255, 240, 6, 32, 64, 788 |

```
162E DATA 6322, 20, 24, 144, 238, 32, 114, 20, 32, 7552
1E27 DATA ES36, 131, 23, 5E, 104, 26, 54, 26, 109, 7545
1628 DATA 6544, 26, 219, 26, 232, 26, 243, 26, 64, 7805
1E29 DRTA 5952, 26, 72, 26, 83, 26, 104, 26, 172, 7487
163] LATA ESER, 26, 141, 25, 155, 26, 127, 26, 254, 7741
1631 DRTA ESEB, 26, 0, 0, 0, 0, 0, 0, 0, 6594
1632 DATA EST5, ロ, ロ, ロ, 日, 日, ロ, 0, 0, 6976
1633 DATA 6984, D, ロ, ロ, ロ, ロ, ロ, ロ, ロ, 5384
```


## Appendix E8:

# Table-Driven Disassembler (Tables) 

APPENDIX E8

DISASSEMELER (TAELES)

THE FOLLOWING DRTR STATEMENTS CONTAIN DECIMAL OEJECT CODE RND CHECKSUMS FOR MEMORY FROM G932 TO $7 E 79$ SUITAELE FOR LOADING WITH THE ERSIC OBJECT CODE LOADER.

| 00 DATA | 2, 127, 65, E5, E8, 65, 68, 6?, |
| :---: | :---: |
| 1701 DATA | 7009, $78,68,65,83,76,66,67,67,7570$ |
| 1702 DATA | 7608, 66, 67, 83, 66, 69, 81, 65, 73,7573 |
| 1703 DATA | 7015, 84, 66, 77, 73, 65, 78, 69, 66, 7595 |
| 1704 IATA | 7024, 80, 76, 66, $82,75,65,86,67,7622$ |
| 1765 DATA | 7832, 66, 85, 83, 67, 76, 67, 67, 76, 7620 |
| 1706 DRTA | 7040, 68, 67, 76, $73,67,76,85,67$, |
| 1707 DATA | 7043, $77,80,67,80,86,67,80,89,7675$ |
| 170 BATA | 7056, 68, 69, 67, 68, 69, 88, 68, 69, |
| 1709 DATA | 7064, 89, 69, $79,82,73,78,67,73$, |
| 1710 DATA | 7972, $78,88,73,78,89,74,77,80$, |
| 1711 DATA | 7080, 74, 83, 82, 76, 65, 65, 76, 68, |
| 1712 DATA | 7®88, 88, $75,68,89,76,83,82,78$, |
| 1713 DATA | 7096, $79,80,79,82,65,80,72,65,7598$ |
| 4 DATA | 7104, 80, 72, ED, 80, 75, E5, 80, 76, |
| 1715 DATA | 7112, 80, 82, 79, 76, 82, 79, 82, 82, 7754 |
| 1715 DATA | 7120, 84, 73, 82, 84, 83, 83, 66, 67, 7742 |
| 17 DATA | 7128, 83, 69, 67, 83, 69, 68, 83, 69, 7719 |
| 1718 DRTA | 7135, 73, 83, 84, 65, 83, 84, 59, 83, 7779 |
| 19 DATA | 7144, 84, 89, 84, 65, 88, 84, 65, 89, 7792 |
| 1720 DATA | 7152, 84, 83, 88, 84, 88, 65, 84, 88, 781 |
| 1721 DATA | 7160, 83, 84, 89, 65, 84, 69, 88, 255, 79 |
| 1722 DATA | 7168, 34, 106, 1, 1, 1, 106, 10, 1, 7428 |
| 1723 DATA | 7176, 112, 106, 10, 1, 1, 106, 10, 1, 75 |
| 1724 DATA | 7184, 31, 105, 1, 1, 1, 106, 10, 1, 7441 |
| 1725 DATA | 7192, 43, 106, 1, 1, 1, 106, 10, 1, 7461 |
| 1725 DATA | 7200, 58, 7, 1, 1, 22, 7, 121, 1, 7448 |
| 1727 DATA | 7208, 118, 7, 121, 1, 22, 7, 121, 1, 760 |
| 1728 DATA | 7216, 25, 7, 1, 1, 1, 7, 121, 1, |



# Appendix E9： 

## Move Utilities

THE FOLLOWING DATA STATEMENTS
CONTAIN DECIMAL OBJECT CODE AND CHECKSUMS FOR MEMORY FROM 6064 TO 63G9
SUITABLE FOR LOADING WITH
THE RASIC OBJECT CODE LOADER．

| D DAT | 6064，ロ，ロ，ロ，ロ，32，8，20，32，6156 |
| :---: | :---: |
| $1 . \mathrm{DA}$ | E0T2，228，20，127，13，10，32，32，32， 6566 |
| 1862 DATA | E0E0，32，32，77，79，86，69，32，84，6571 |
| 1803 DATA | 6988，79，79，76，46，13，10，10，255， 6656 |
| 1804 DATA | 6056，32，233，21，32，155，24，174，85， 6882 |
| 1805 DATA | 6104，21，56，173，84，21，237，82，21， 6799 |
| 180 E DRIT | 6112，141，176，23，176，2，202，56，138，7026 |
| 1807 DATA | 6120，237，83，21，141，177，23，176，3，E981 |
| 1803 DATA | E128，169，日，96，160，3，185，日，0，6741 |
| 1509 DATA | 6135，72，136，16，249，56，173，83，21， 6942 |
| 1810 DATA | 6144，205，179，23，144，64，200，24，173， 7164 |
| DA | 6152，82，21，205，178，23，144，54，208， 7067 |
| 1212 DAt | 6160，14，160，0，104，153，日，日，200， 6791 |
| DA | 6158，192，4，208，247，169，255，96，32， 7371 |
| 1814 DA | 6176，164，24，160，0，174，177，23，240， 7136 |
| DA | 6154，14，177，0，145，2，200，208，249， 7179 |
| 1815 DA | 6192，230，1，230，3，202，208，242，136， 7444 |
| DA | 6200，200，177，0，145，2，204，176，23， 7127 |
| g | 6208，208，246，76，17，24，173，177，23， 7152 |
| 1819 DATR | 6216，240，72，172，177，23，173，176，23， 7272 |
| 182 L DATA | 6224，56，233，255，176，1，136，170，132， 7383 |
| 1821 DA | 6232，3，138，24，109，82，21，133，0， 6742 |
| 1822 DATA | 6240，144，1，200，152，105，83，21，133， 7083 |
| 1823 D | 6248，1，138，24，109，178，23，133，2， 6856 |
| 1824 DATA | 6256，144，2，230，3，165，3，109，179， 7091 |
| 1825 DATA | 6264，23，133，3，174，177，23，160，255， 7212 |
| 1826 DATA | E272，177，日，145，2，136，208，249，177， 7366 |
| 1827 DATA | 6280，0，145，2，198，1，198，3，202， 7029 |
| 1828 DATA | 6288，208，236，32，164，24，172，176，23， 732 |


| 1829 DATA | 6236，177，0，145，2，136，192，255，208， 7411 |
| :---: | :---: |
| 1830 DATA | 6304，247，76，17，24，173，82，21，133， 7077 |
| 1831 DATA | 6312，0，173，83，21，133，1，173，178， 7074 |
| 1832 DATA | 6329，23，133，2，173，179，23，133，3， 6989 |
| 1833 DATA | 6328，95，32，8，20，32，228，20，127， 6891 |
| 1834 DATA | 6336，13，10，83，69，84，32，68，69， 6764 |
| 1835 DATA | 6344，83，84，73，78，65，84，73，79， 6963 |
| 1836 DATA | $6352,78,32,65,78,68,32,80,82,6867$ |
| 1837 DRTA | 6360，59，83，83，32，81，46，255，32， 7041 |
| 1835 DATA | 6368， $7,18,173,5,18,141,178,23,6531$ |
| 1839 DATA | 6376，173，E，18，141，179，23，96，0， 7012 |
| 1840 DATA | E3B4，ロ，日，ロ，ロ，D，D，ロ， 0,6334 |
| 1841 DATA | 6392，ロ，ロ，ロ，ロ，日，ロ，ロ，ロ， 6392 |
| 1842 END |  |

# Appendix EIO: 

## Simple Text Editor

APPENDIX EIO
A SIMPLE TEXT EDITOR

THE FOLLONING DATA STATEMENTS
CONTAIH DECIMAL OBJECT CODE AND
CHECKSUMS FOR MEMORY FROM 7680 TO 8191
SUITABLE FOR LOADING WITH
THE ERSIC OBJECT CODE LOADER.

| Data | 7680, 255, 1, 32, 15, 30, 32, 55, 30, 8130 |
| :---: | :---: |
| 1 DATA | 7688, 32, 200, 30, 24, 24, 144, 246, 32, 8420 |
| 1902 DATA | 7696, 8, 20, 32, 228, 20, 127, 13, 10, 8154 |
| 1583 DATA | 7704, 10, 83, 69, 84, 32, 85, 80, 32, 8179 |
| 1904 DATA | 7712, 59, 68, 73, 84, 32, 66, 85, 70, 8259 |
| 1905 BATA | 7720, 70, 69, 82, 46, 13, 10, 10, 255, 8275 |
| 1905 DATA | 7728, 32, 233, 21, 32, 160, 23, 96, 32, 8357 |
| 1907 DATA | 7736, 196, 17, 32, 43, 17, 174, 3, 16, 8234 |
| 1908 DATA | 7744, 160, 3, 32, 19, 17, 32, 43, 17, 8667 |
| 1999 DATA | 7752, 32, 118, 17, 32, 196, 17, 32, 34, 8290 |
| 1910 DATA | 7760, 30, 32, 211, 17, 32, 118, 17, 32, 8249 |
| 1911 DATA | 7750, 137, 30, 32, 211, 17, 96, 32, 18, 8341 |
| 1912 DATA | 7776, 21, 173, 3, 16, 74, 170, 202, 202, 8637 |
| 1913 DAT | 7784, 32, 25, 19, 202, 16, 259, 173, 3, 8505 |
| 1914 Data | 7792, 16, 141, 0, 30, 32, 148, 18, 32, 8209 |
| 1915 DATA | 7800, 155, 17, 32, 127, 17, 32, 13, 19, 8212 |
| 1915 DATA | 7806, 206, 0, 30, 16, 239, 32, 43, 21, 8355 |
| 1917 DAT | 7816, 96, 173, 3, 16, 74, 233, 2, 32, 8445 |
| 1918 DATA | 7824, 129, 17, 173, 1, 30, 201, 1, 208, 8584 |
| 1919 DATA | 7832, 5, 165, 73, 24, 144, 2, 169, 79, 8497 |
| 20 data | 7840, 32, 155, 17, 159, 2, 32, 129, 17, 8393 |
| 1921 DATA | 7848, 173, 7, 16, $32,155,17,169,2,8419$ |
| 1922 DATA | 7856, 32, 129, 17, 173, 6, 18, 32, 163, 8426 |
| 1923 DATA | 7864, 17, 173, 5, 18, 32, 163, 17, 96, 8385 |
| 1924 DATA | 7872, 6, 3, 62, ED, 16, 127, 81, 日, 8227 |
| 925 DATA | 7880, 32, 224, 18, 205, 198, 30, 208, 23, 8816 |
| 1925 DATA | 7888, 72, 32, 224, 18, 205, 152, 30, 208, 8675 |
| 1927 DRTA | 7696, 4, 104, 104, 104, 96, 141, 199, 30, 8678 |
| 1928 DATA | 7904, 104, 32. 231, 30, 173, 199, 30, 205, 8908 |



# Appendix EII： 

## Extending the Visible Monitor

APFENDIX E11

EXTENDING THE UISIBLE MONITOR

THE FOLLONING DATA STATEMENTS CONTAIN DECIMAL OBJECT CODE AND CHECKSUMS FOR MEMORY FROM 4272 TO 4351 SUITAELE FOR LOADING WITH THE EASIC OBJECT CODE LOADER．

| 260日 Data | 4272，201，60，208，9，173，0，20，73， 5036 |
| :---: | :---: |
| 2001 DRTA | 4230，255，141，0，20，96，201，85，208， 5286 |
| 2002 DATA | 4288，9，173，2，20，73，255，141，2， 4963 |
| 2093 DATA | 4296，20，96，201，72，208，13，173，0， 5079 |
| 2004 DATA | 4304，20，208，4，32，87，21，96，32，4804 |
| 2005 DATA | 4312，174，21，55，201，77，200，4，32，5125 |
| 2006 DATA | 4320，180，23，96，201，63，208，13，173， 5277 |
| 2007 DATA | 4328，0，20，208，4，32，9，25，56，4722 |
| 2008 DATA | 4336，32，38，25，56，201，84，208，4，5024 |
| 2003 DATA | 4344，32，2，30，96，36，日，日，0，4500 |
| 2010 ETH |  |

# Appendix EI2: 

# System Data Block for the Ohio Scientific C-IP 

fPFENDIX E1Z SYSTEM DATA BLOCK FOR OSI CIP

```
THE FOLLOWING DATA STATEMENTS COMTAIN DECIMAL OEJECT CODE RND CHECKSUMS FOR MEHORY FROM 4036 TO 4115 SUITABLE FOR LOADING WITH THE EASIC OBJECT CODE LOADER.
```

```
210g DATA 4096, 101, 208, 32, 24, 24, 211, 32, 16, 4744
2101 DATA 4104, 237, 254, 45, 191, 177, 252, 15, 16, 5292
2102 DATf
2103 END
```

OK

# Appendix EI3: <br> <br> System Data Block for the PET 2001 

 <br> <br> System Data Block for the PET 2001}

APPENDIX E13
SYSTEM DATA RLOCK FOR THE PET $2 B 01$

THE FOLLOWING DATA STATEMENTS CONTAIN DECIMRL OBJECT CODE AIND CHECKSUMS FOR MEMORY FROM 4035 TO 4151 SUITAELE FOR LOADING WITH THE BASIC OBJECT CODE LOADER.

```
2100 DATA 4096, 0. 128, 40, 39, 24, 131, 32, 30, 4520
2101 DATA 4104, 42, 15, 210, 255, 15, 16, 16, 15, 4591
2102 BATA 4112, 95, 41, 127, 56, 201, 54, 144, 17, 4858
2103 DATA 4120, 201, 96, 144, 10, 162, 14, 141, 76, 4964
2104 DATA 4128, 232, 233, 32, 24, 144, 3, 55, 233, 5085
2105 DATA 4136, 64, 96, 32, 228, 255, 41, 127, 240, 5219
210E DRTA 4144, 24S, S6, 0, 0, , , , 0, 0, 4483
Z107 END
```

OK

## Appendix El4:

## System Data Block for the Apple II

RPFENIIIX E14<br>SYSTEM DATA ELOCK FOR THE RPFLE II

THE FOLLONING DATA STATEMENTS CONTAIN DECIMAL OBJECT CODE AND CHECKSUMS FOR MEMORY FROM 4696 TO 4127 SUITHELE FOR LOFIDING WITH the besid ogject code londer.

```
2100 DATA 4095, 0, 4, 128, 39, 7, 7, 160, 222, 4663
2101 DATA 4104, 29, 16, 25, 16, 16, 15, 16, 15, 4246
Z1日2 DATA 4112, 95, 9, 120, 95, 32, 12, 253, 41, 4779
2103 DATA 4120, 127, 95, 9, 128, 32, 253, 251, 96,5112
```

OK

## Appendix EI5： <br> System Data Block for the Atari 800

APFENDIX E1S
SYSTEM DATA BLOCK FOR THE ATARI BDU

THE FOLLOWING UATA STATEMENTS
CONTAIN DECIMAL OBJECT CODE AND CHECKSUMS FOR MEMORY FROM 3712 TO 4223 SUITABLE FOR LOADING WITH THE BASIC OBJECT CODE LORDER．
2100 DATA
2101 DATA
2102 DATA
2103 DATA
2104 DATA
2105 DATA
2106 DATA
2107 DATA
2108 DATA
2109 DATA
2118 DATA
2111 DATA
2112 DATA
2113 DATA
2114 DATA
2115 DATA
2116 DATA
2117 DATA
2118 DATA
2119 DATA
2120 DATA
2121 DATA
2122 DATA
2123 DATA
2124 DATA
2125 DATA
2126 DATA
2127 DATA
2128 DATA

3712，32，196，17，173，179，23，72，173， 4577 3720，178，23，72，173，85，21，72，173， 4517 $3728,84,21,72,173,83,21,72,173,4427$ $3736,82,21,72,32,43,17,165,0,4168$ $3744,141,178,23,165,1,141,179,23,4595$ $3752,32,118,17,165,0,141,82,21,4328$ 3760，165， $1,141,83,21,174,3,16,4364$ $3768,172,4,16,32,60,17,165,0,4234$ 3775，141，84，21，165，1，141，85，21， 4435 $3784,32,214,23,172,4,16,162,6,4407$ $3792,32,60,17,174,3,16,160,1,4255$ $3800,32,19,17,104,141,82,21,104,4320$ 3808．141，83，21，104，141，84，21，104，4507 $3816,141,85,21,104,141,178,23,104,4613$ 3824，141，179，23，32，211，17，96，0， 4523 3832，日，日，日，日，日，日，ロ，ロ， 3832
$3840,108,106,59,0,0,107,43,42,4305$ $3848,111,0,112,117,13,105,45,61,4412$ 3856，118，日，99，日，0，98，120，122， 4413 3864,52, 日， $51,54,27,53,50,49,4200$ 3872，44，32，46，110，0，109，47，0， 4260 3880，114，0，101，121，9，116，119，113， 4573 3888，57， $0,48,55,8,56,60,52,4234$ 3896，102，104，100，日，日，103，115，97， 4517 3904，76，74，58，0，0，75， $91,94,4372$ $3912,79,8,80,85,13,73,45,61,4348$ 3920，86，0，67，0，0，66，88，90，4317 $3928,52,0,51,54,27,37,34,33,4216$ 3936．90，32，53，78，0，77，63，0， 4369


OK

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# Beyond Games: Systems Software for Your 6502 Personal Computer 

By Ken Skier

Use your 6502 personal computer for more than games! Learn how it works and how to make it work for you. This book, for Apple, Atari, Ohio Scientific and PET computer owners who know little or nothing about bits, bytes, hardware, and software, presents a guided tour of your computer. Beginning with basic concepts such as what is memory? and what is a program?, Beyond Games moves through a fast but surprisingly complete course in assembly language programming. Having mastered these fundamentals, the reader is introduced to many useful subroutines and programming tools, such as screen utilities, print utilities, a machine language monitor, a hexadecimal dump tool, a move tool, a disassembler, and a simple, screen-based text editor.

## About the Author

Ken Skier, systems analyst for Wang Laboratories, Inc, designs software for word processing and other applications concerning the office of the future. A Massachusetts Institute of Technology graduate, he co-founded the M.I.T. Writing Program, where he teaches science fiction writing. He lives in Cambridge, Massachusetts, with his wife Cynthia and a nameless white cat.


[^0]:    *Butterfield, et al, The First Book of Kim, Rochelle Park, NJ: Hayden Book Company, 1977.

